

Lecture 9

Augmenting Trees (contd.)

Finding the Element with i th Rank

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To find the element with i th rank in T call **Select**($T.\text{root}, i$).

Finding the Element with i th Rank

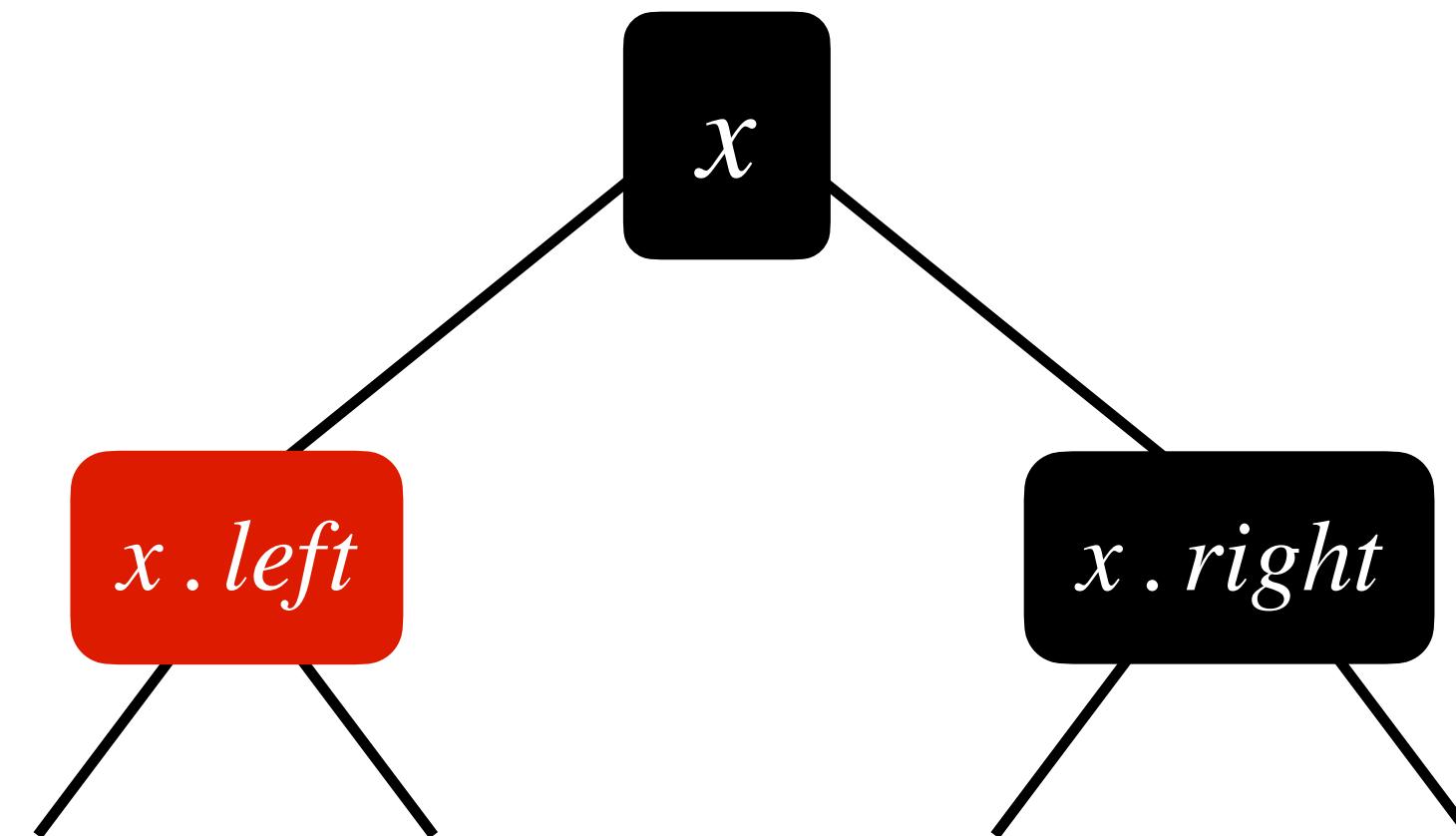
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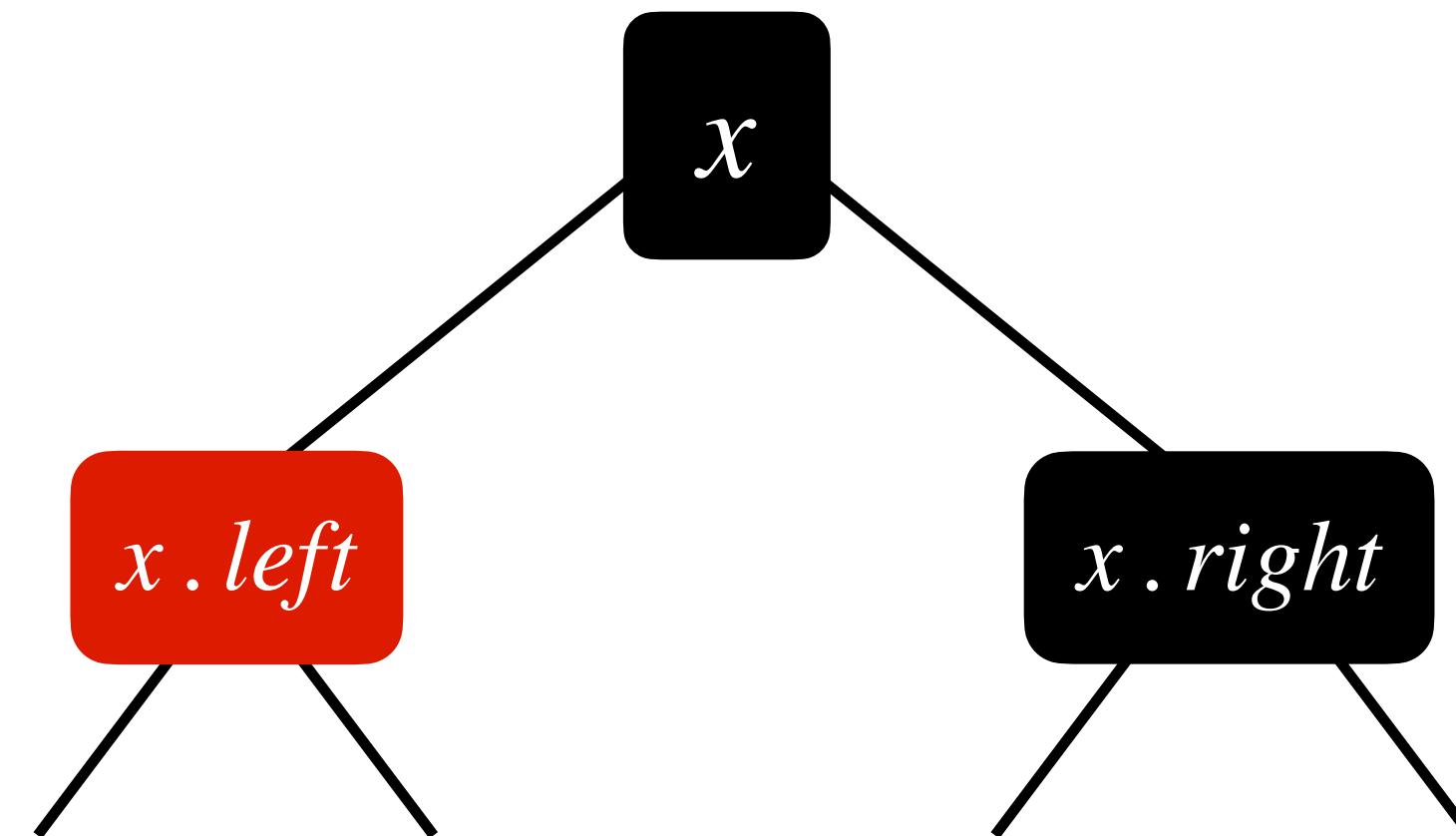


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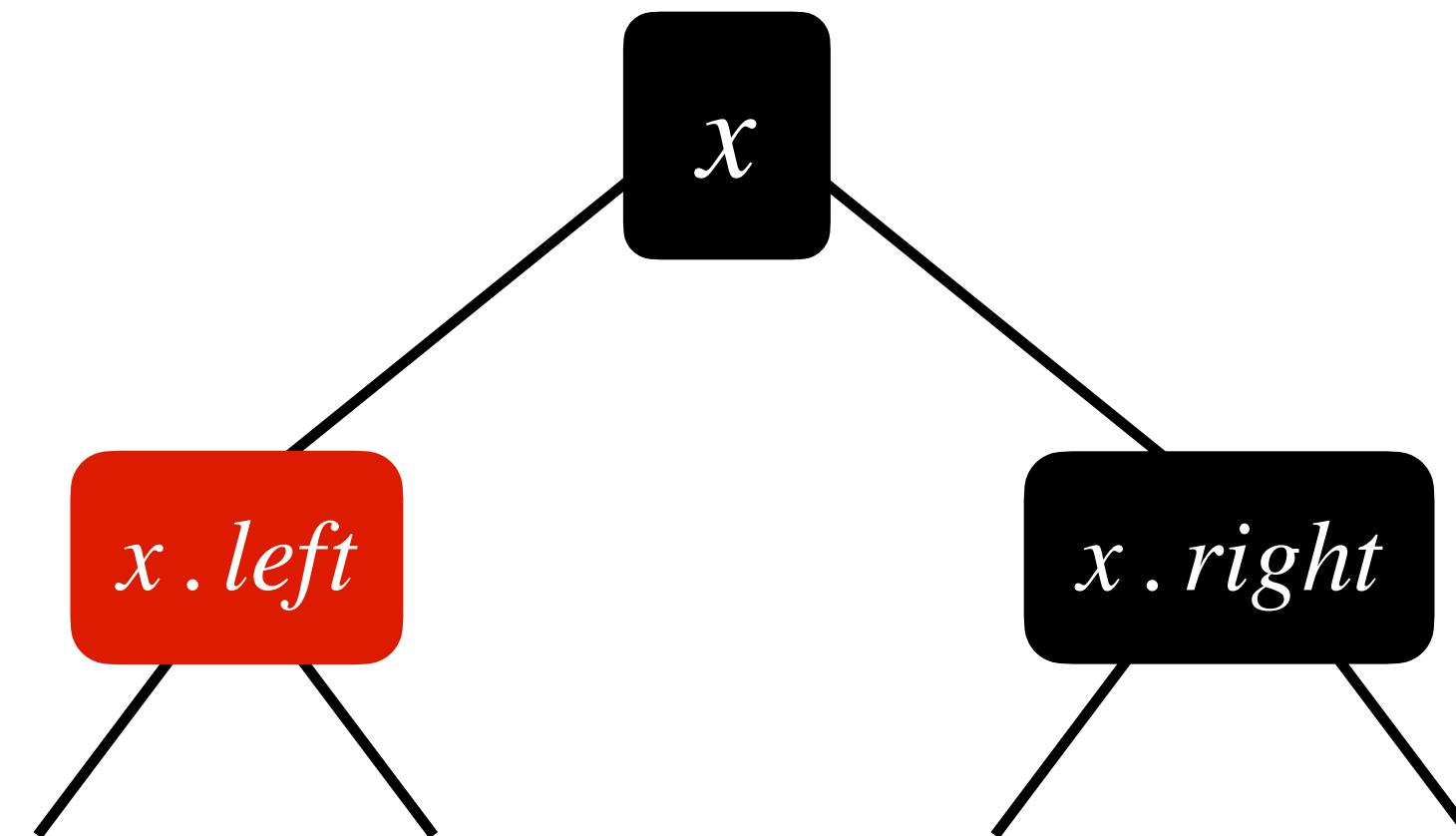
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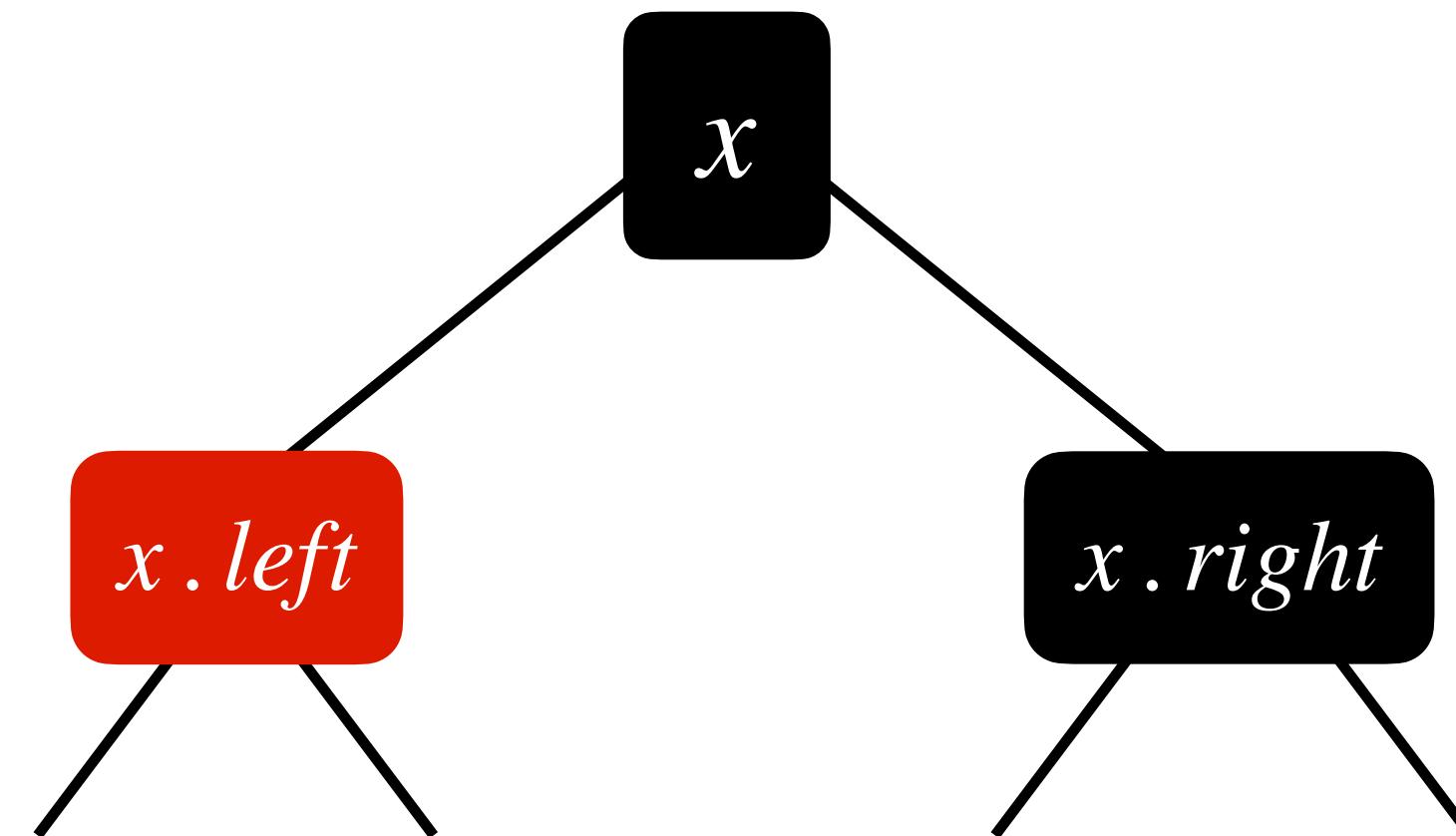
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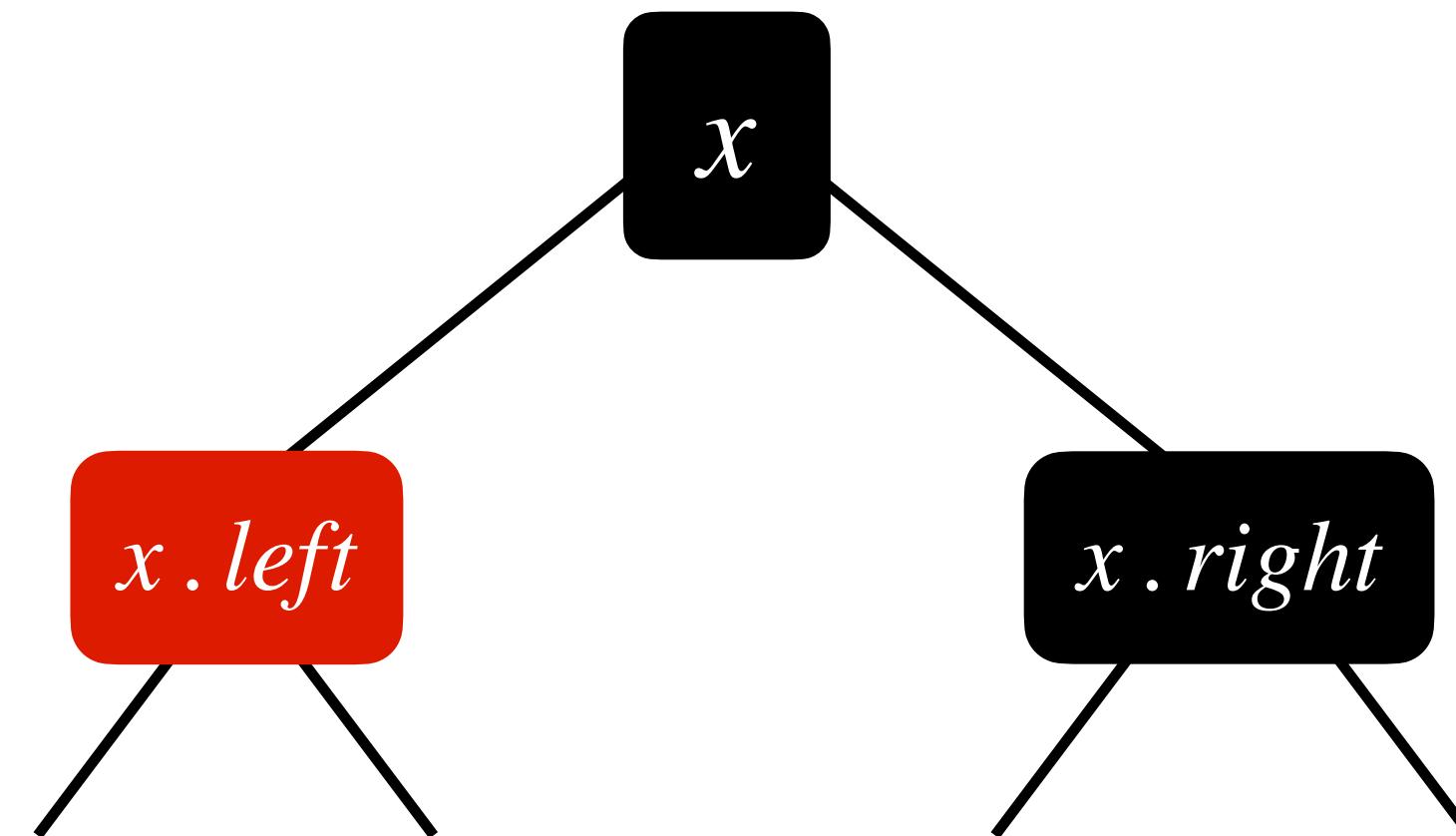


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2. if $i == r$



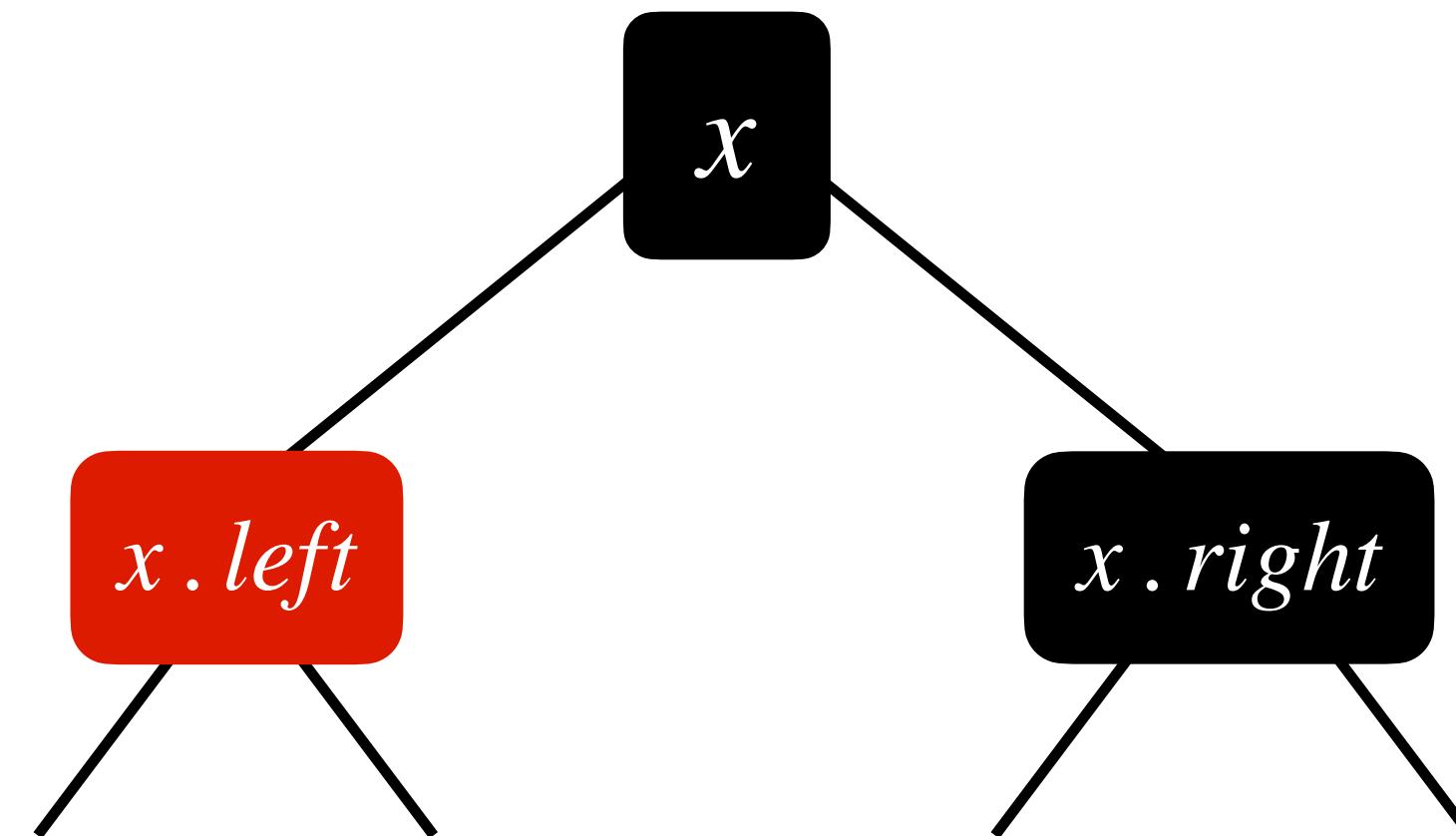
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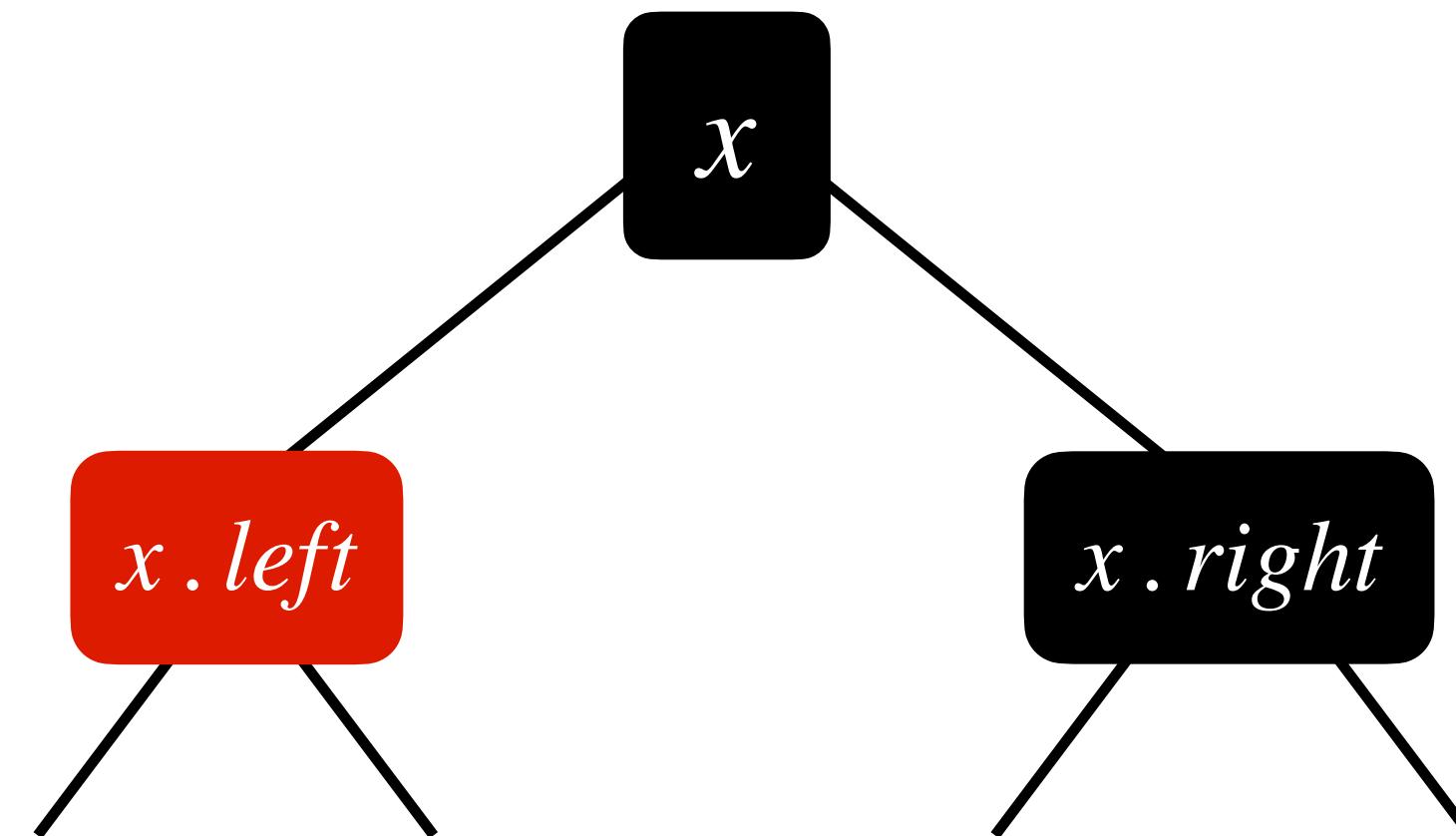
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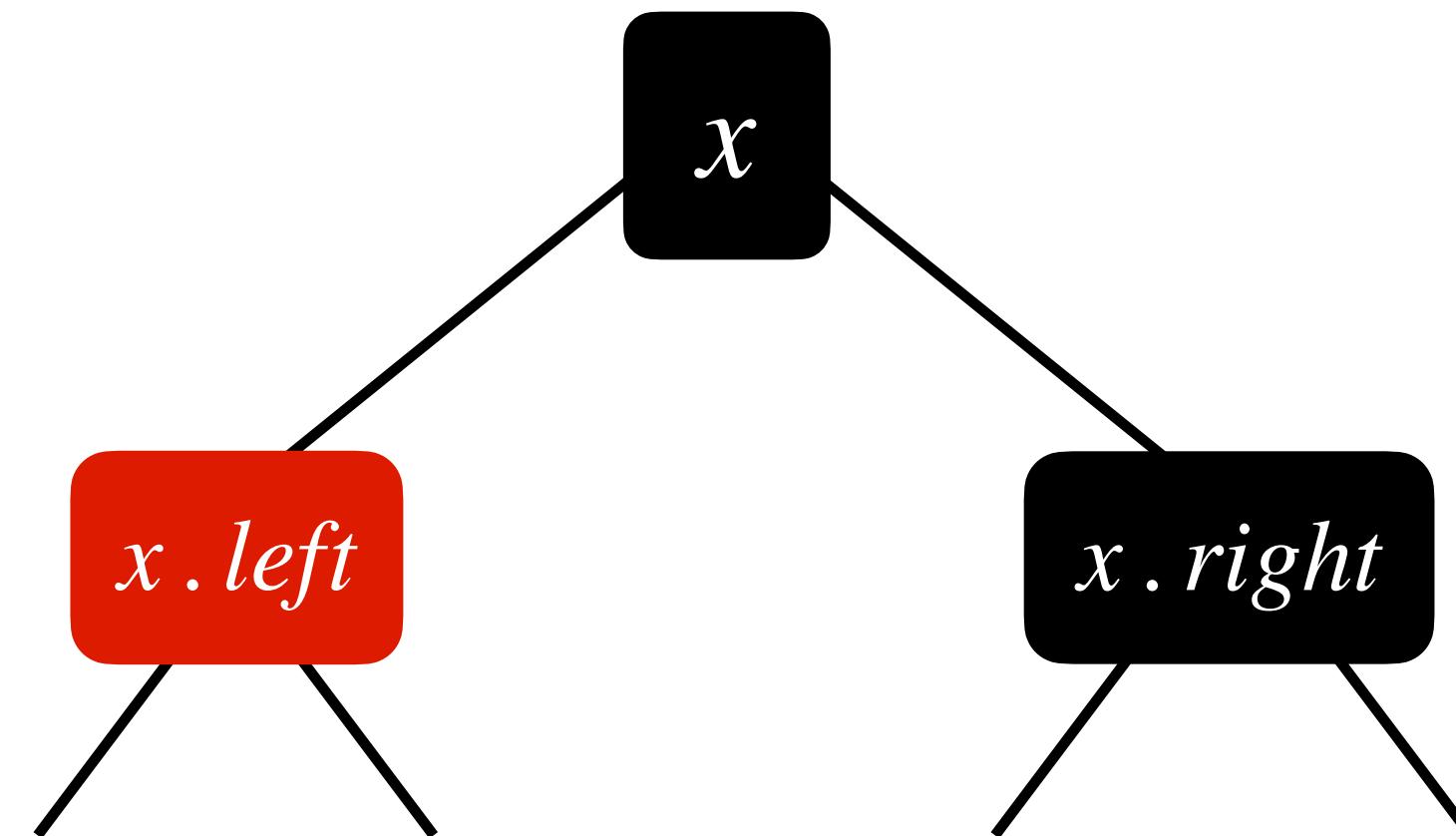


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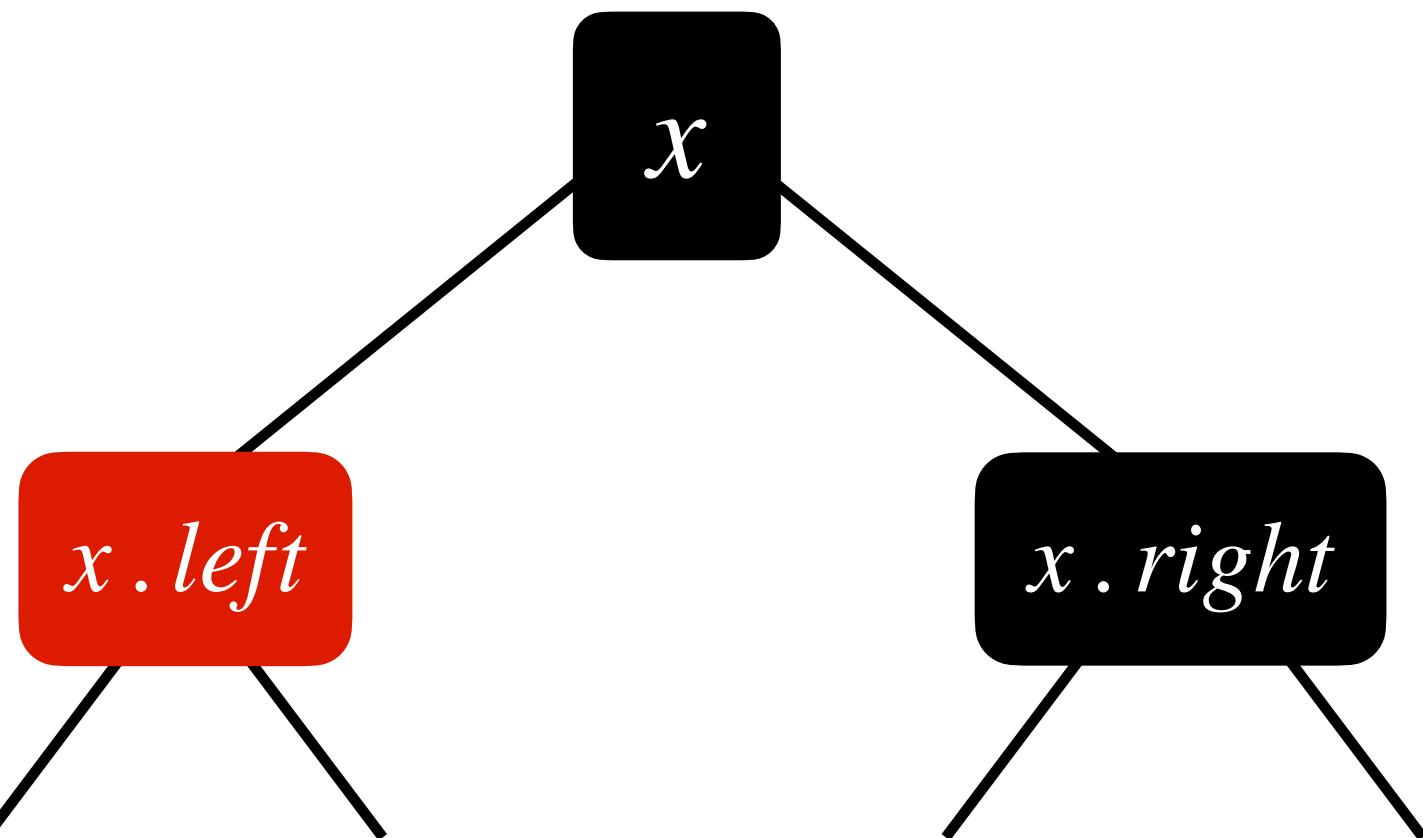
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Element with i th rank
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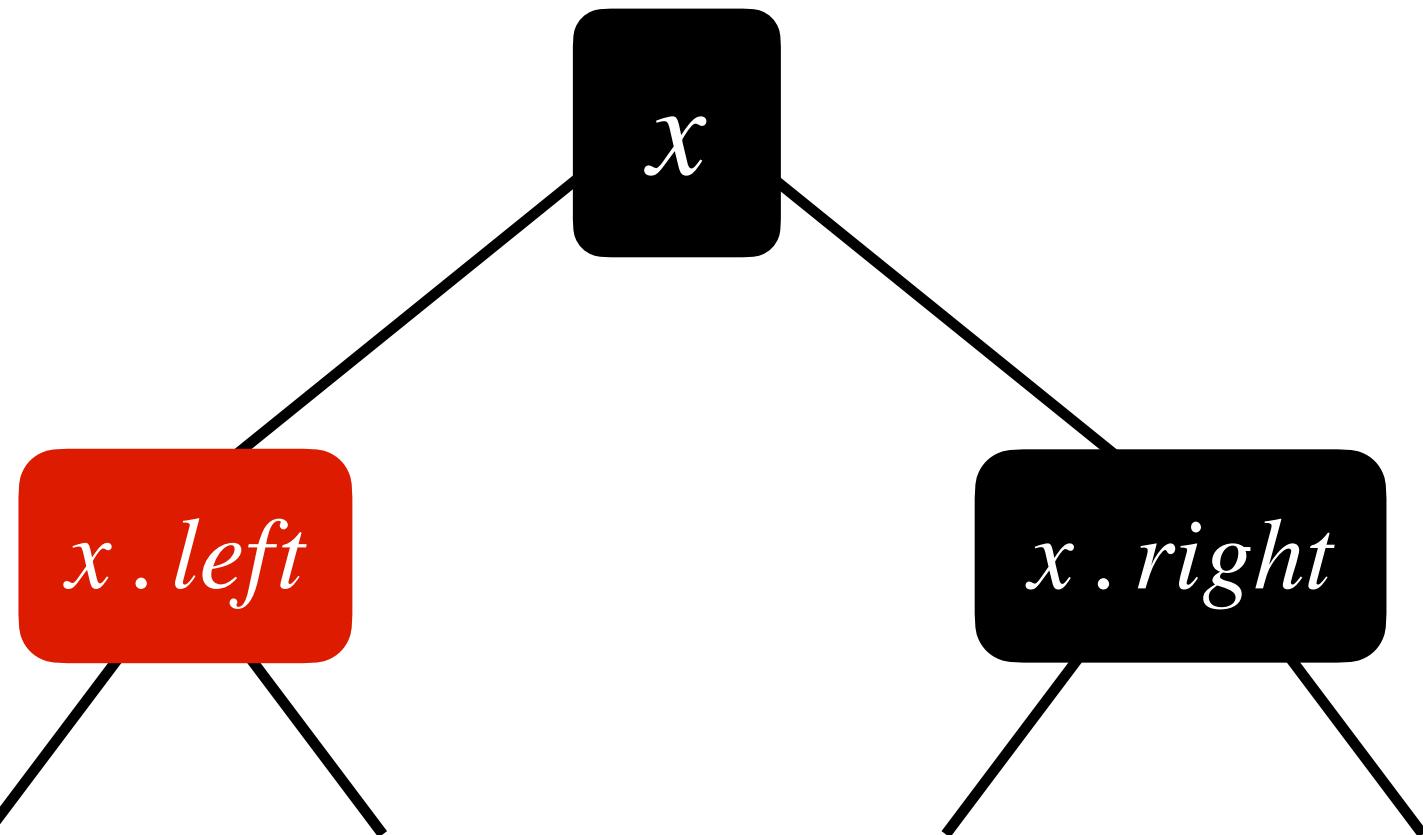
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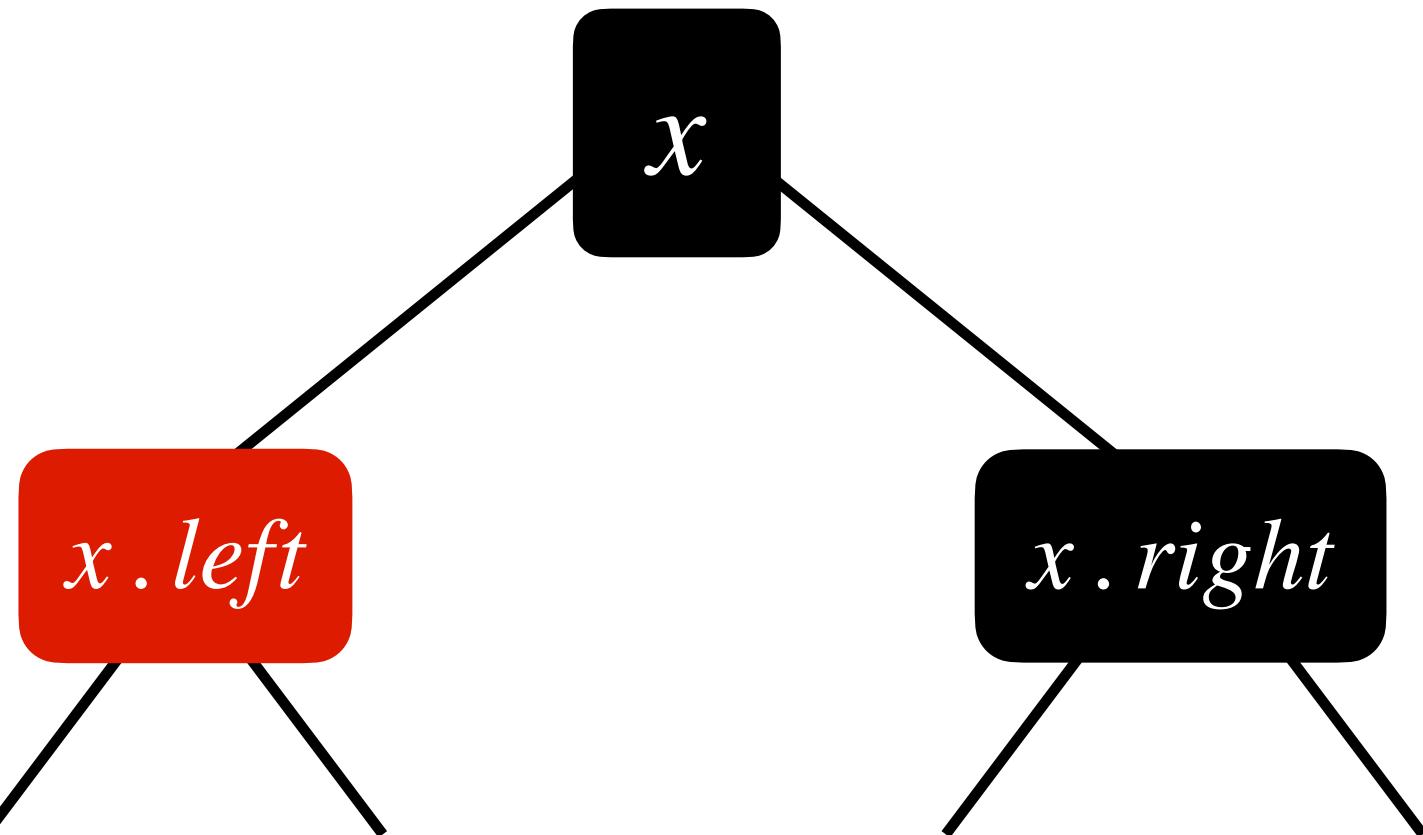
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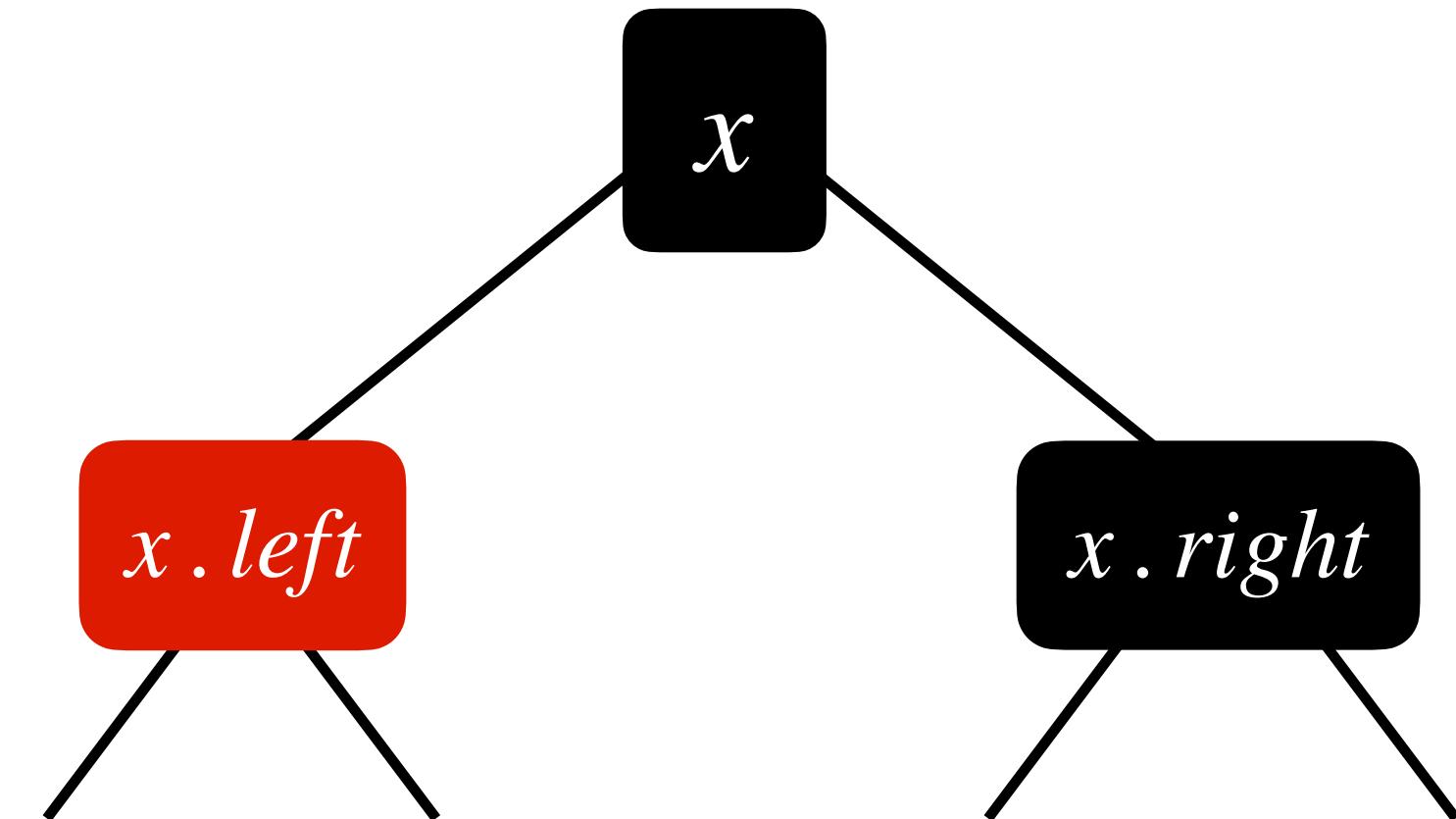


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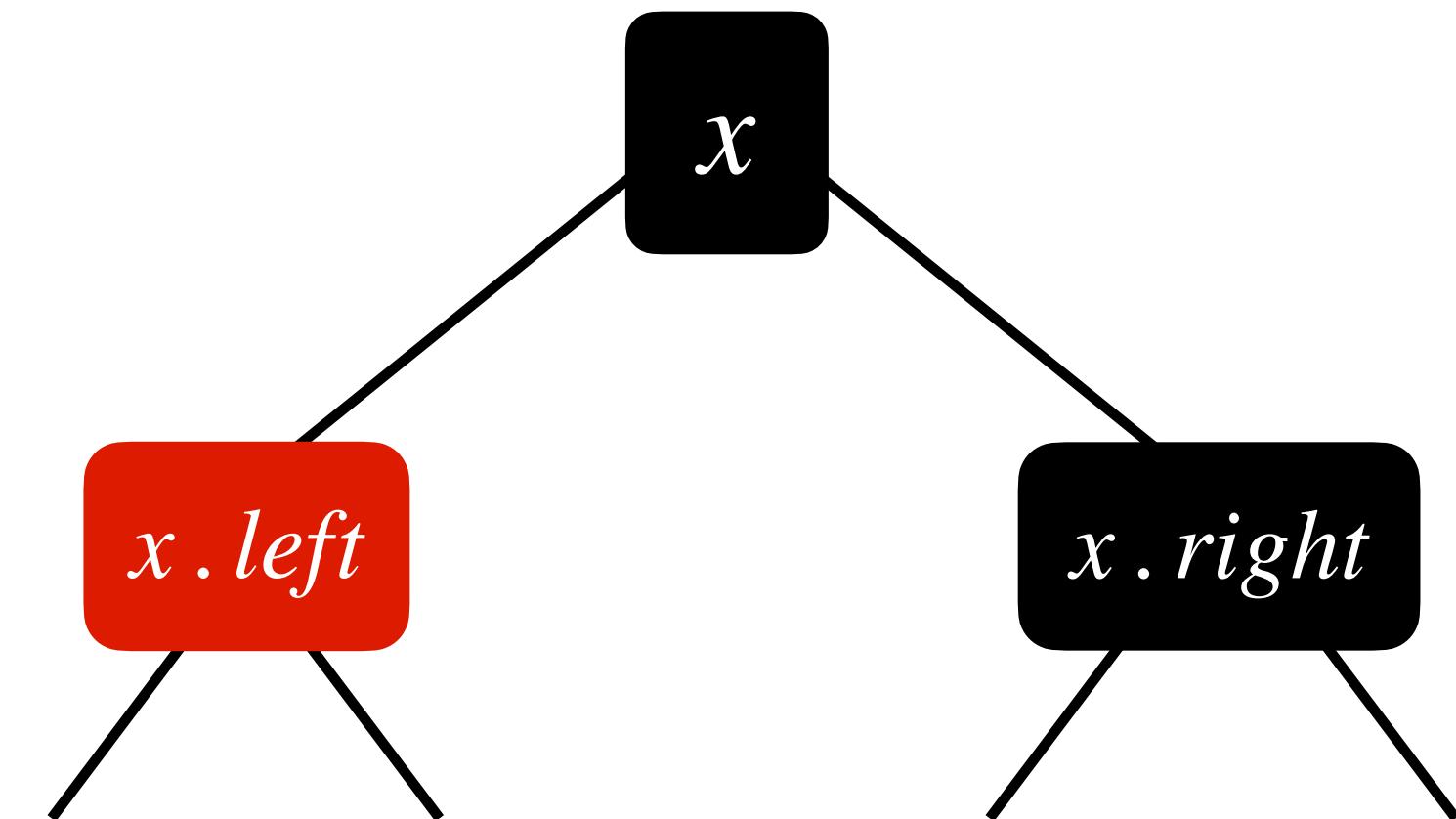
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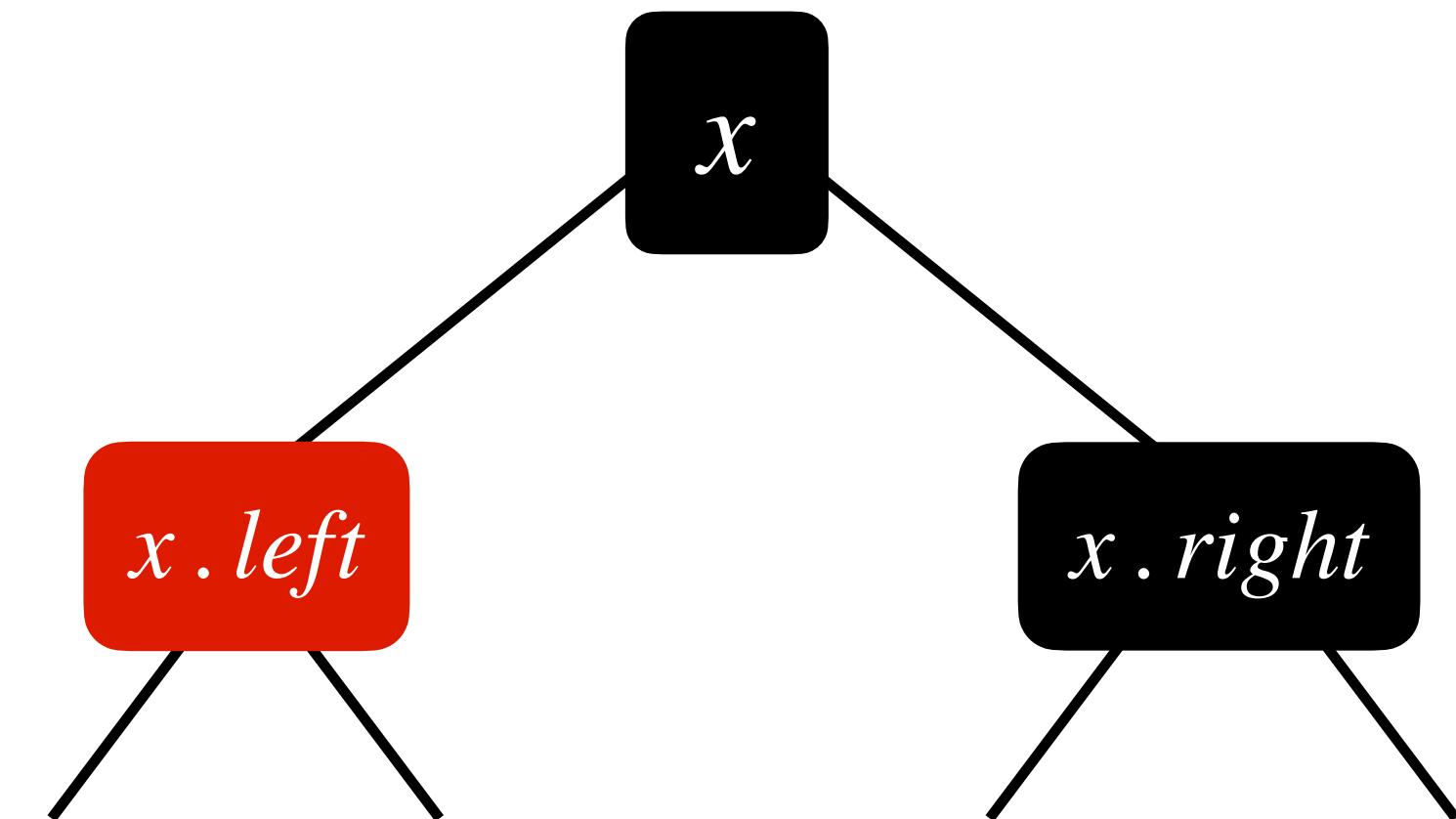
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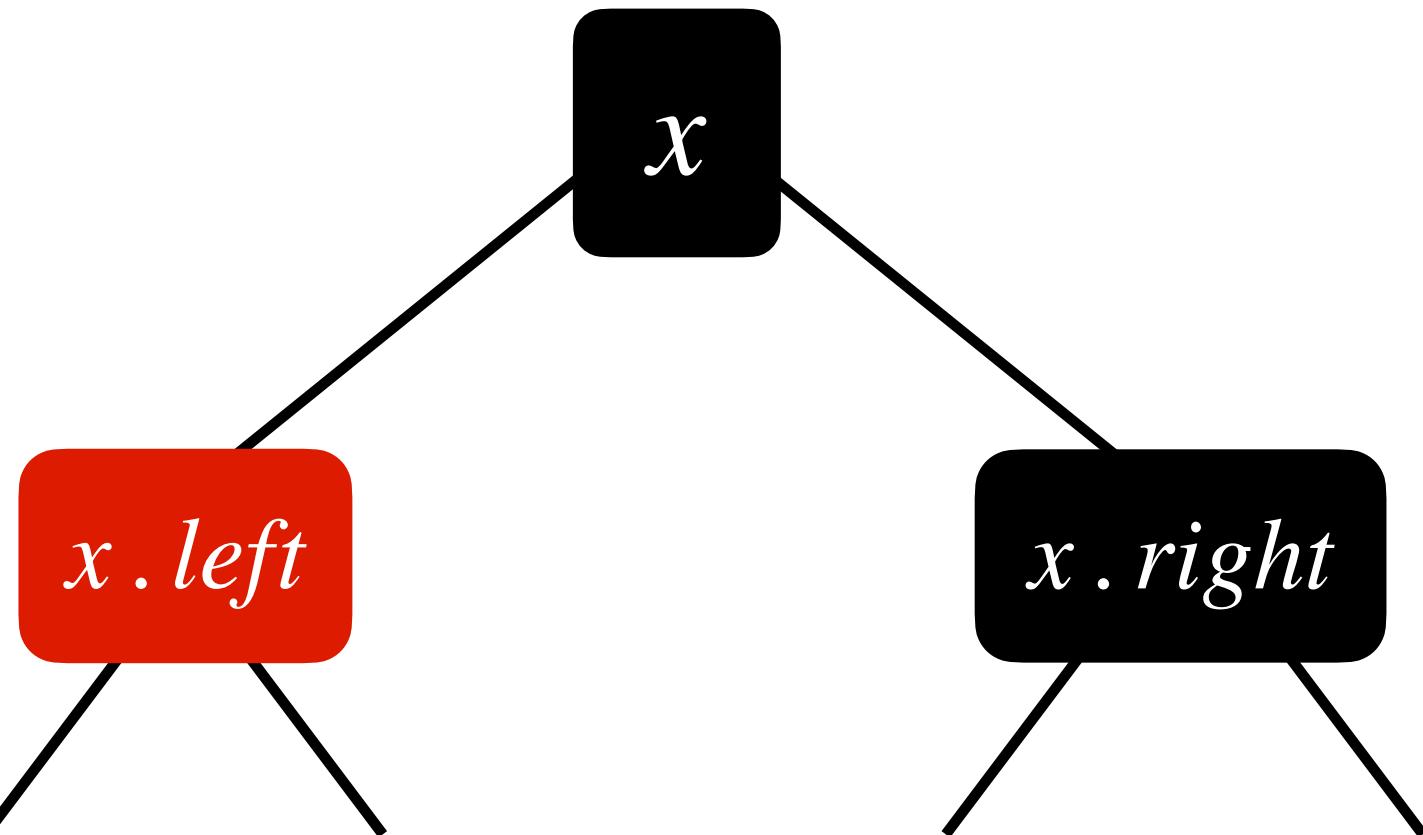
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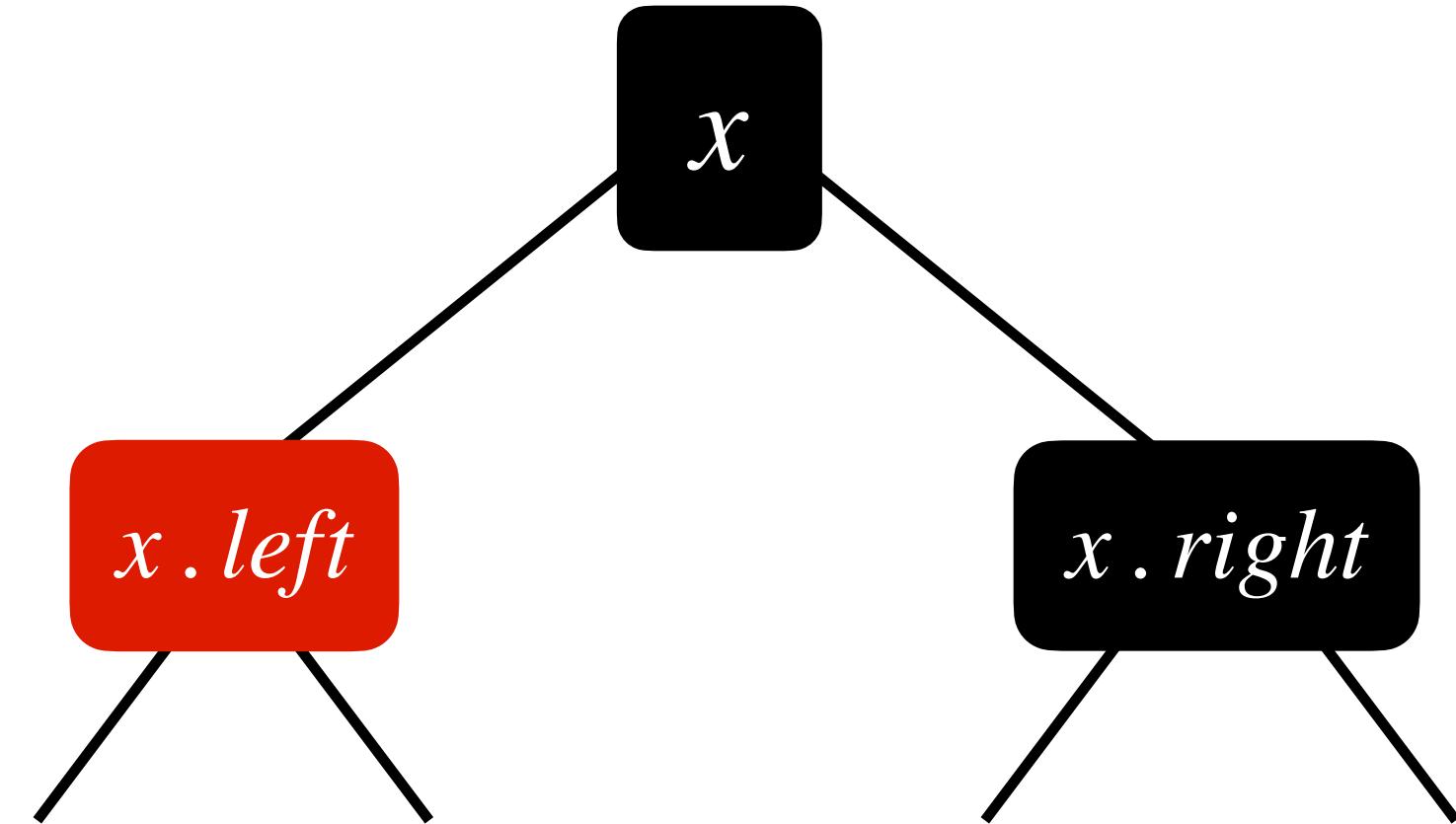


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4. **else if** $i < r$
5. **return** **Select**($x.\text{left}$, i)
6. **else**
7. **return** **Select**($x.\text{right}$, $i - r$)

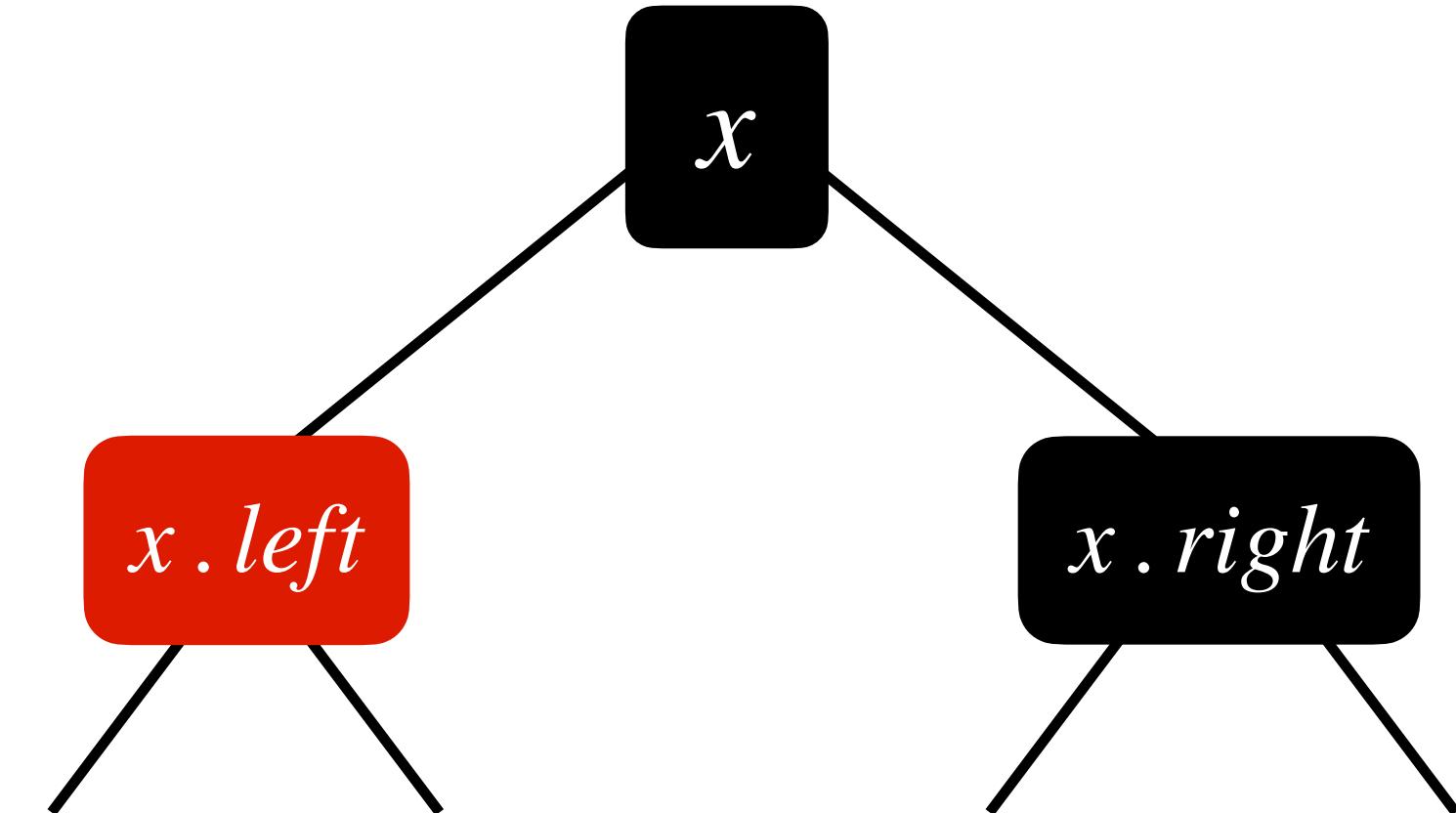


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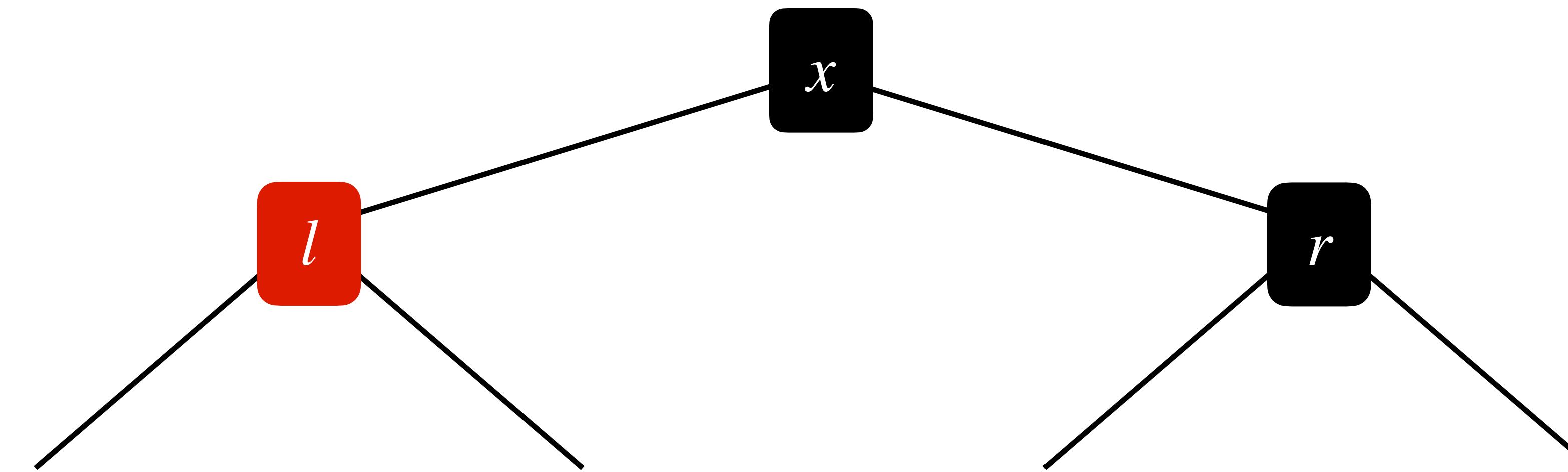
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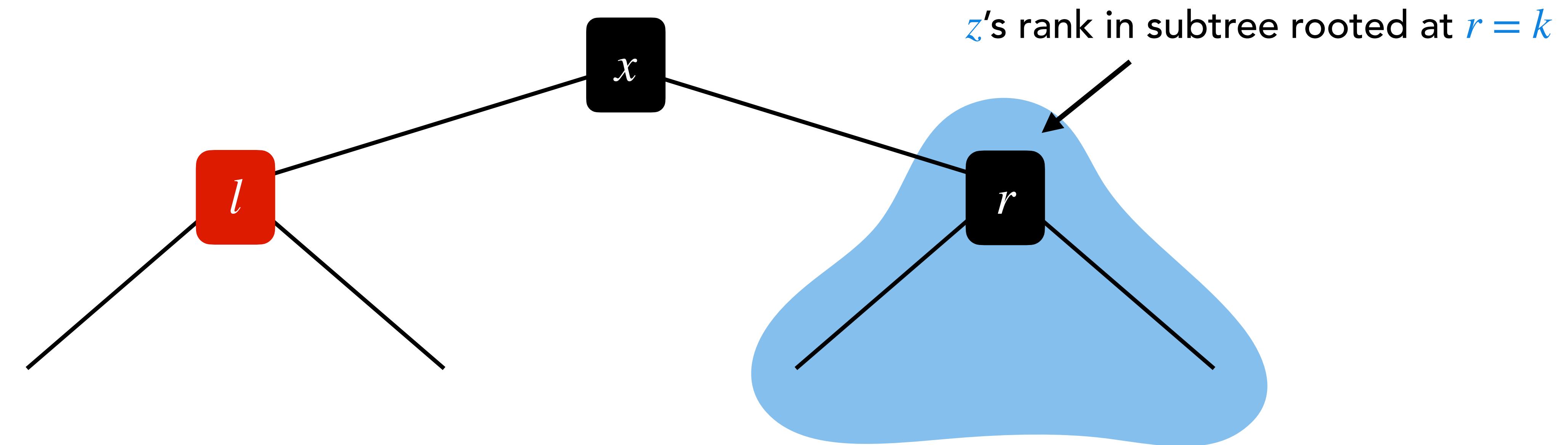


Time Complexity: $\Theta(h) = \Theta(\log n)$ as with every recursive call algorithm goes one level down.

Finding The Rank of an Element

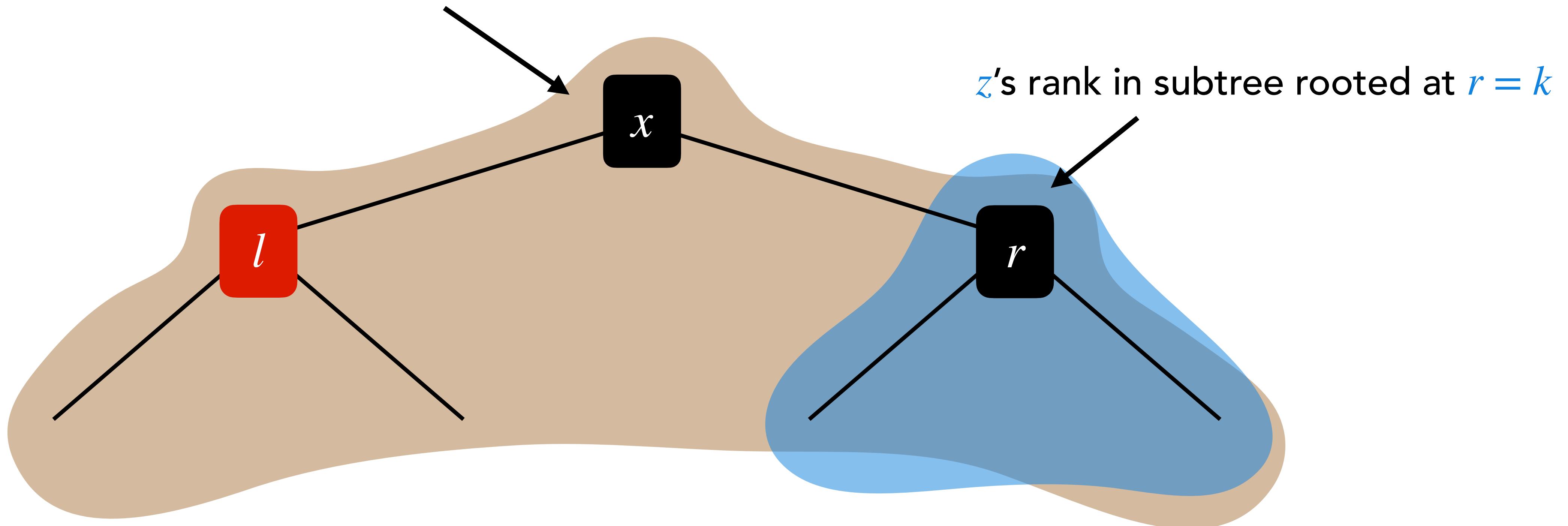


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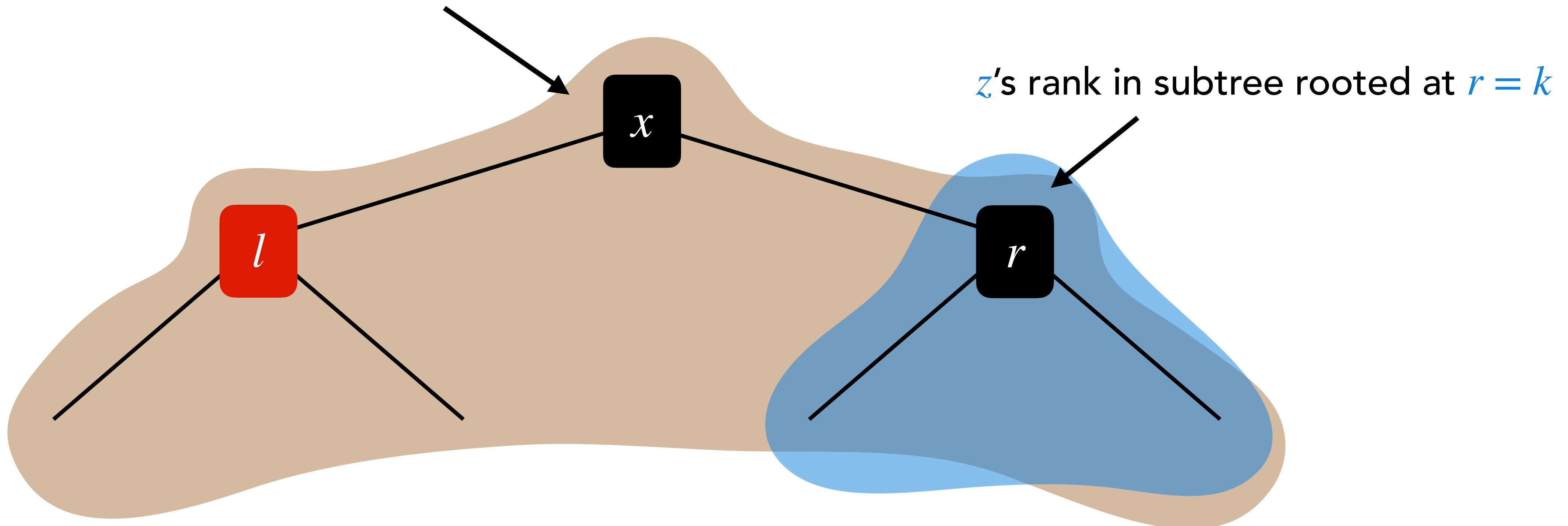
z 's rank in subtree rooted at $x =$



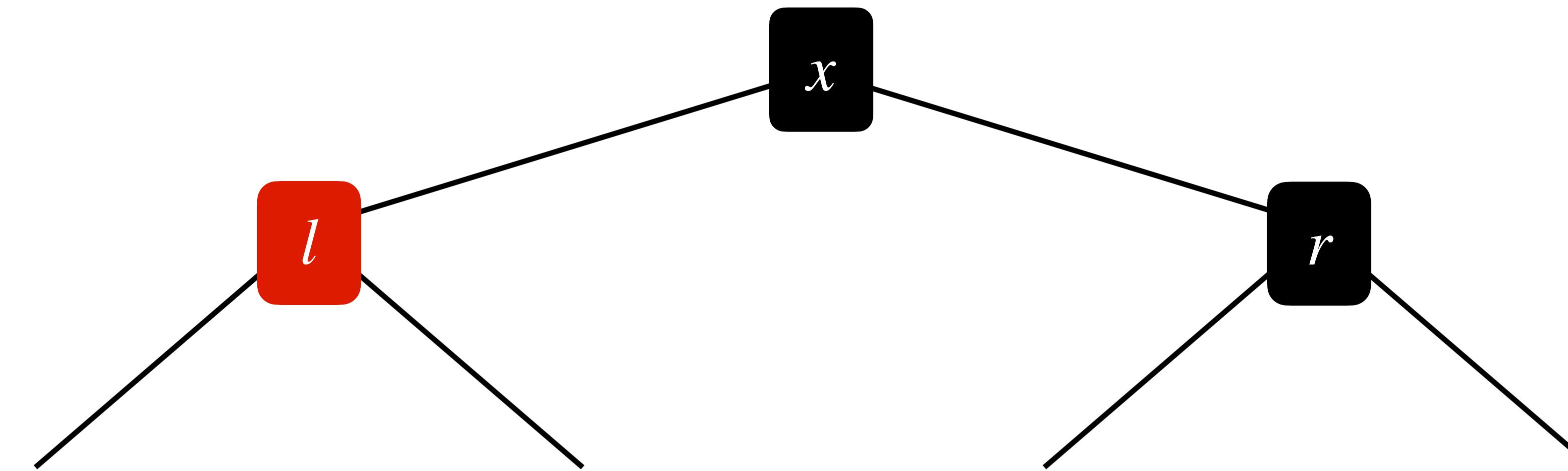
z 's rank in subtree rooted at $r = k$

Finding The Rank of an Element

z 's rank in subtree rooted at $x = k + l.size + 1$

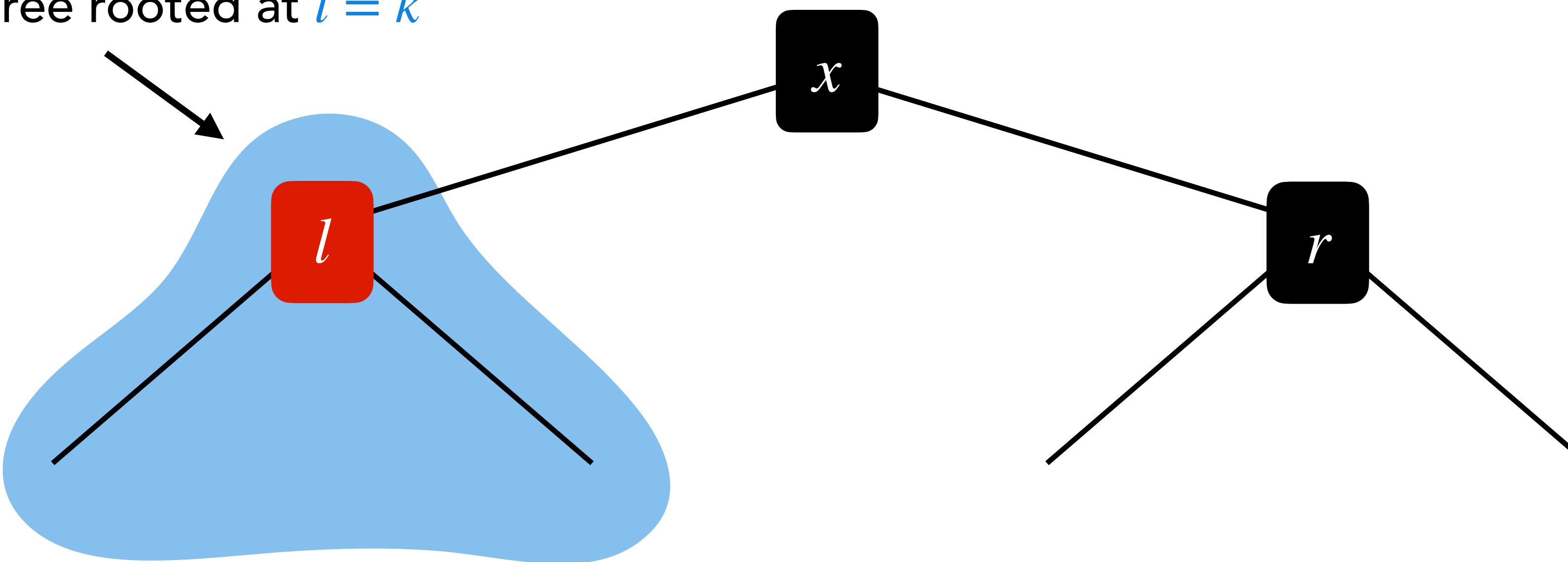


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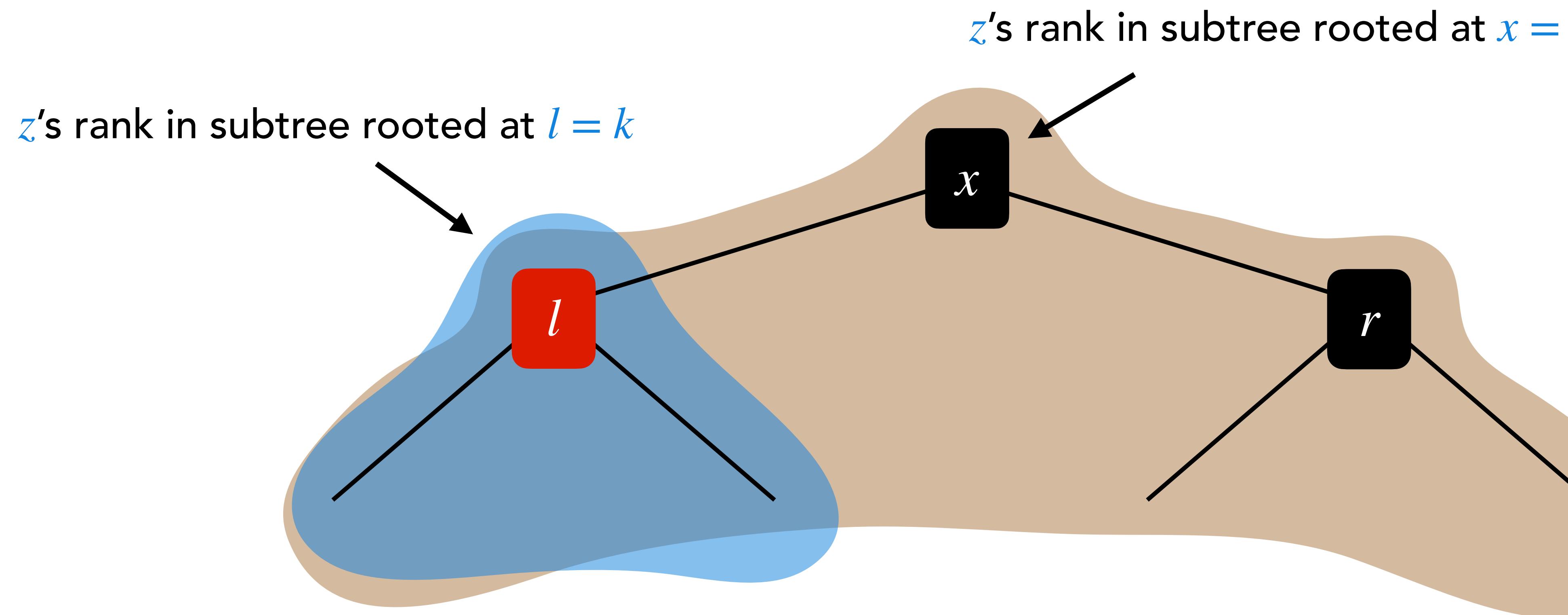


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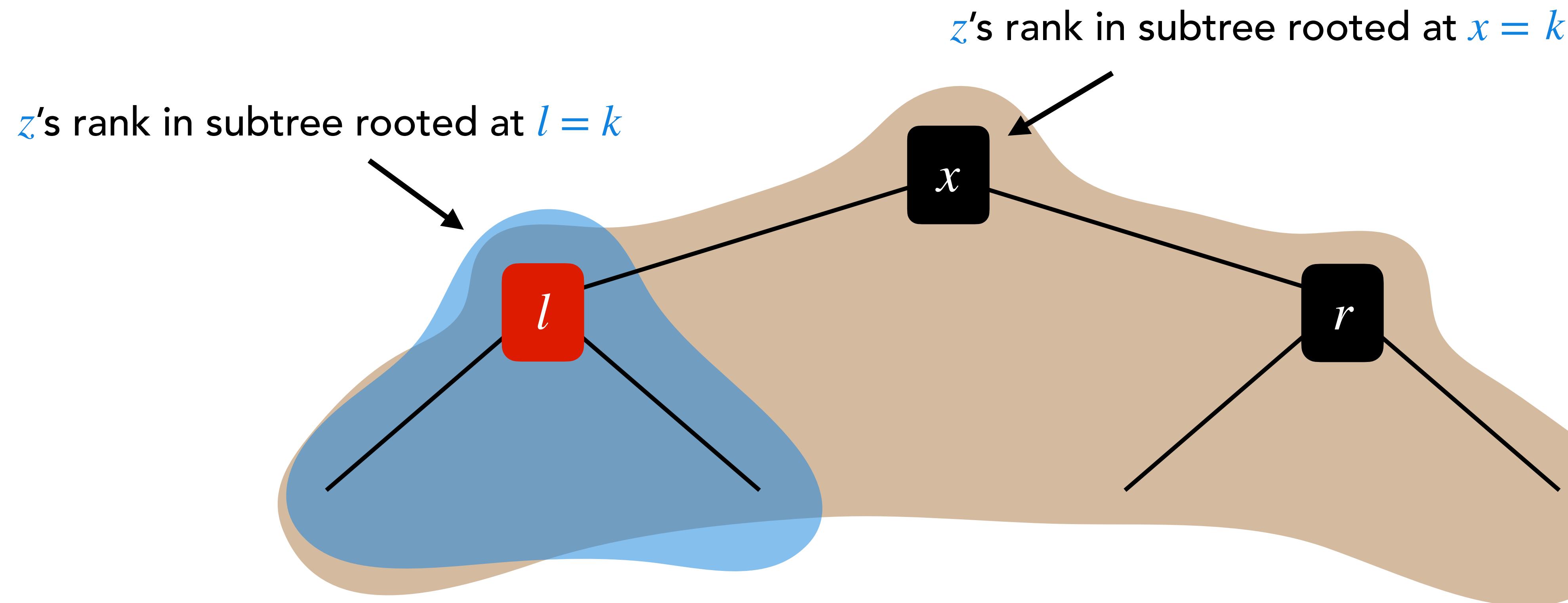
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Finding The Rank of an Element



Finding The Rank of an Element



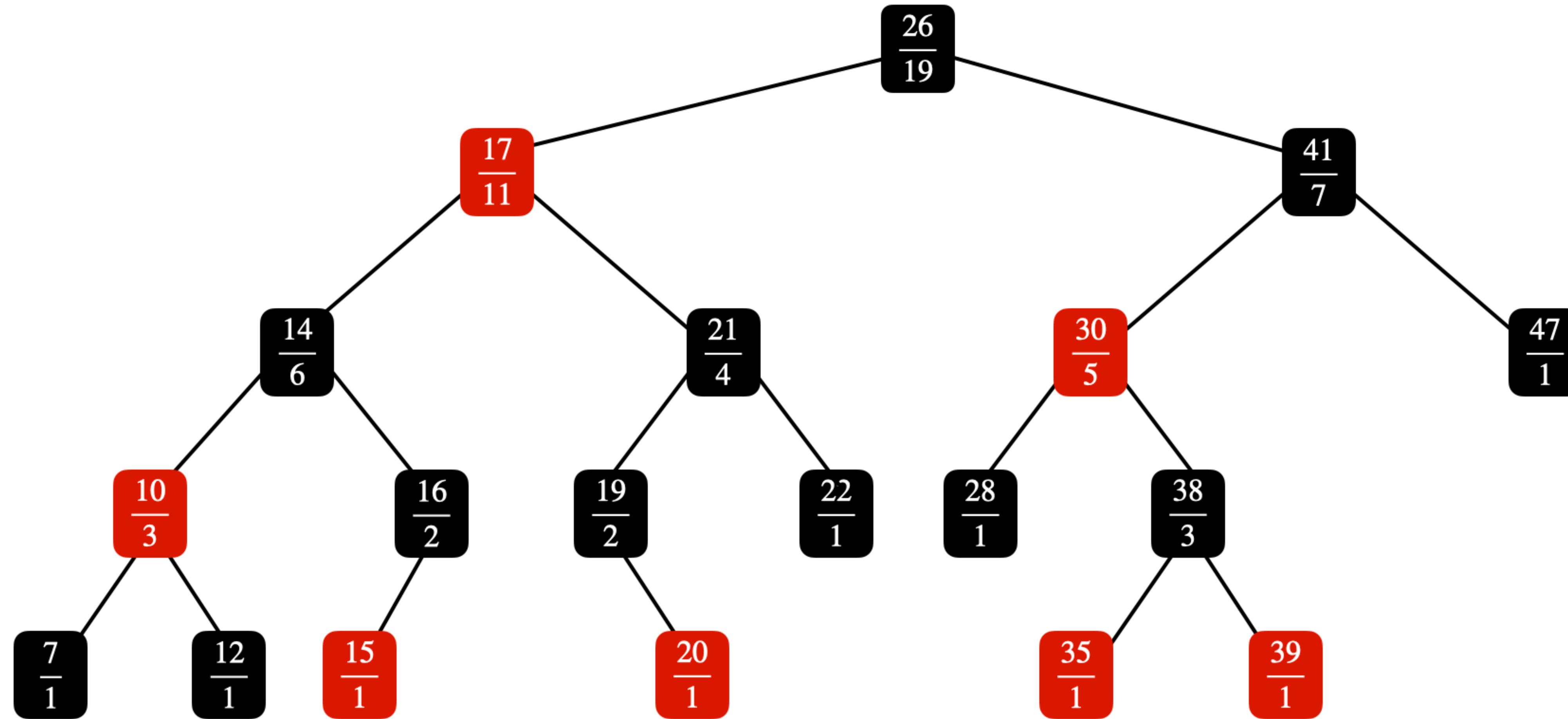
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Example: Find the rank of 38 in the below set or RB-tree.

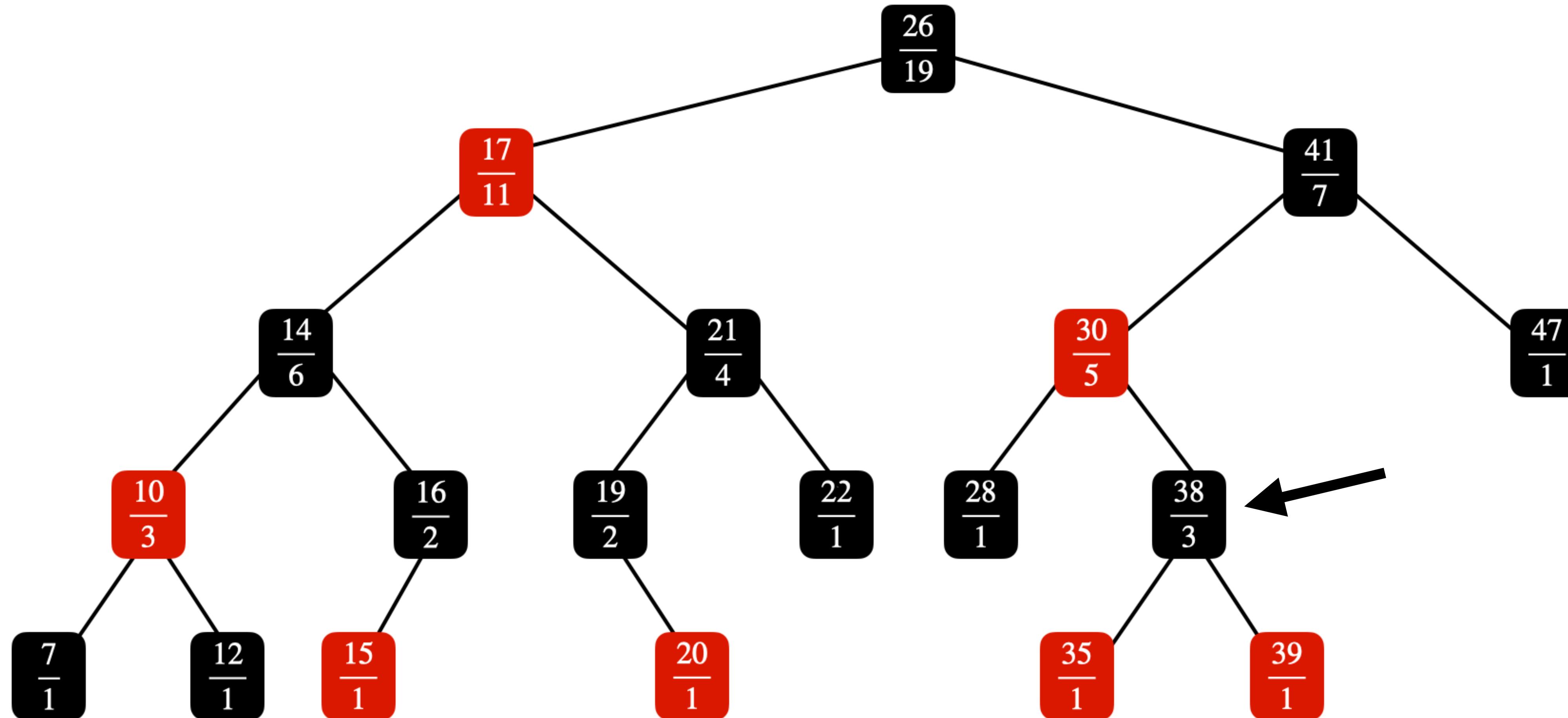
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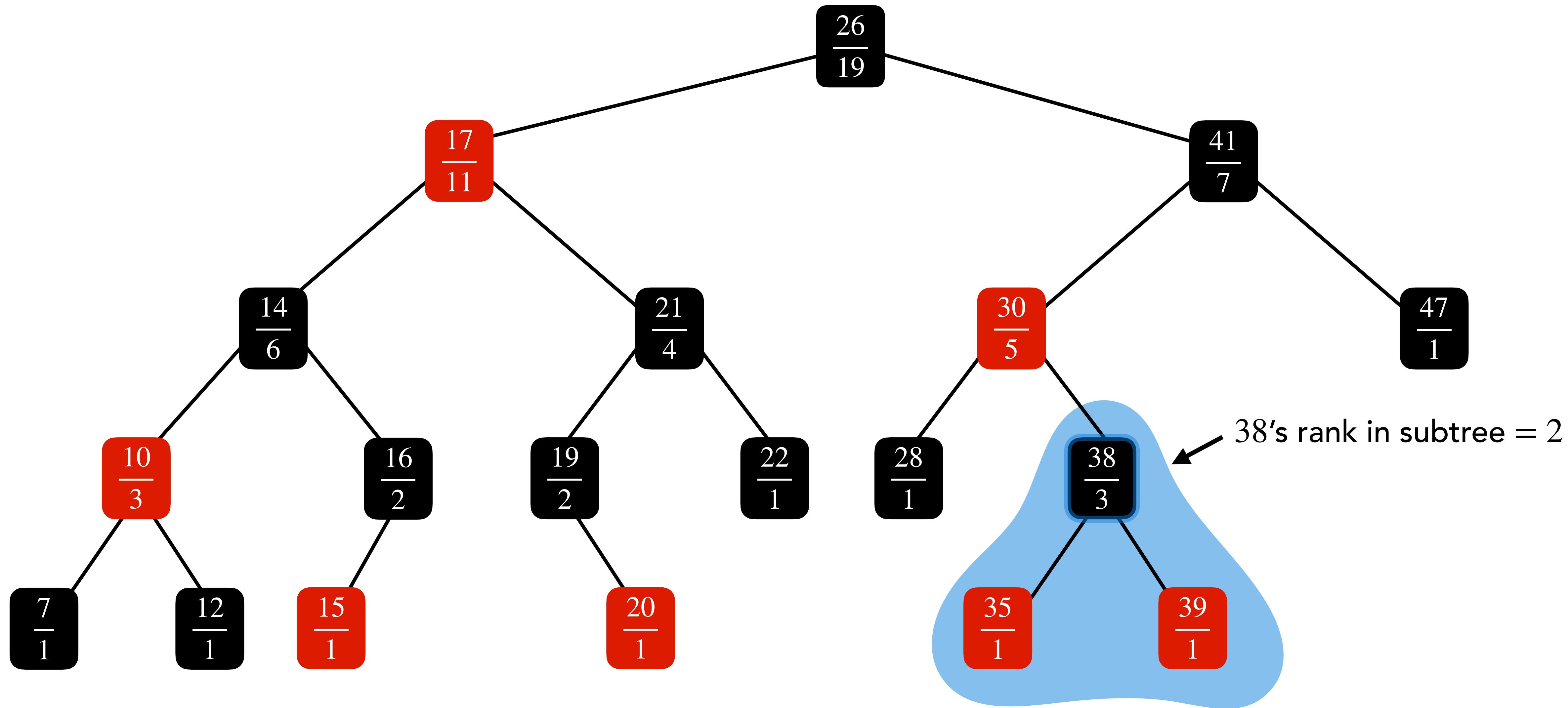


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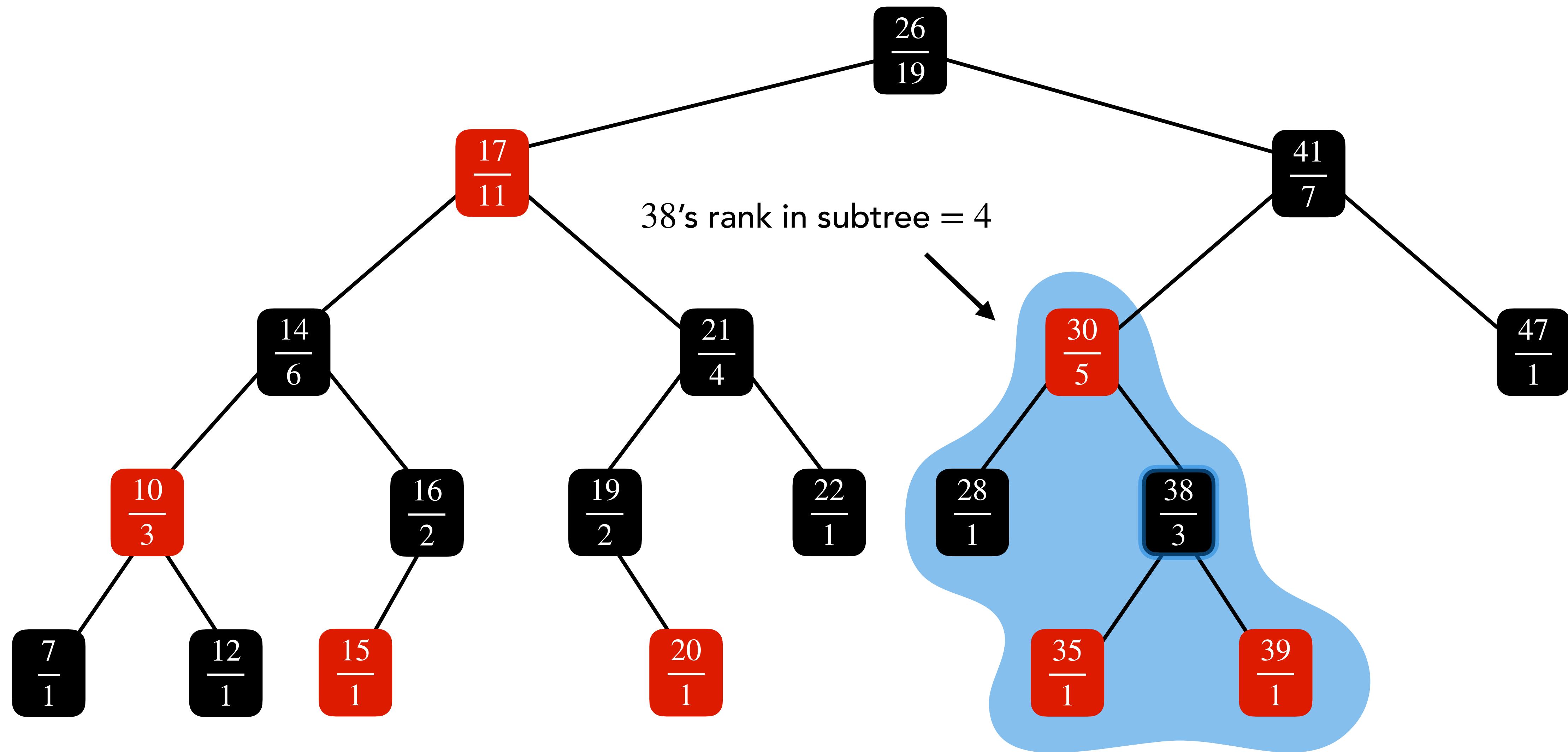
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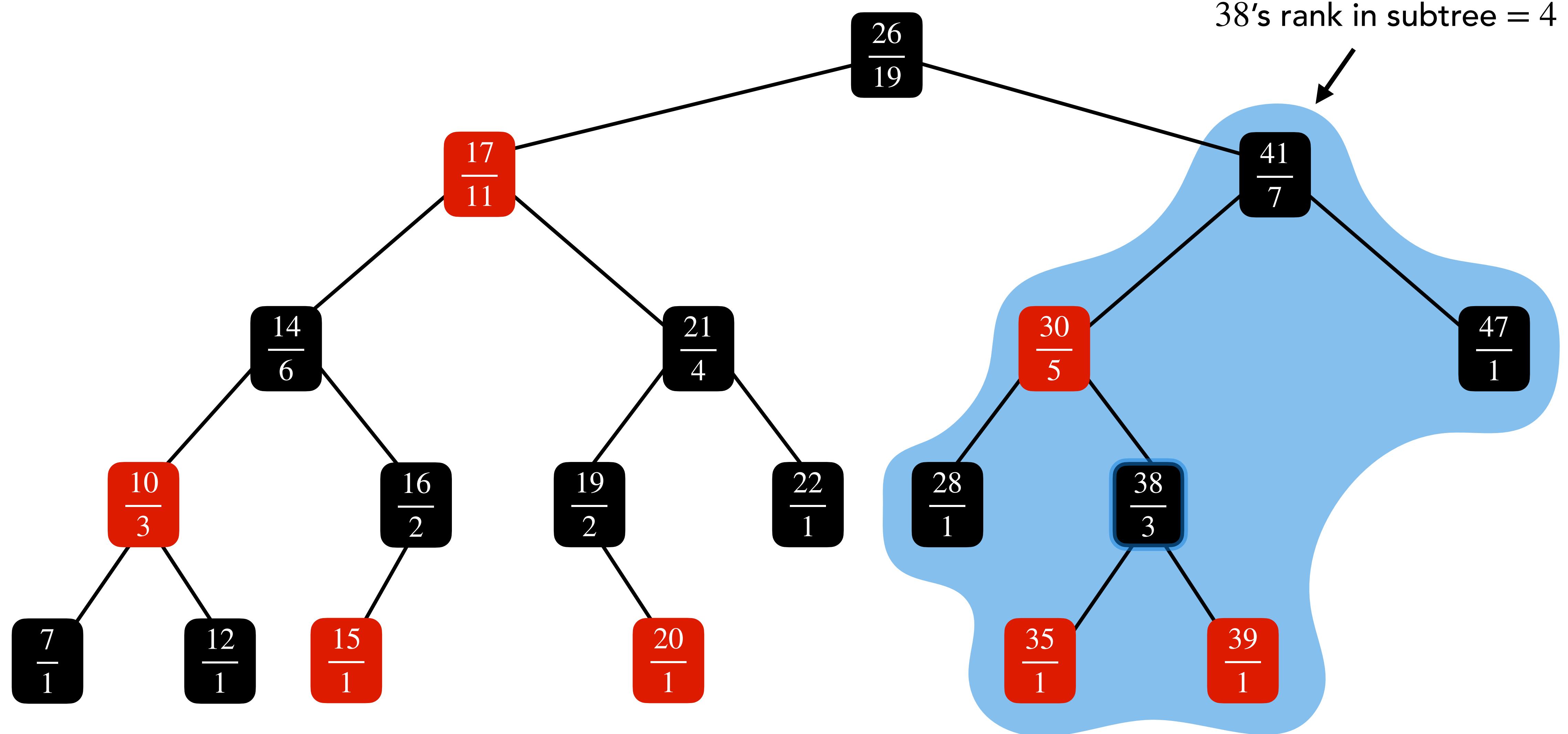
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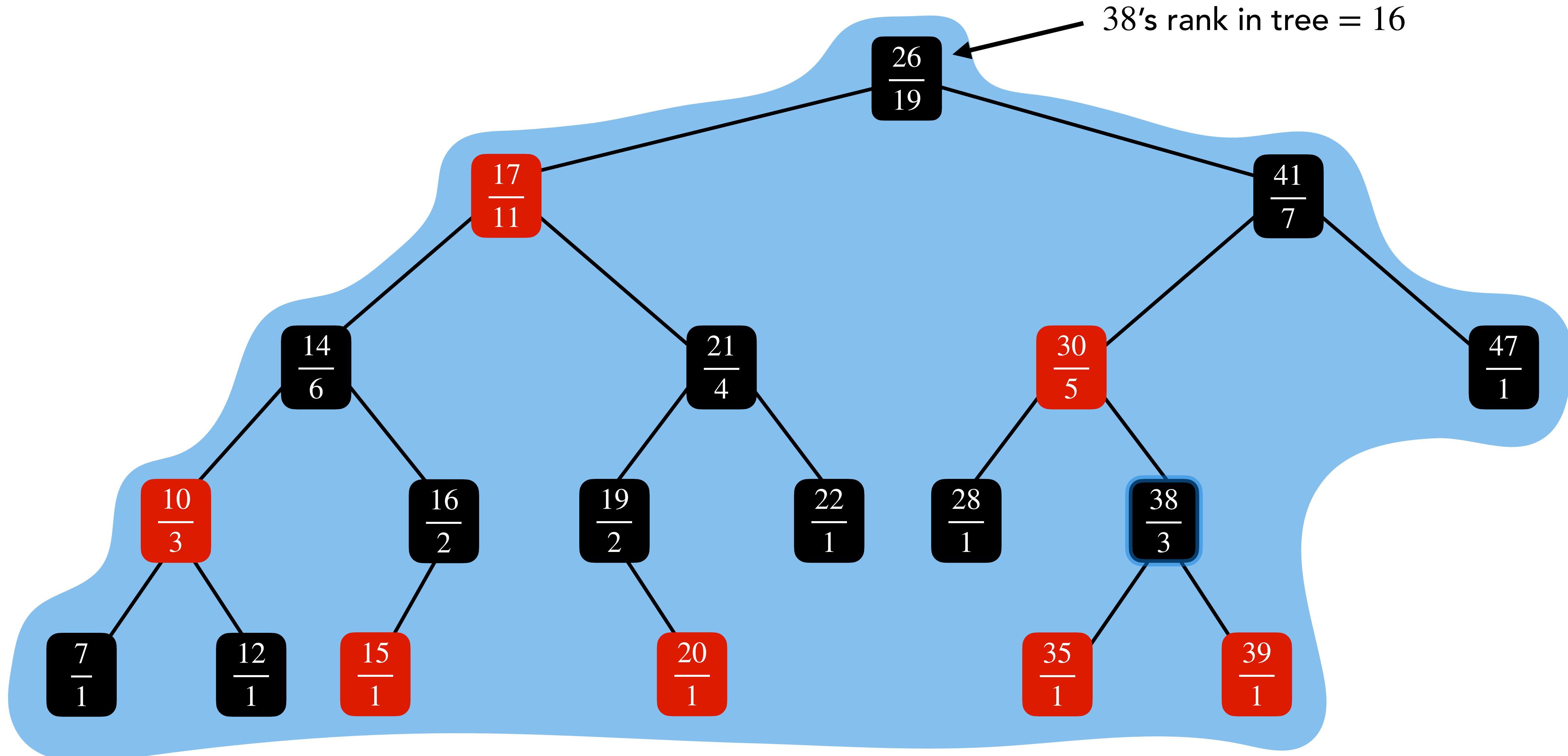
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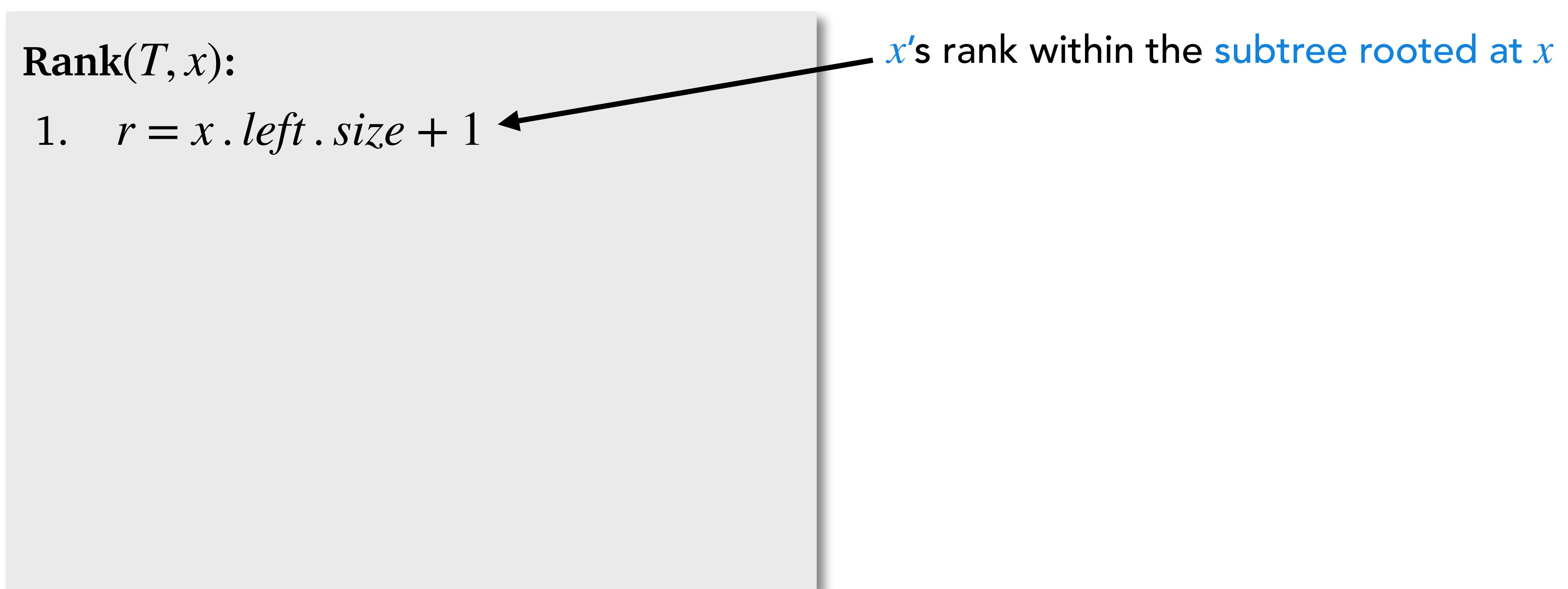
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2. $y = x$

x 's rank within the subtree rooted at x

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3. **while** $y \neq T.\text{root}$

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x' 's rank within the subtree rooted at x

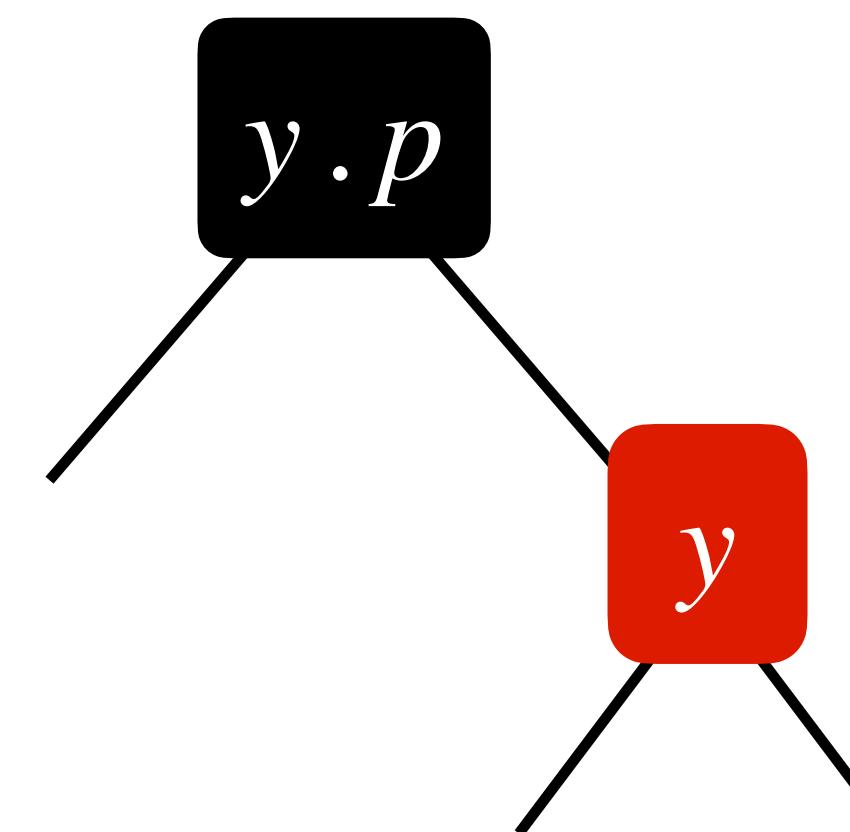
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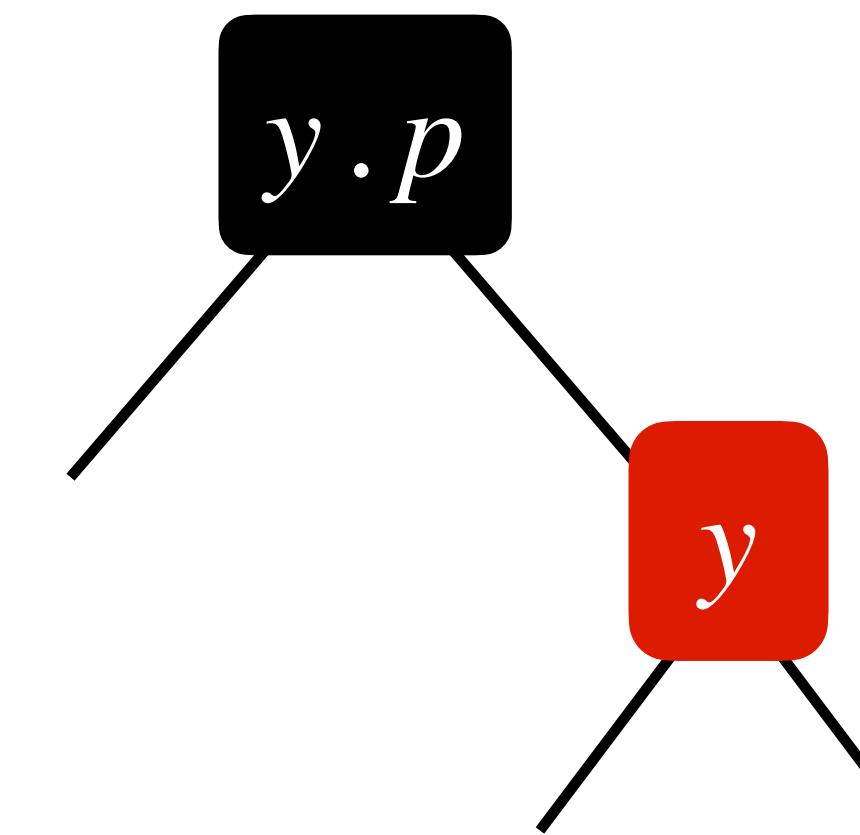
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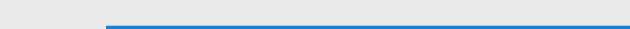


Updating y 's rank when y is the right child of its parent.

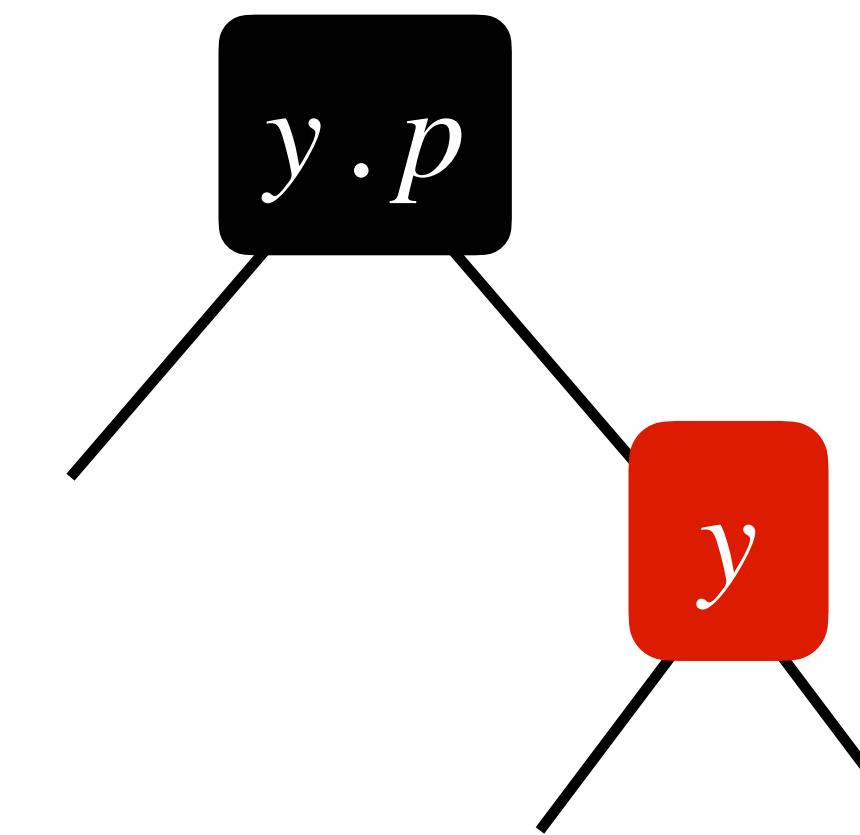
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5. $r =$ 

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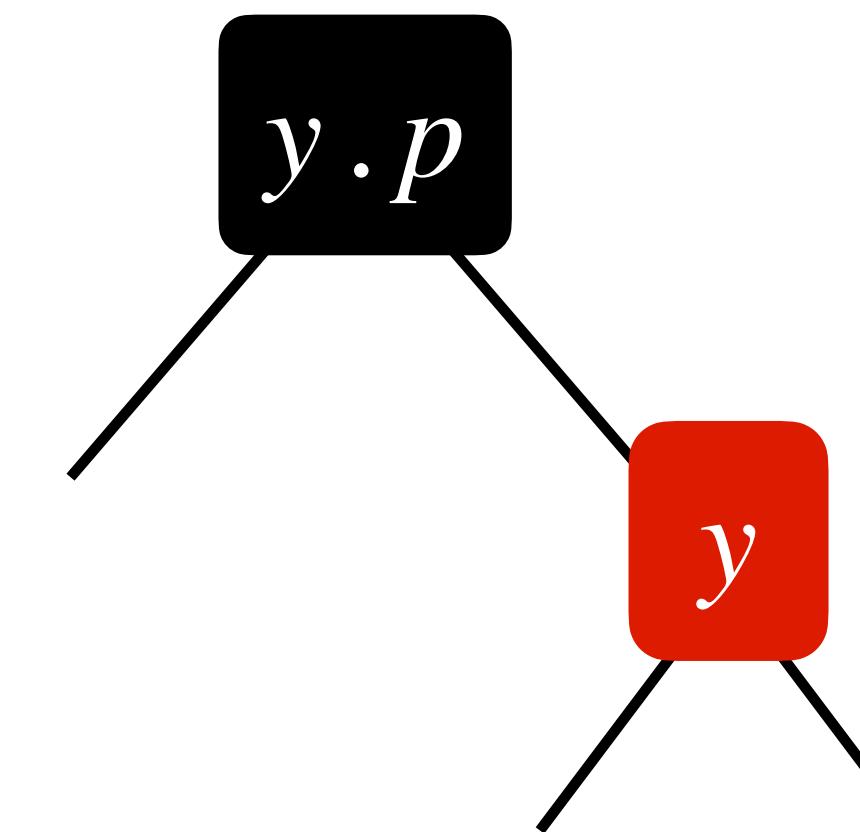
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1. $r = x.left.size + 1$
2. $y = x$
3. **while** $y \neq T.root$
4. **if** $y = y.p.right$
5. $r = r + y.p.left.size + 1$

x 's rank within the subtree rooted at x



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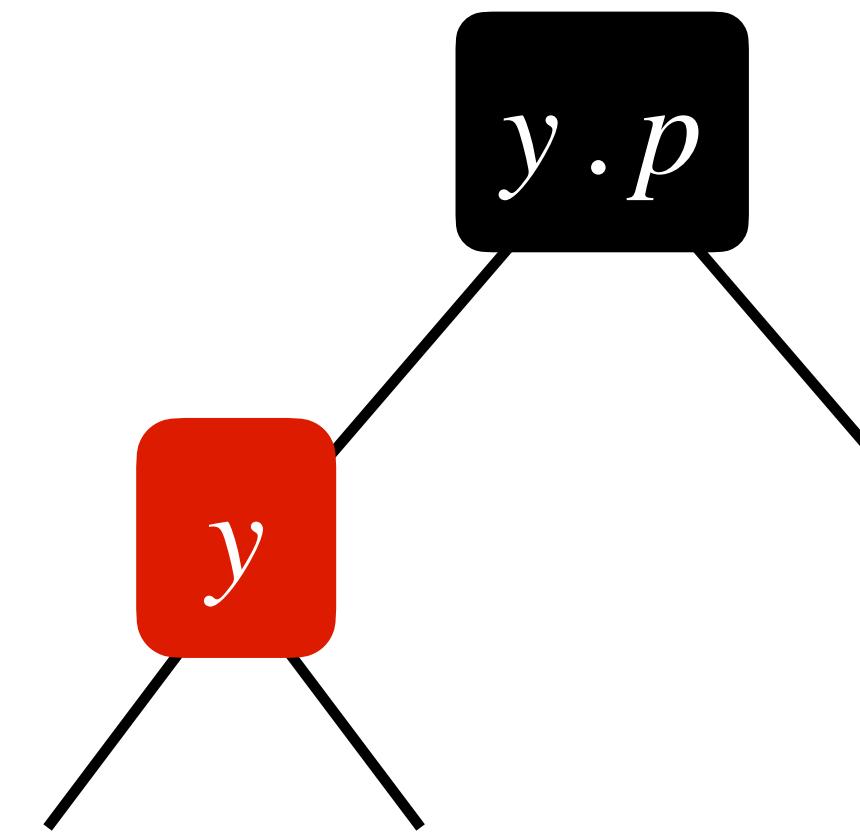
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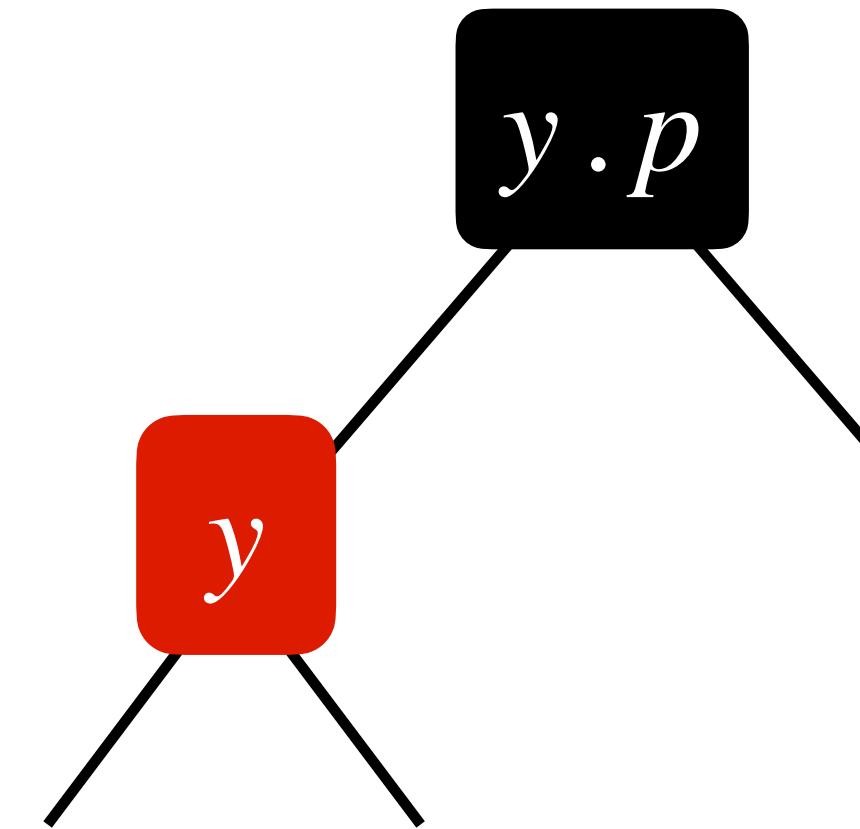
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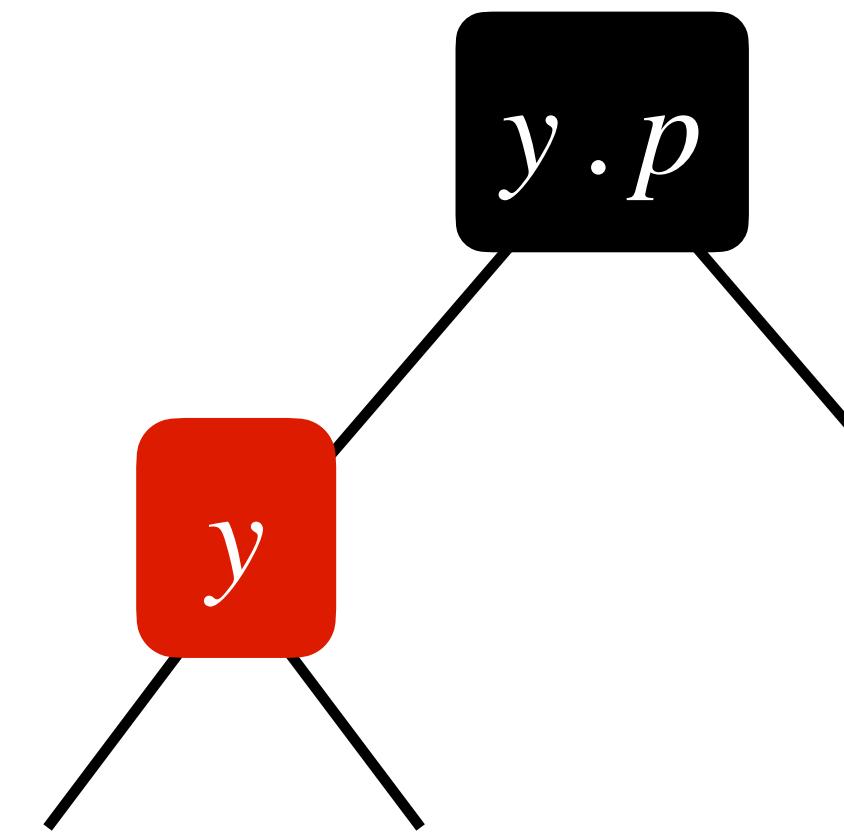
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6. $y = y.p$
7. **return** r

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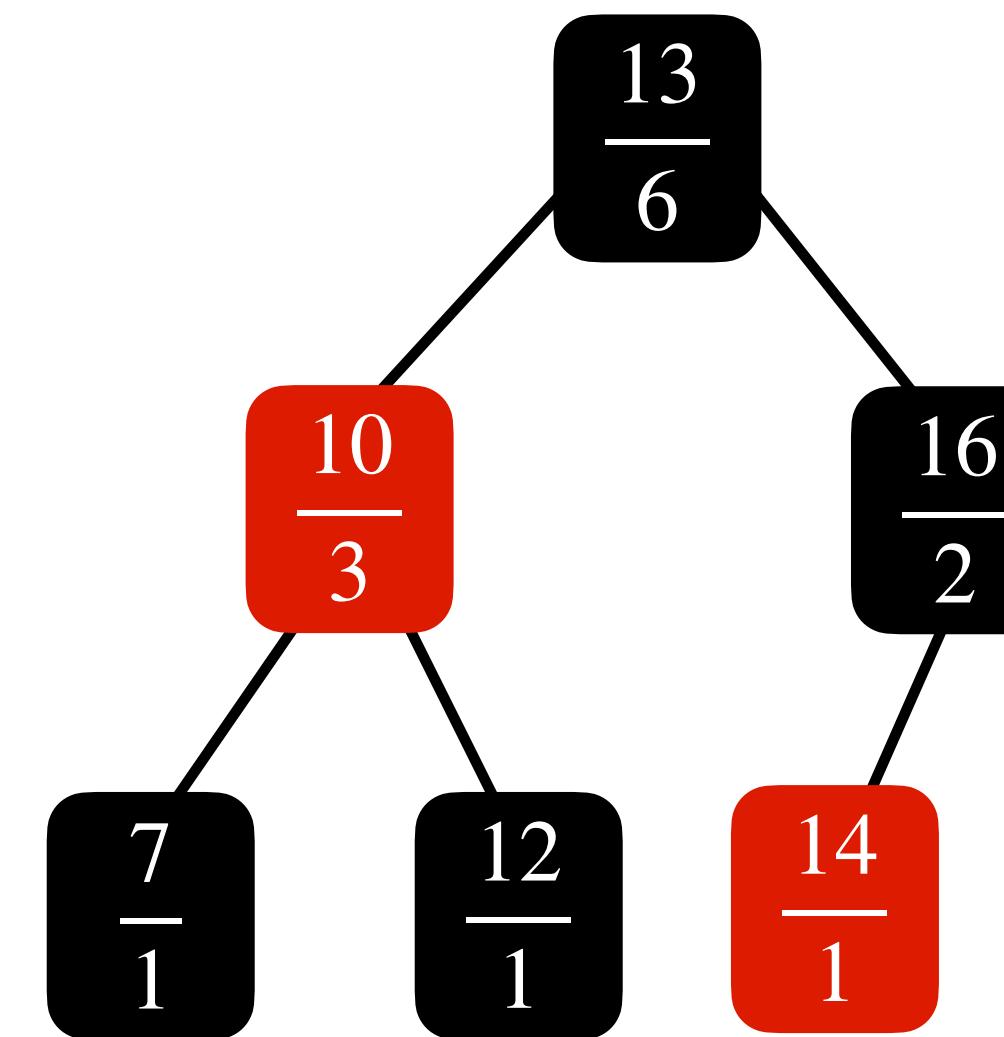
Maintaining Subtree Sizes: Insertion

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Insert 15 in this tree

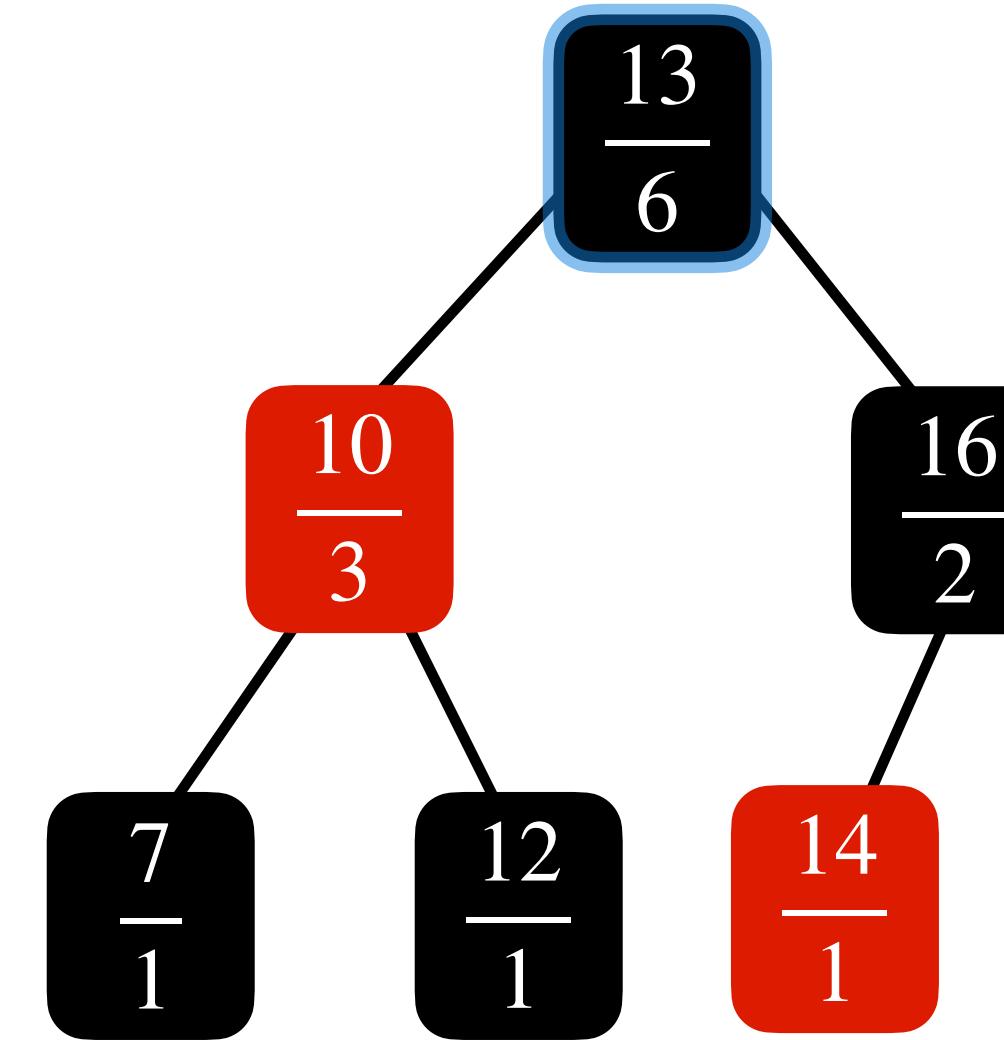
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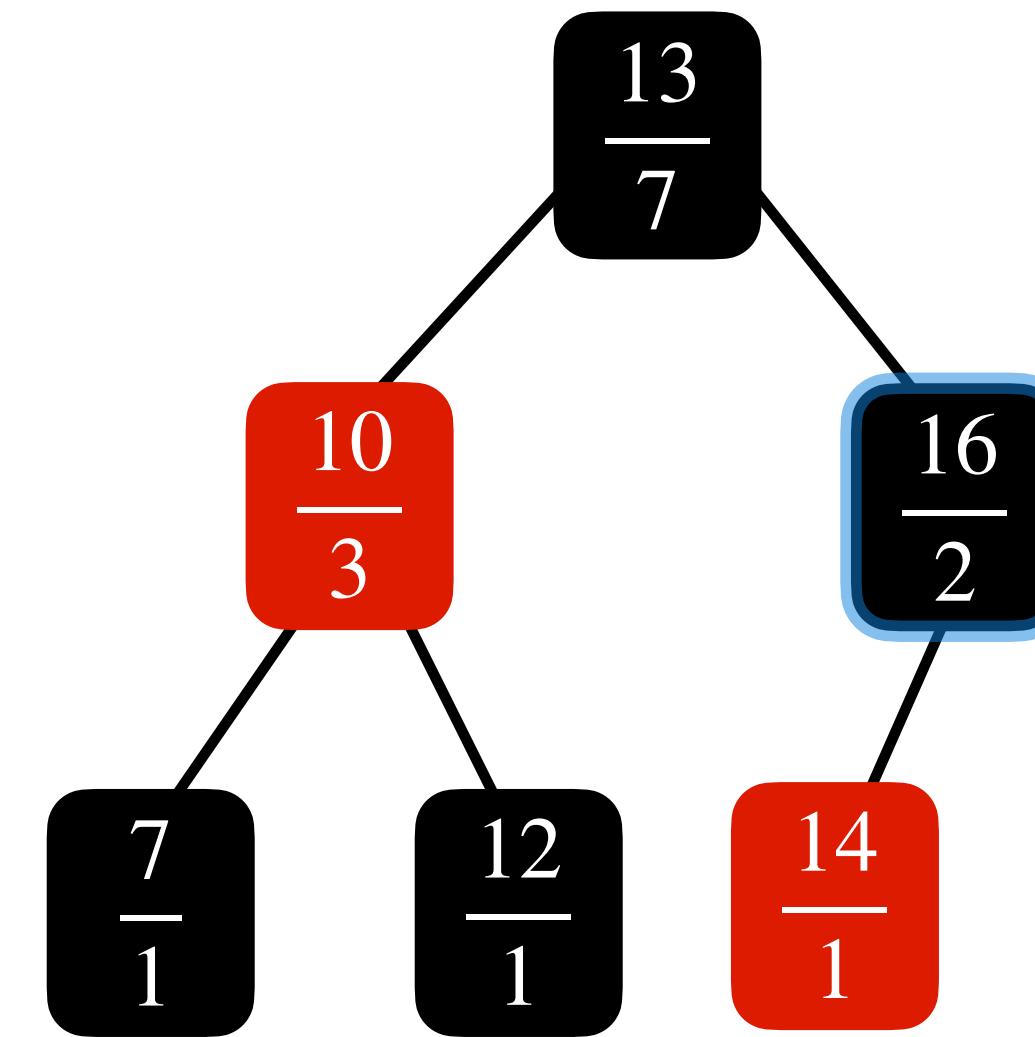
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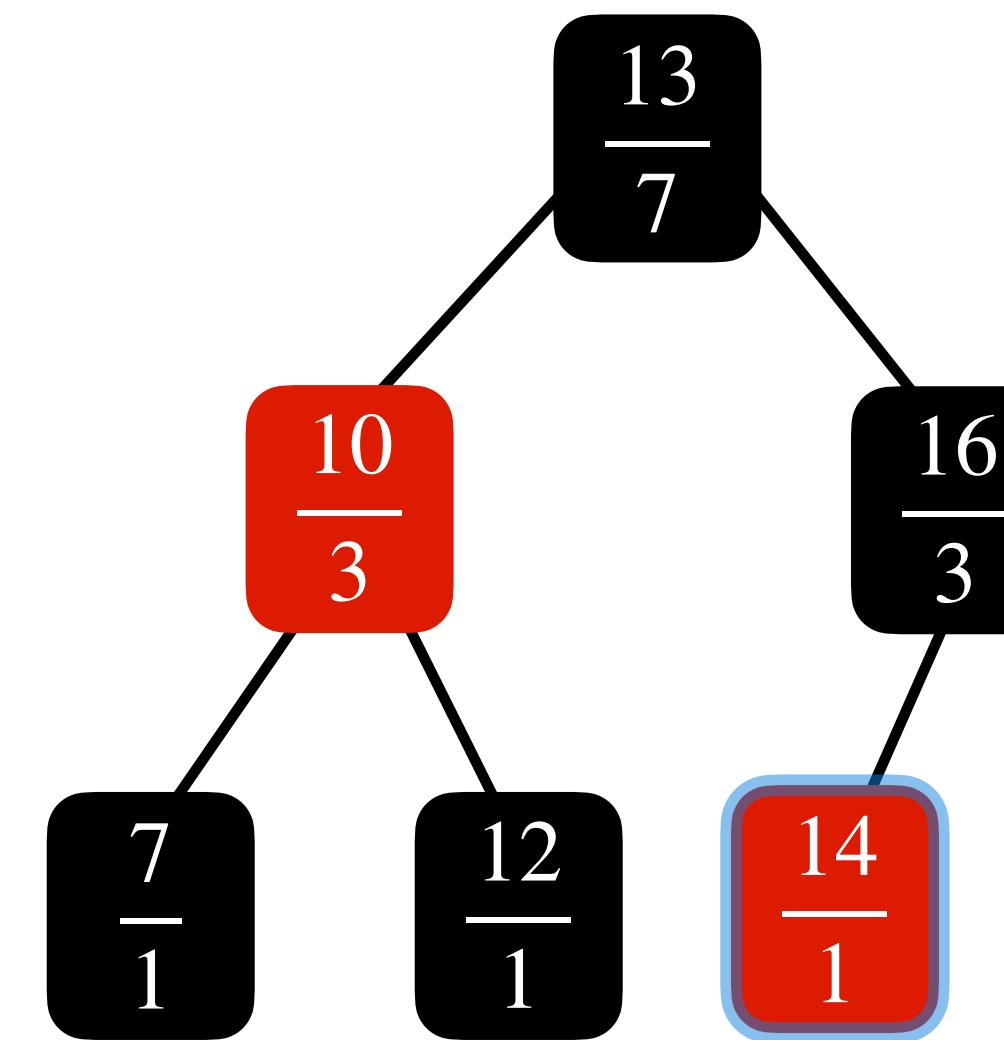
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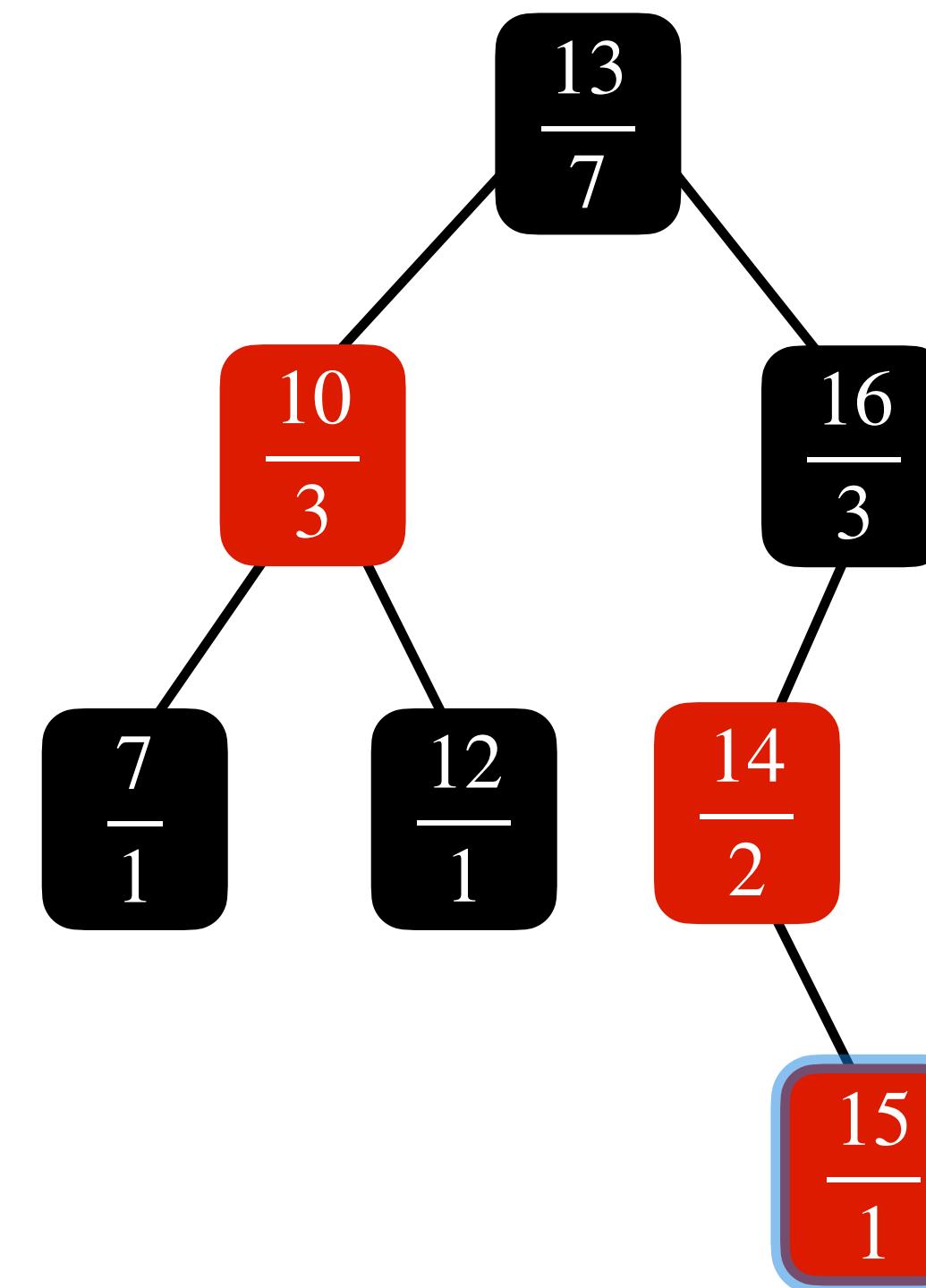
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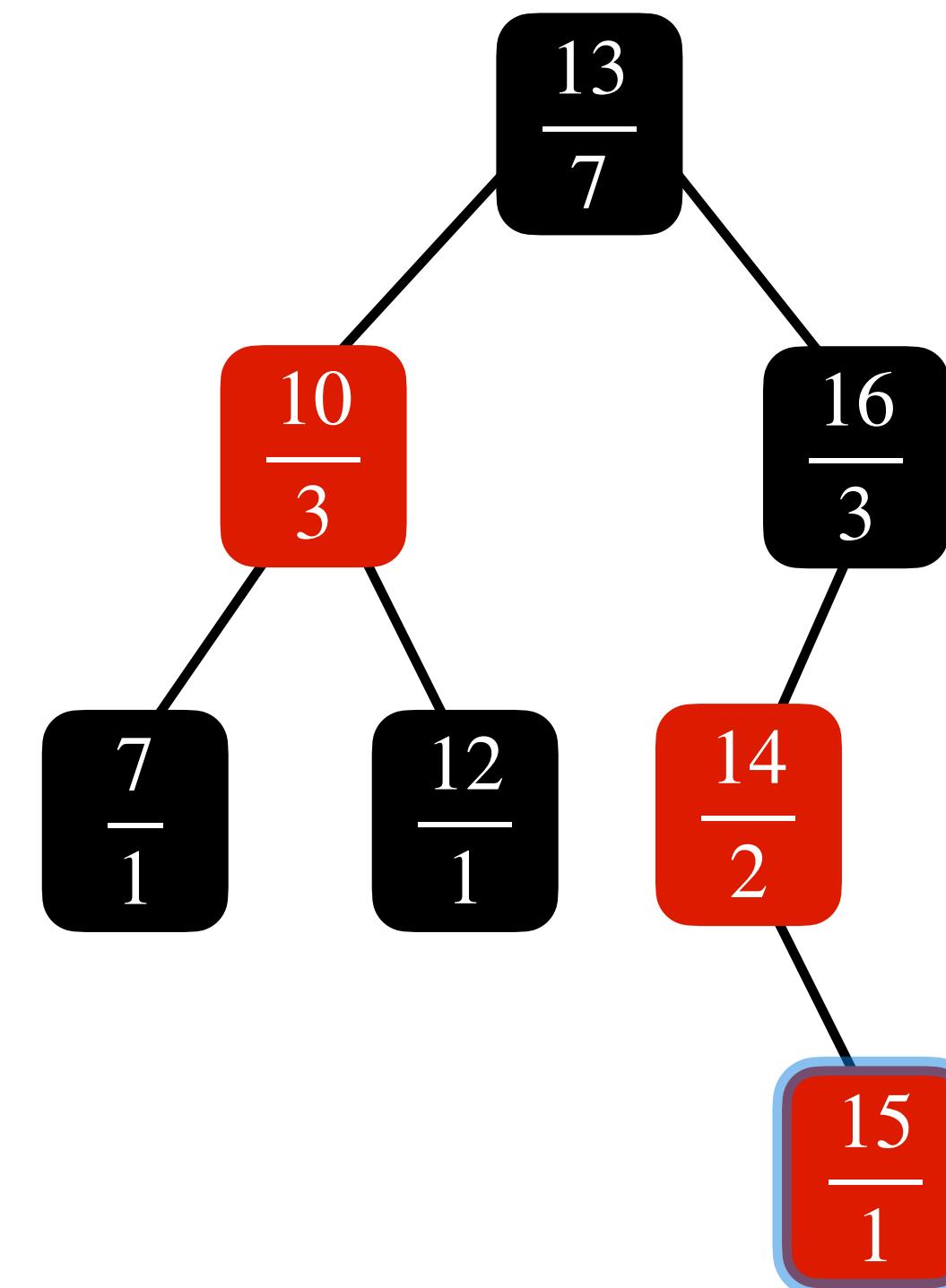
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Maintaining Subtree Sizes: Insertion

Insertion phase:

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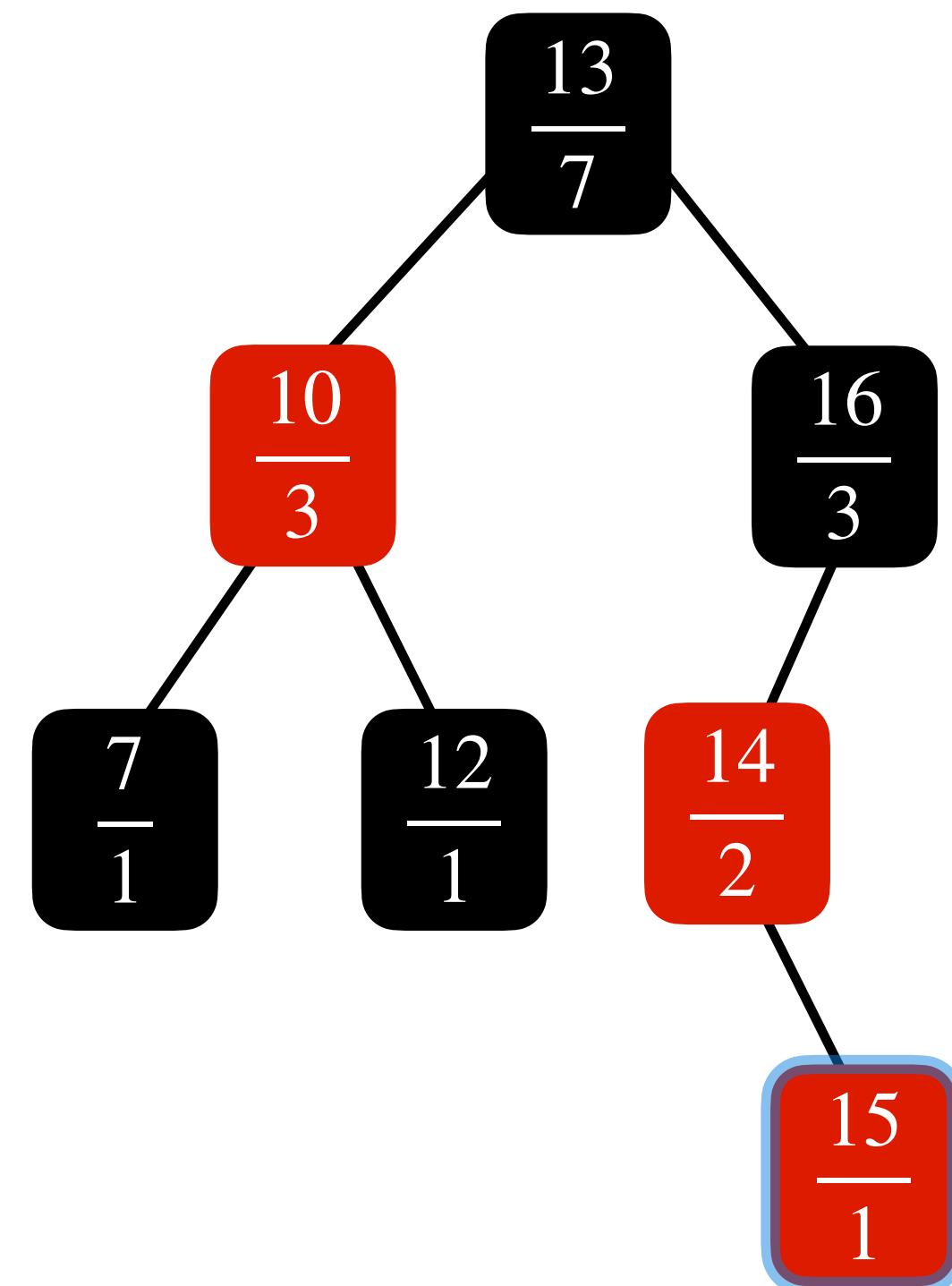


Maintaining Subtree Sizes: Insertion

Insertion phase:

Keep adding 1 to the sizes of every node we visit

Insert 15 in this tree

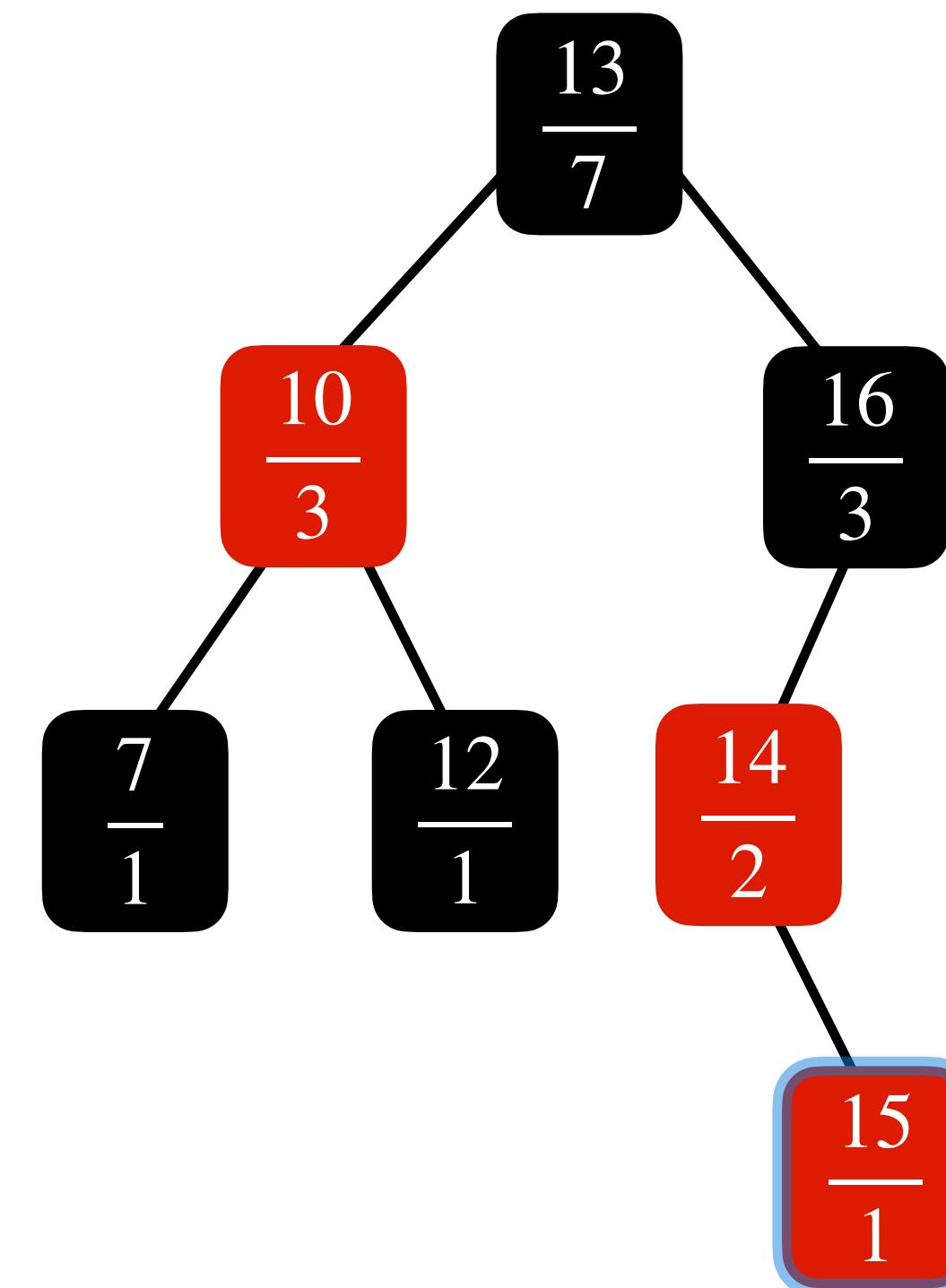


Maintaining Subtree Sizes: Insertion

Insertion phase:

Keep adding 1 to the sizes of every node we visit while searching for the correct leaf

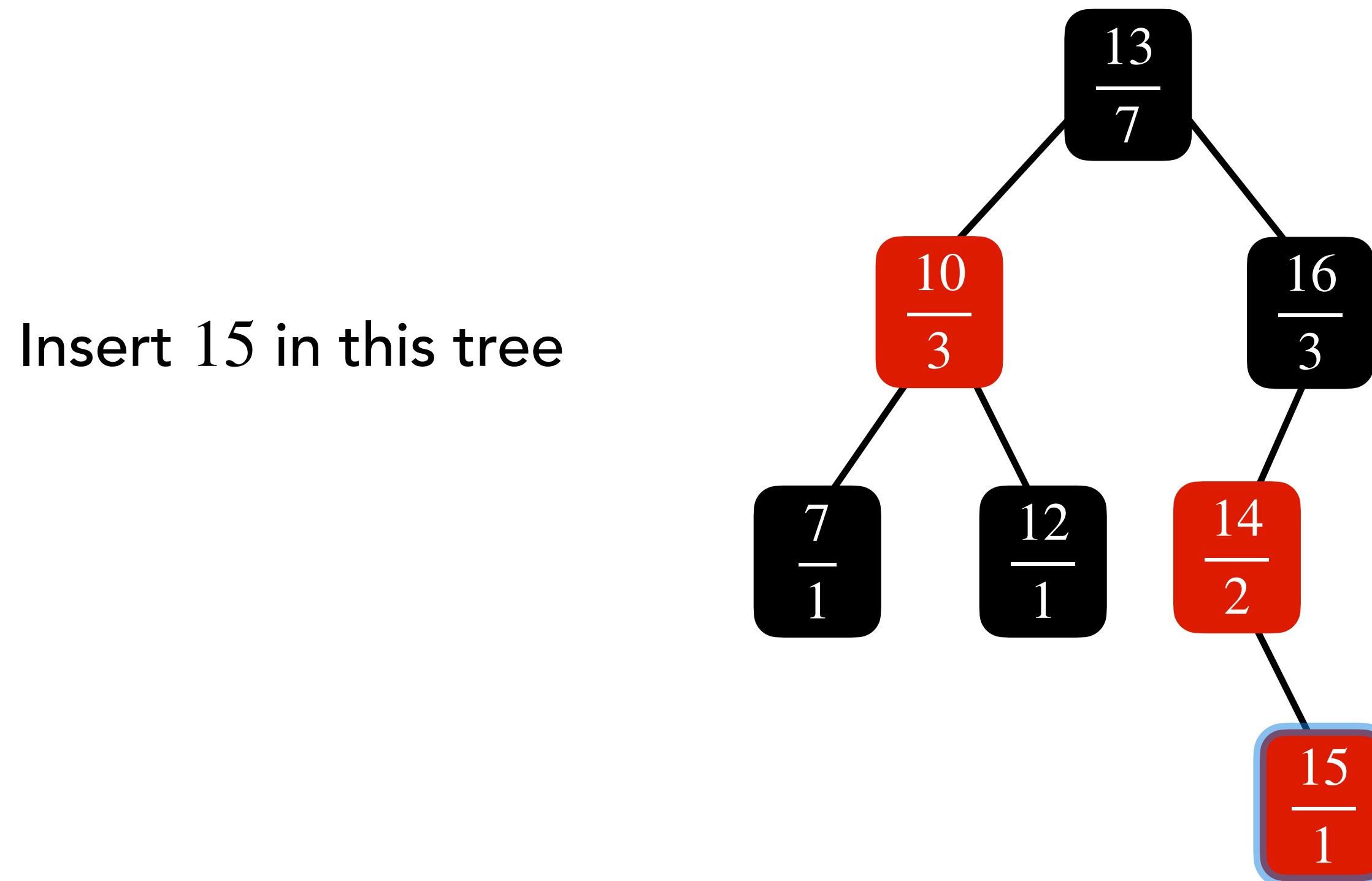
Insert 15 in this tree



Maintaining Subtree Sizes: Insertion

Insertion phase:

Keep adding 1 to the sizes of every node we visit while searching for the correct leaf where new node can be inserted.



Maintaining Subtree Sizes: Insertion

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Fix-up phase:

Maintaining Subtree Sizes: Insertion

Fix-up phase:

- Fix-ups involve only rotations and recolouring.

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Fix-up phase:

- Fix-ups involve only **rotations** and **recolouring**.
- Recolouring doesn't require changing sizes.

Maintaining Subtree Sizes: Insertion

Fix-up phase:

- Fix-ups involve only **rotations** and **recolouring**.
- Recolouring doesn't require changing sizes.
- During **rotations** size changes are doable in constant time.

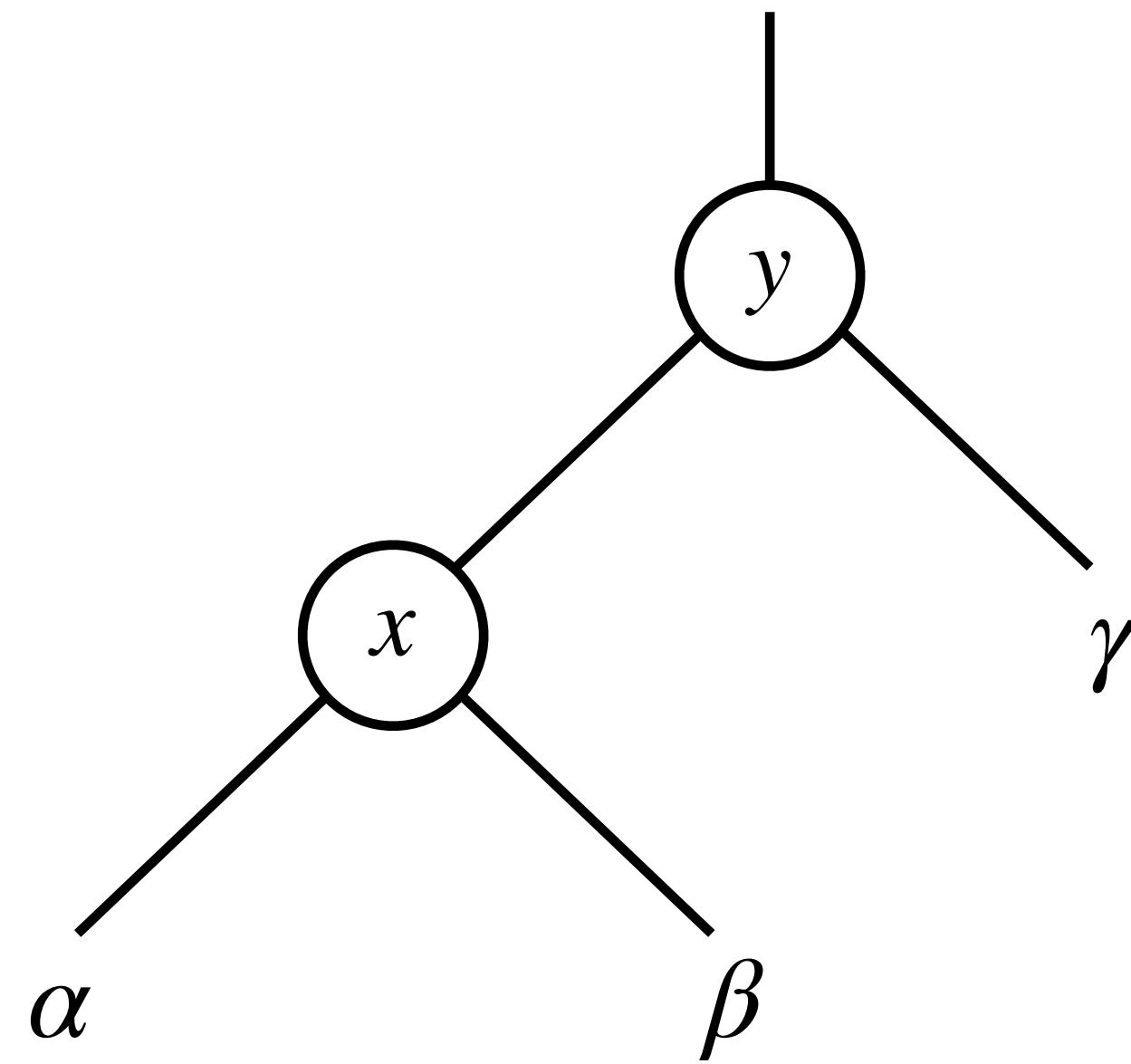
Maintaining Subtree Sizes: Rotations

Maintaining Subtree Sizes: Rotations

Maintaining subtree sizes during rotations is easy.

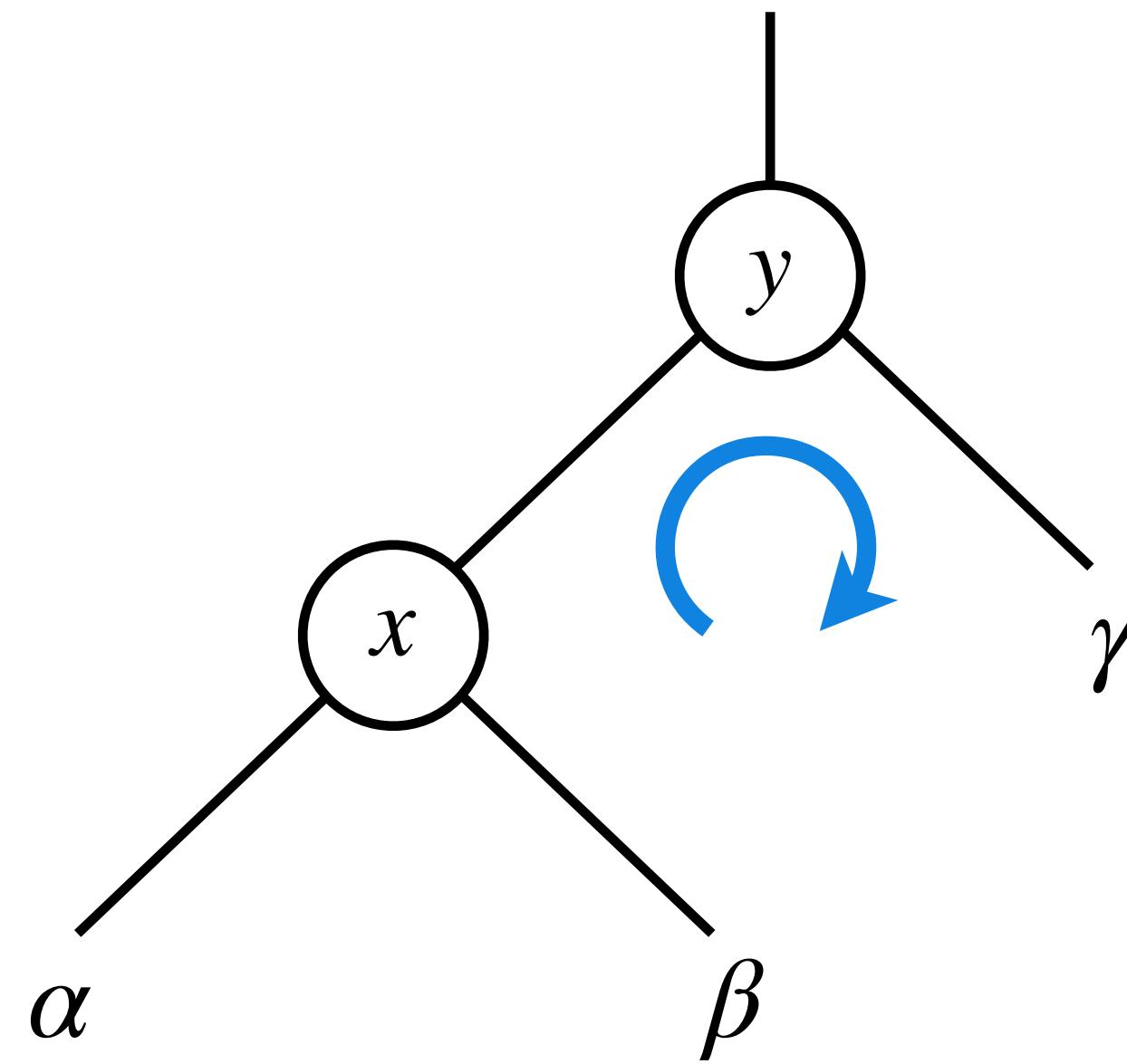
Maintaining Subtree Sizes: Rotations

Maintaining subtree sizes during rotations is easy.



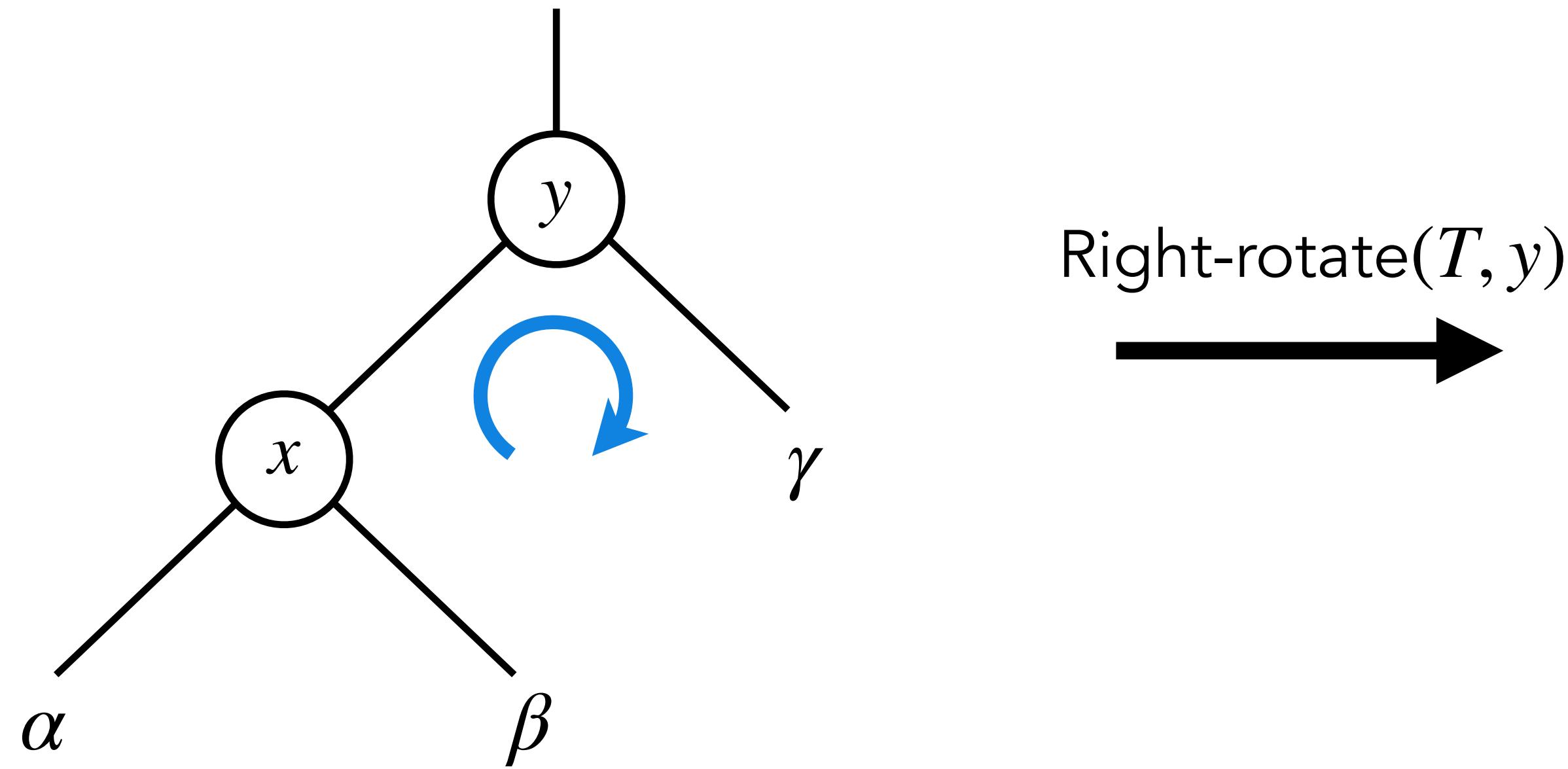
Maintaining Subtree Sizes: Rotations

Maintaining subtree sizes during rotations is easy.



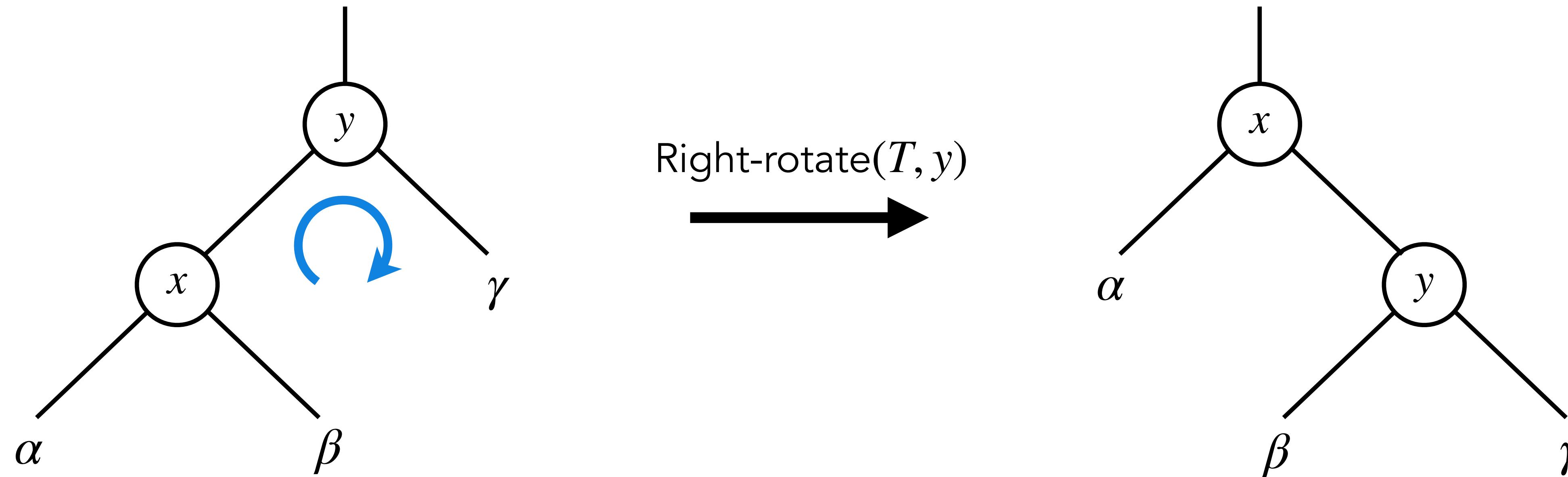
Maintaining Subtree Sizes: Rotations

Maintaining subtree sizes during rotations is easy.



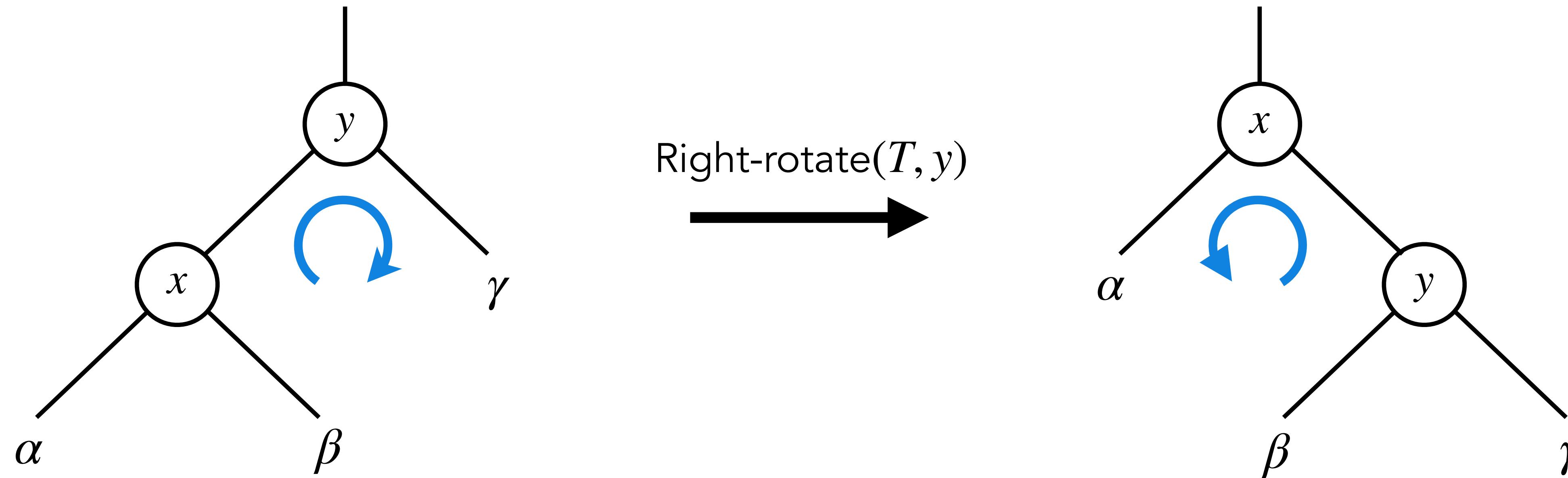
Maintaining Subtree Sizes: Rotations

Maintaining subtree sizes during rotations is easy.



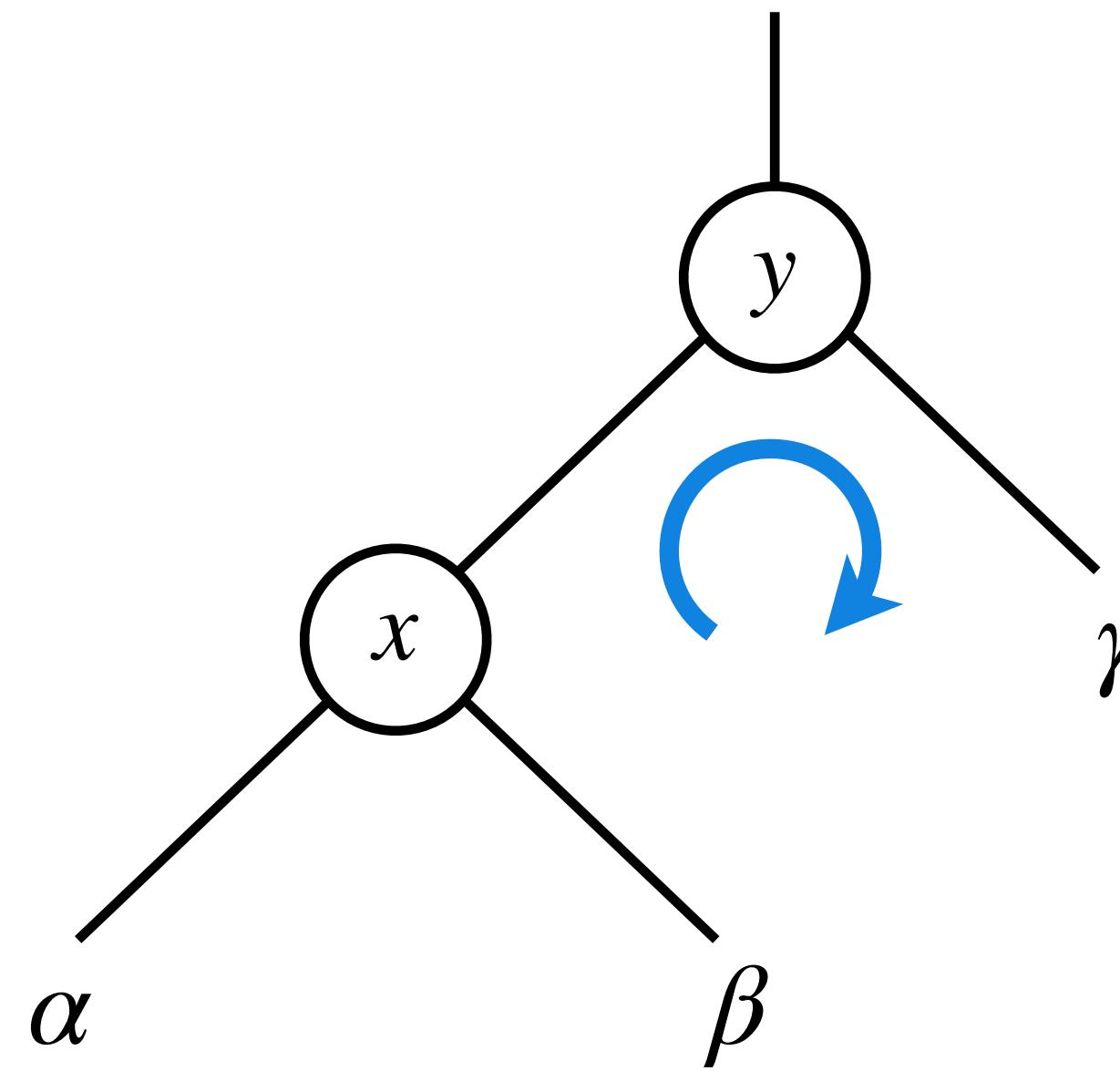
Maintaining Subtree Sizes: Rotations

Maintaining subtree sizes during rotations is easy.

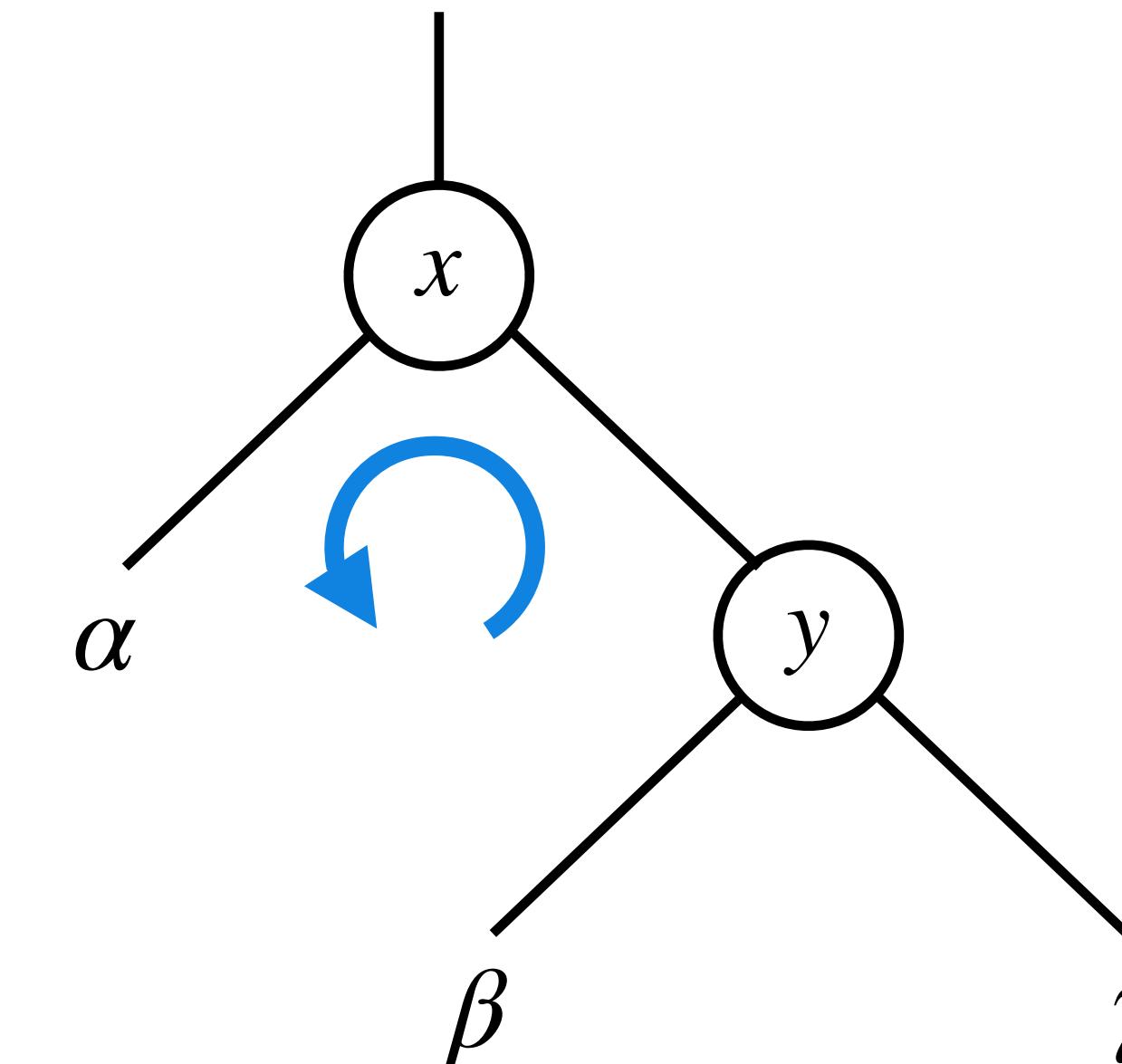
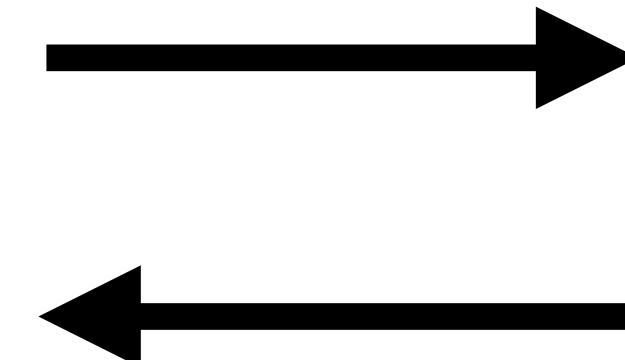


Maintaining Subtree Sizes: Rotations

Maintaining subtree sizes during rotations is easy.

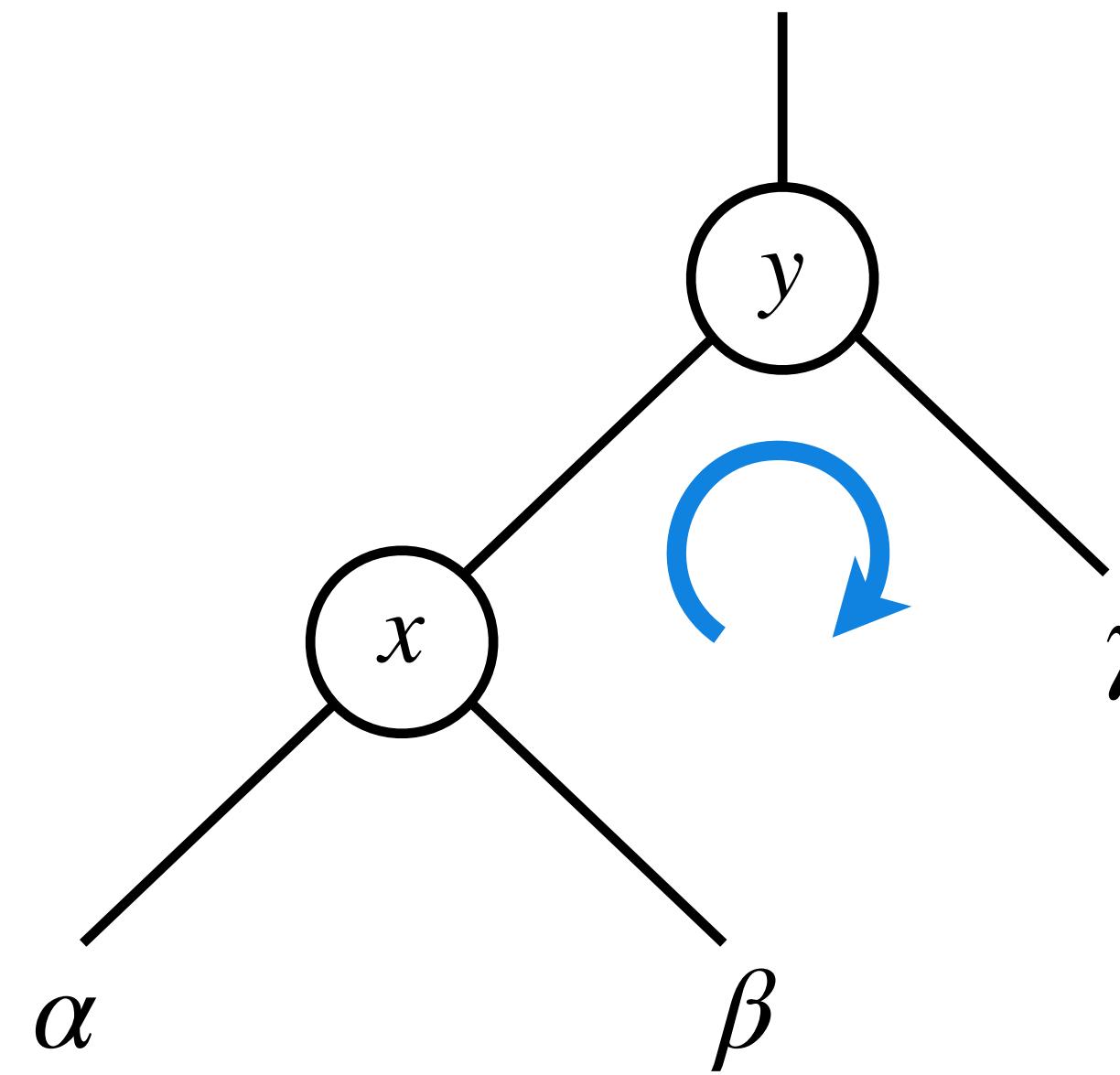


Right-rotate(T, y)

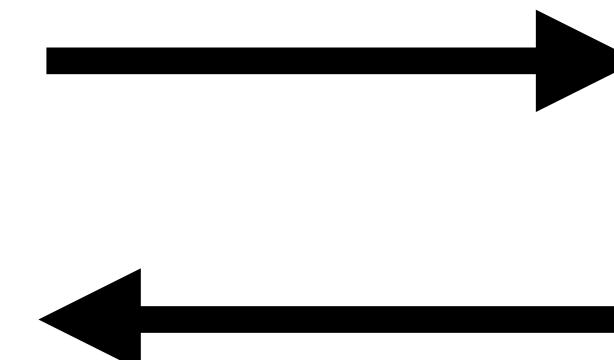


Maintaining Subtree Sizes: Rotations

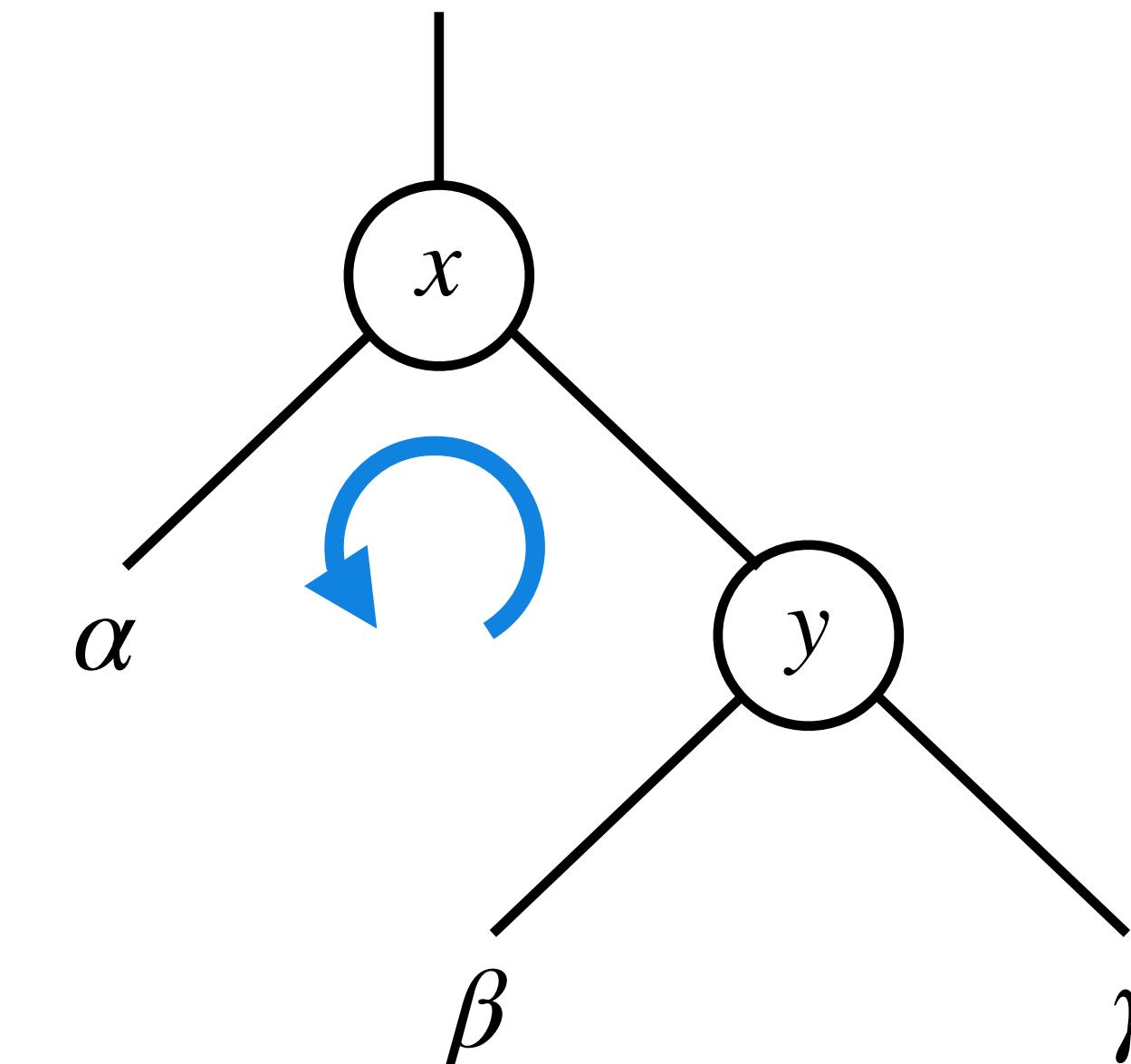
Maintaining subtree sizes during rotations is easy.



Right-rotate(T, y)



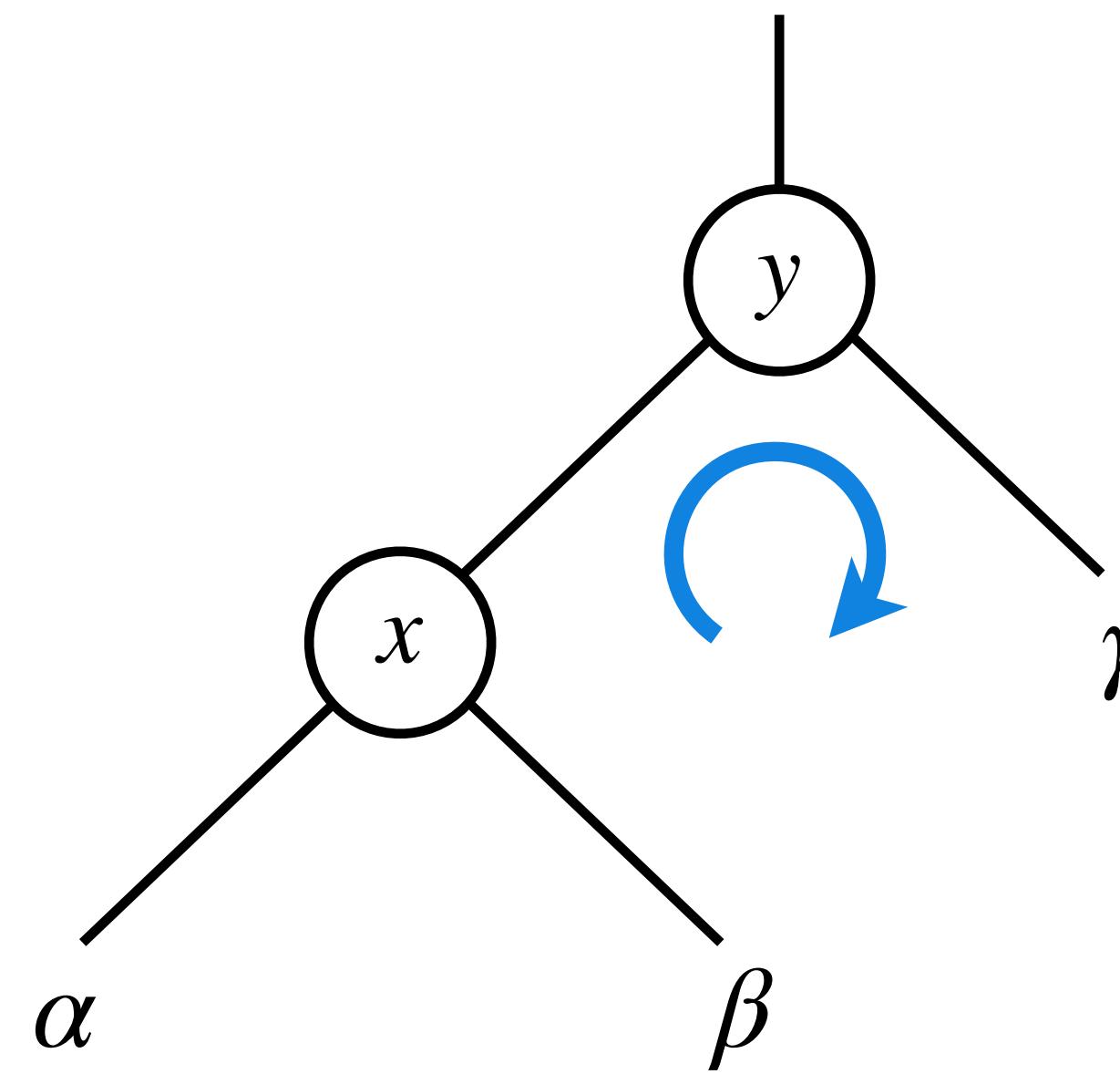
Left-rotate(T, x)



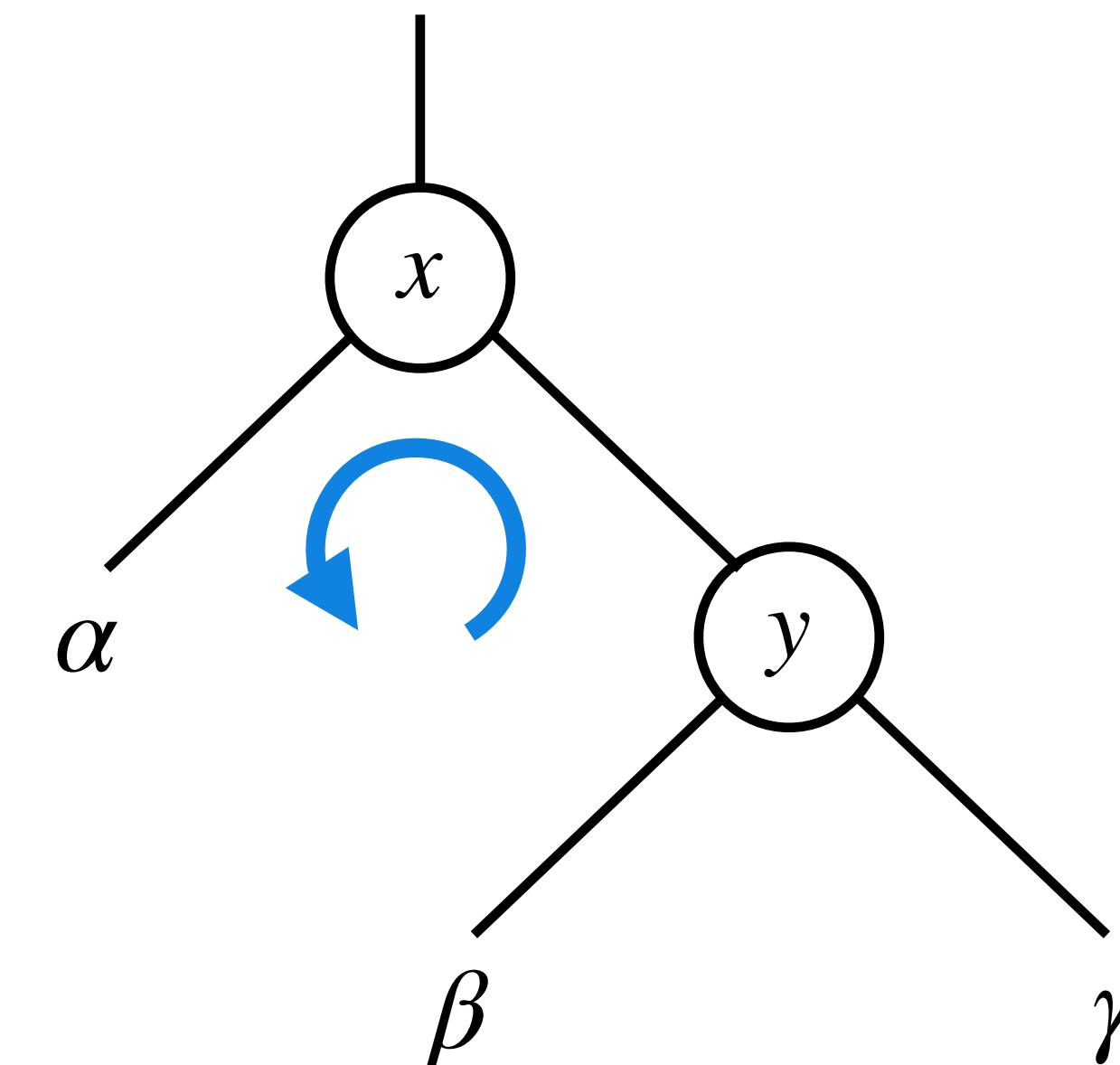
After $\text{Left-rotate}(T, x)$, do the following:

Maintaining Subtree Sizes: Rotations

Maintaining subtree sizes during rotations is easy.



Right-rotate(T, y)

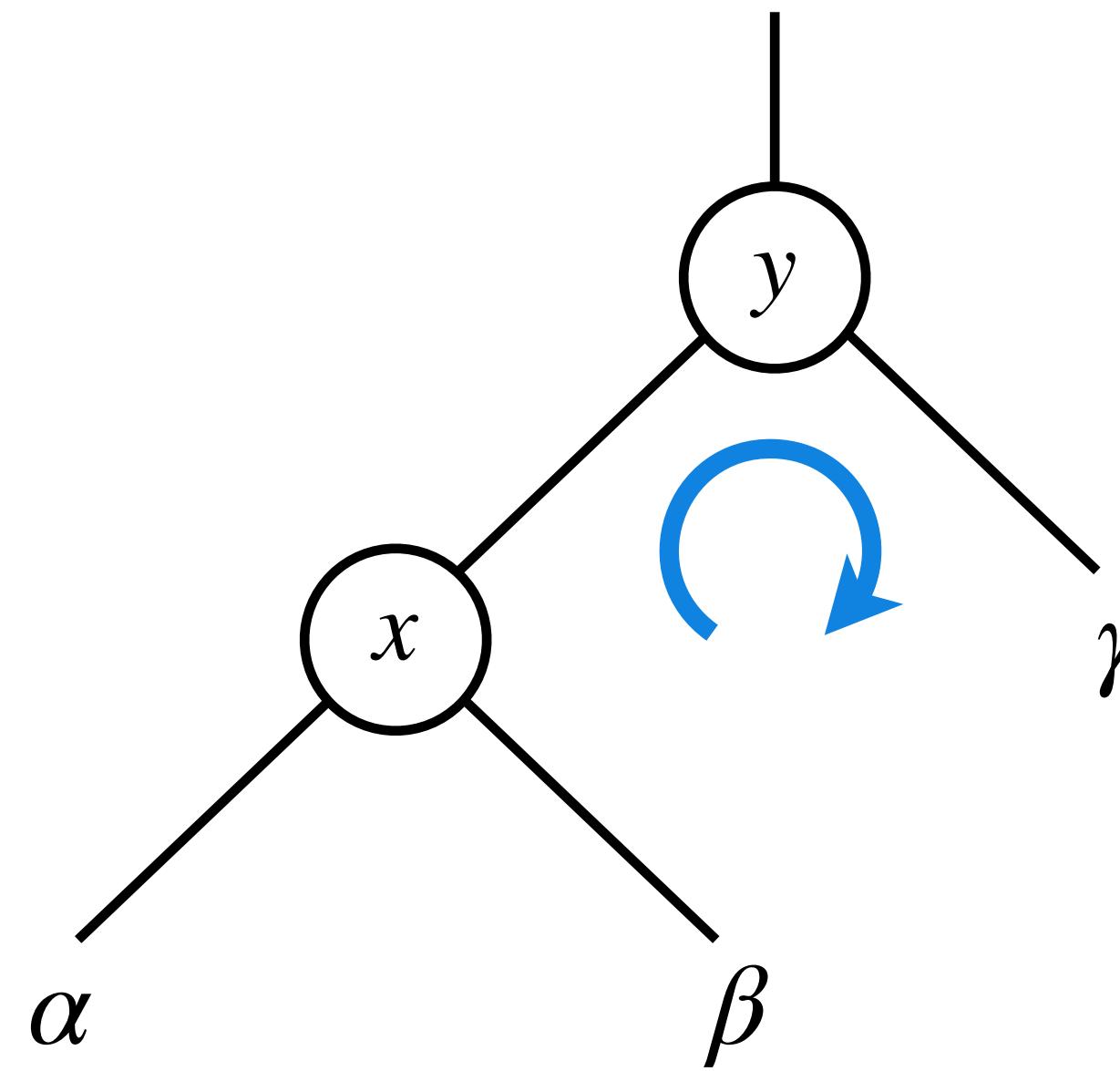


After $\text{Left-rotate}(T, x)$, do the following:

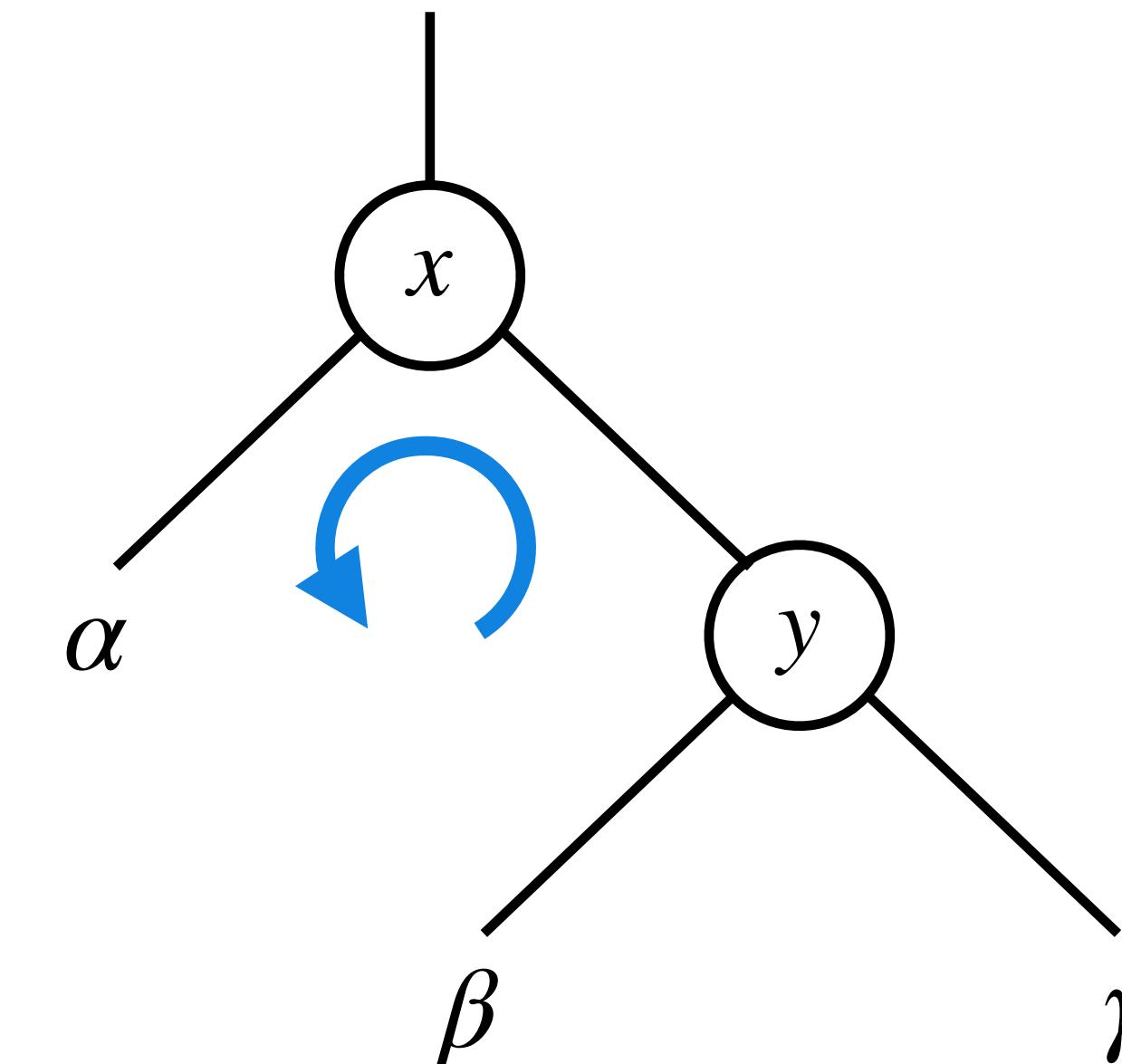
1) $y.size =$

Maintaining Subtree Sizes: Rotations

Maintaining subtree sizes during rotations is easy.



Right-rotate(T, y)
Left-rotate(T, x)

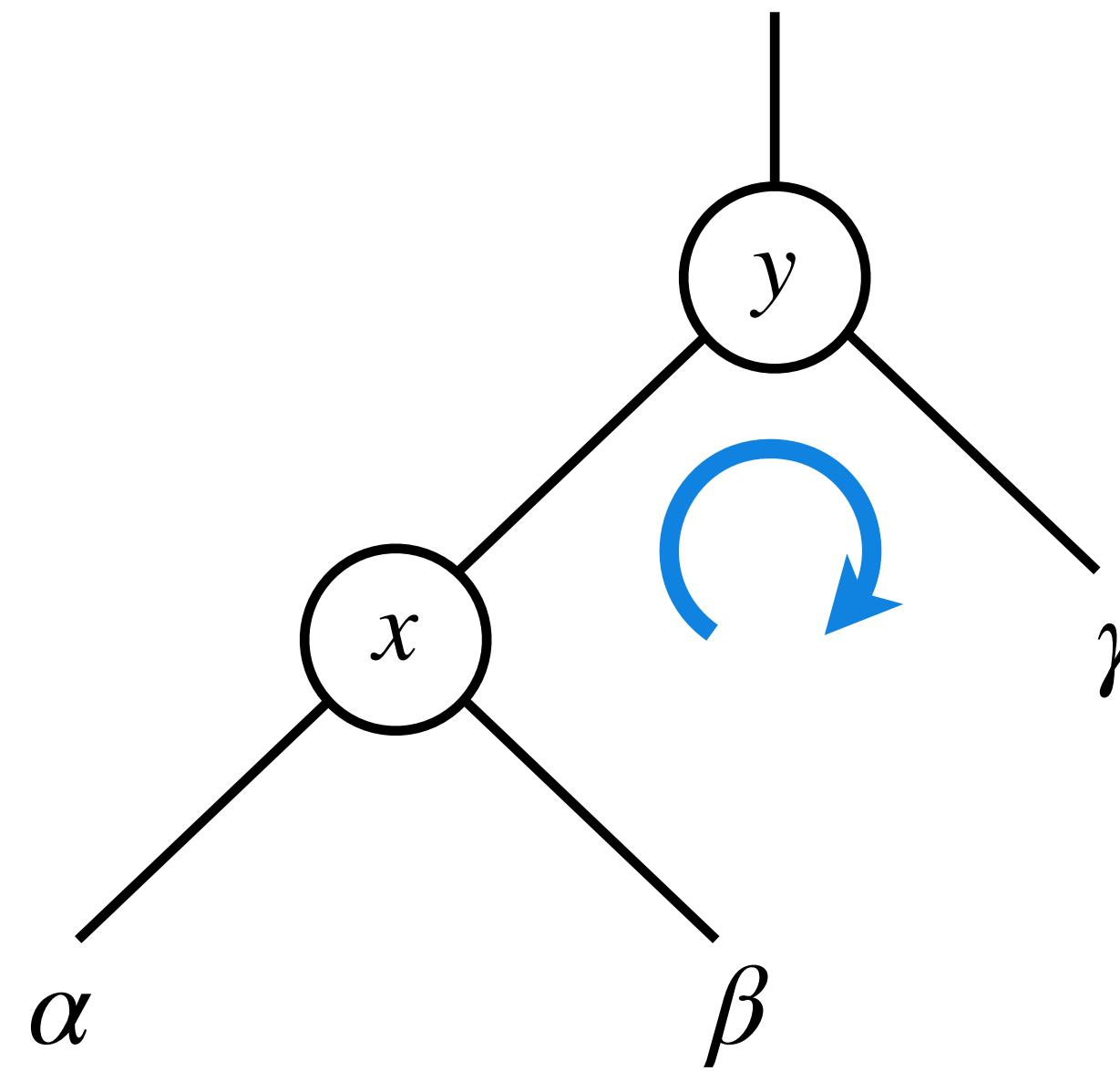


After $\text{Left-rotate}(T, x)$, do the following:

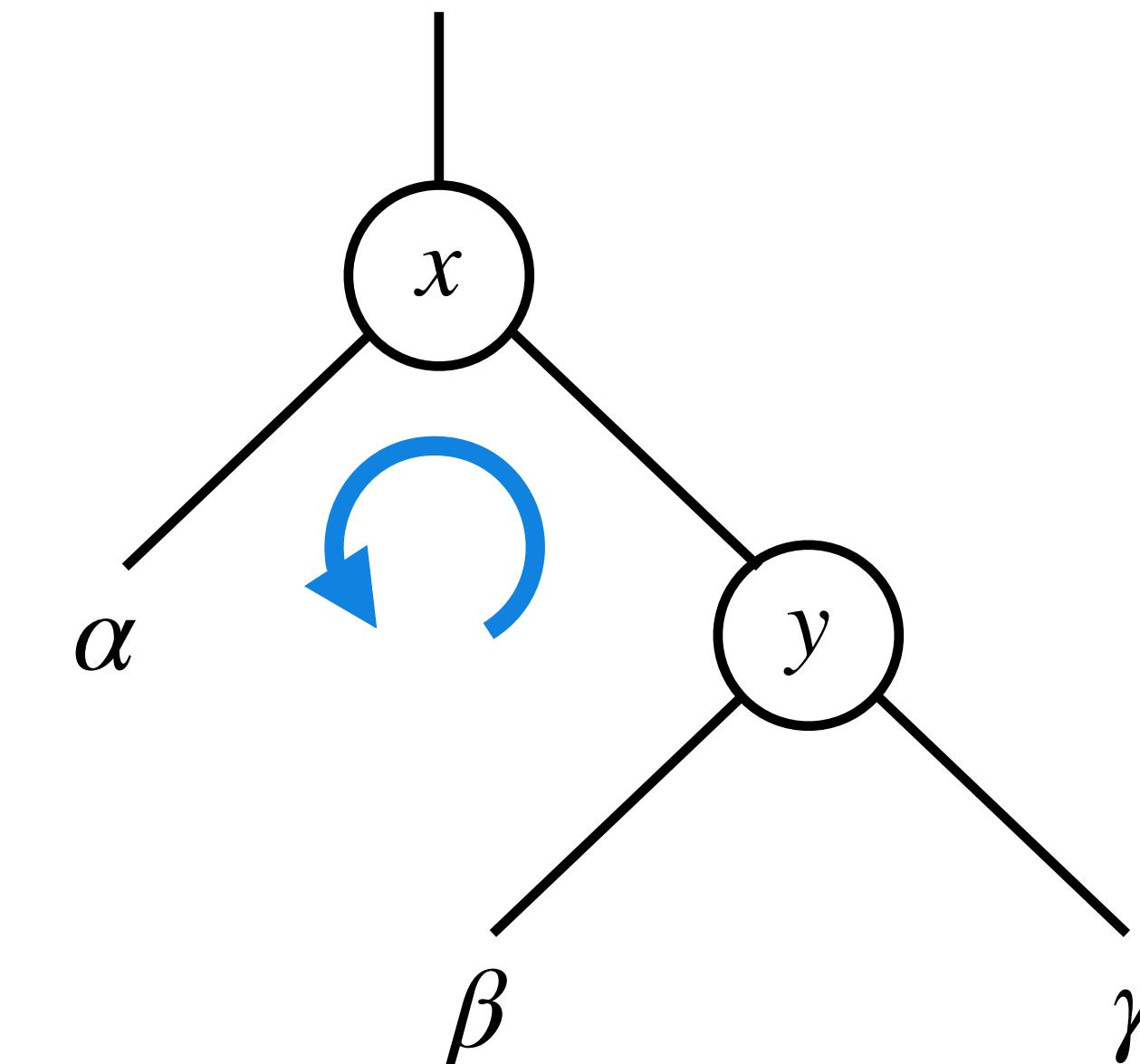
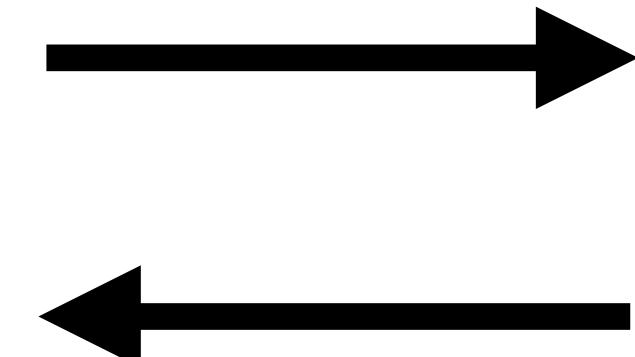
- 1) $y.size =$
- 2) $x.size =$

Maintaining Subtree Sizes: Rotations

Maintaining subtree sizes during rotations is easy.



Right-rotate(T, y)

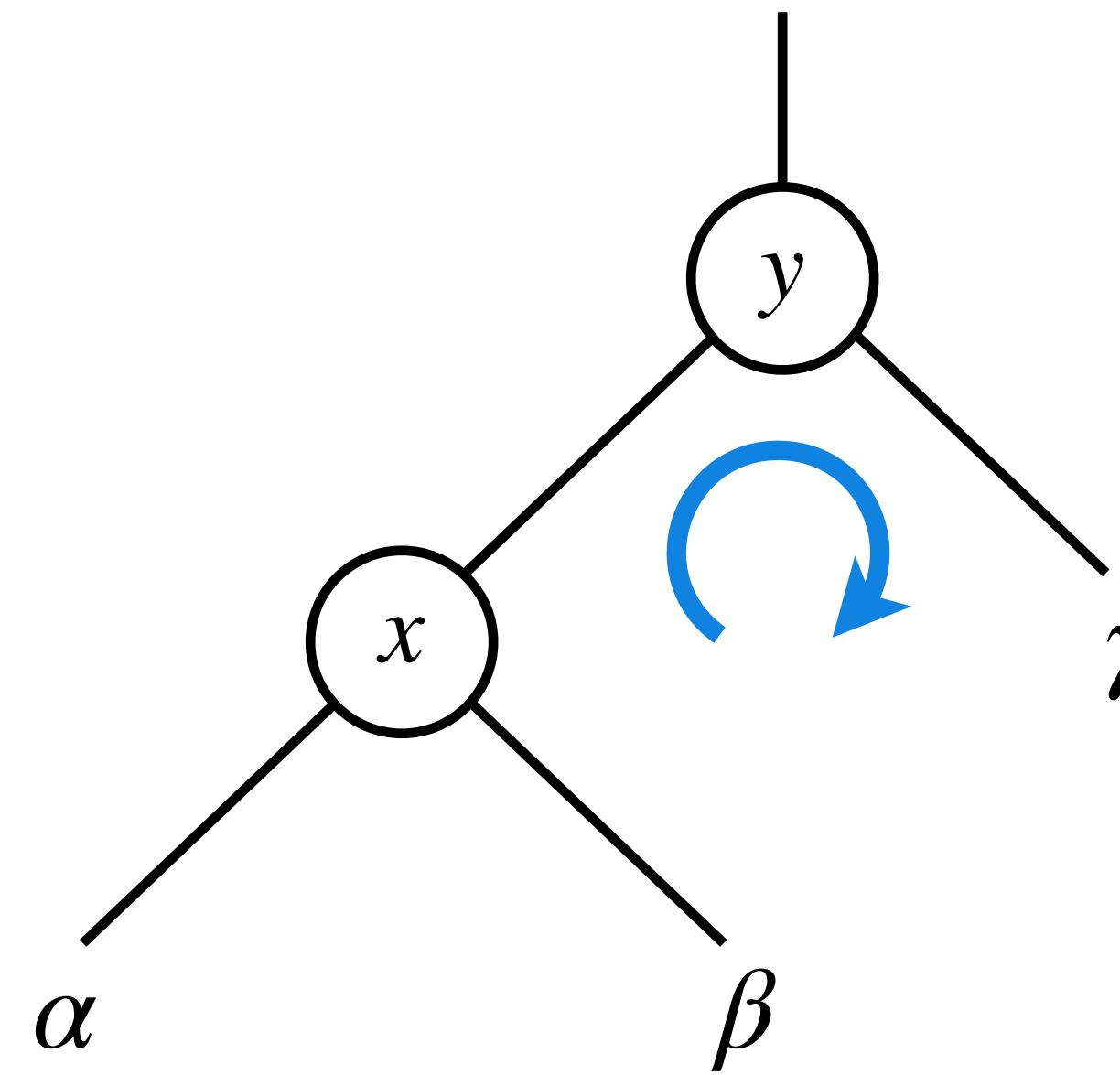


After $\text{Left-rotate}(T, x)$, do the following:

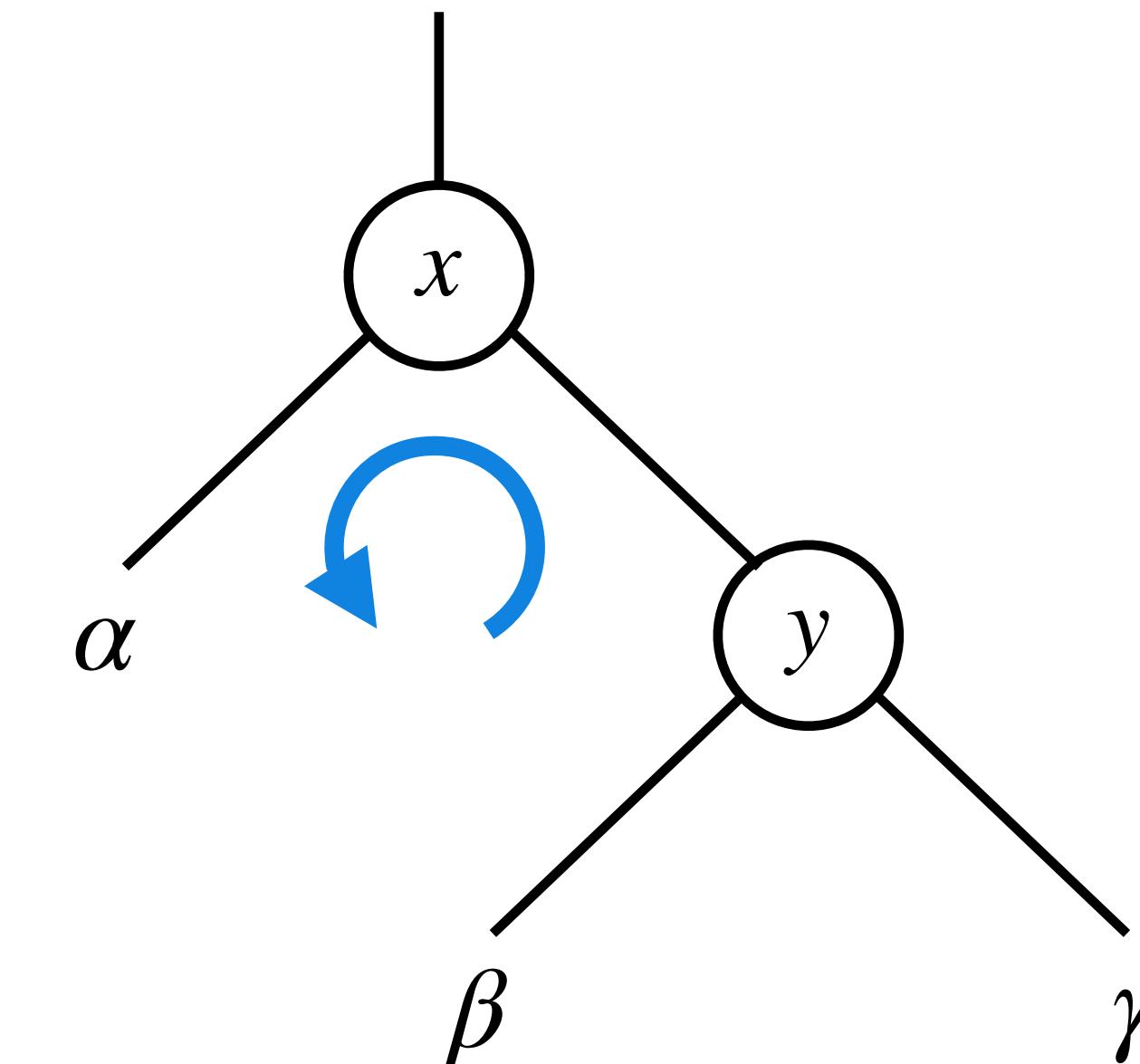
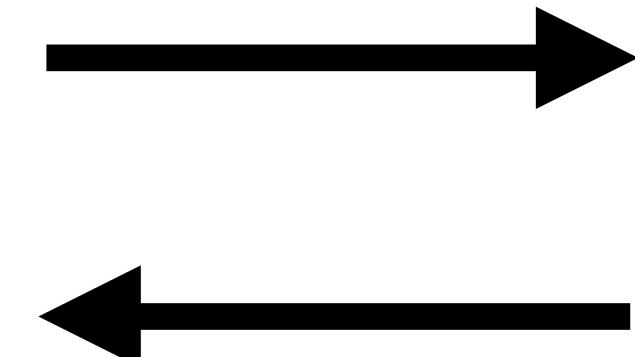
- 1) $y.size = x.size$
- 2) $x.size =$

Maintaining Subtree Sizes: Rotations

Maintaining subtree sizes during rotations is easy.



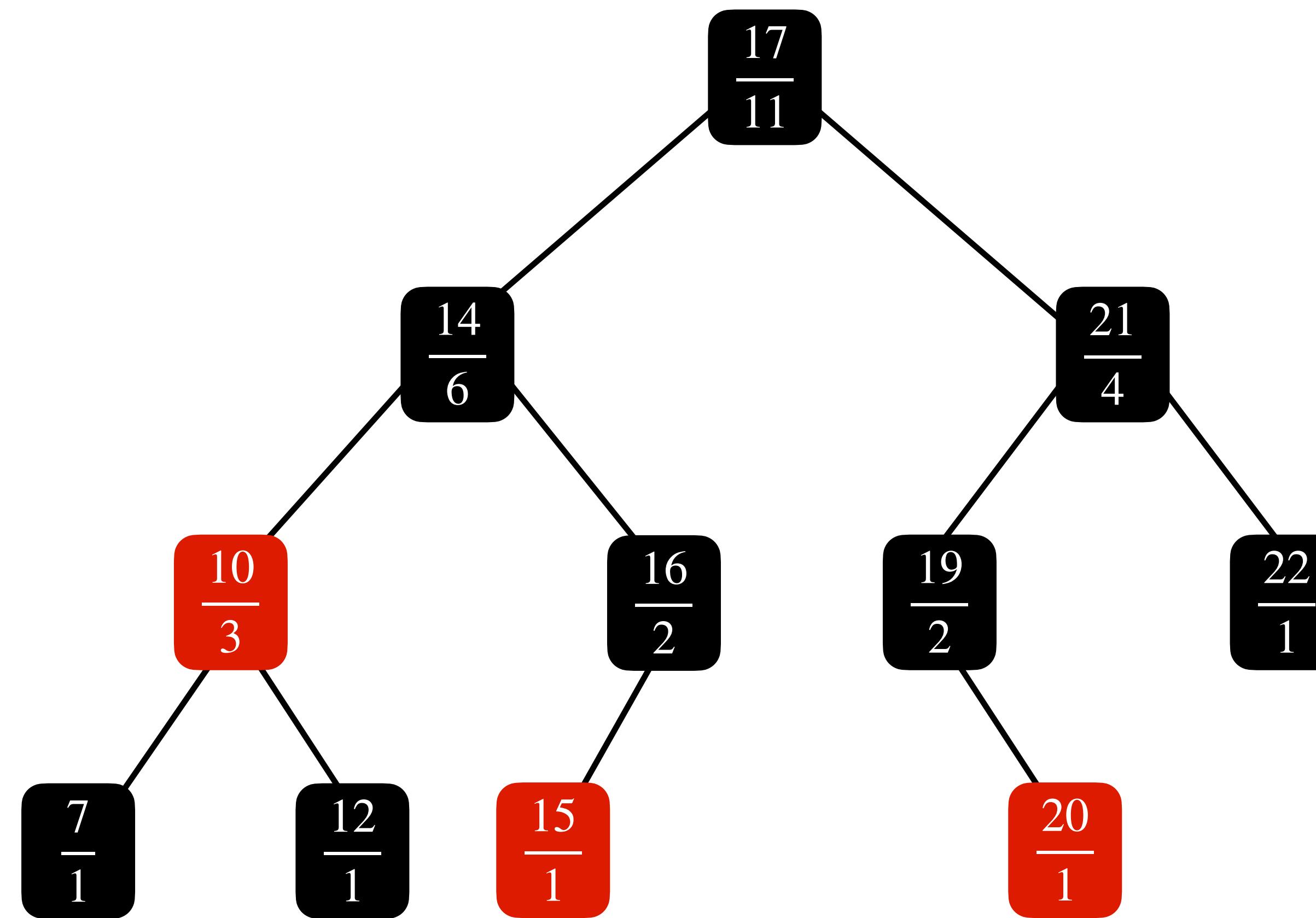
Right-rotate(T, y)



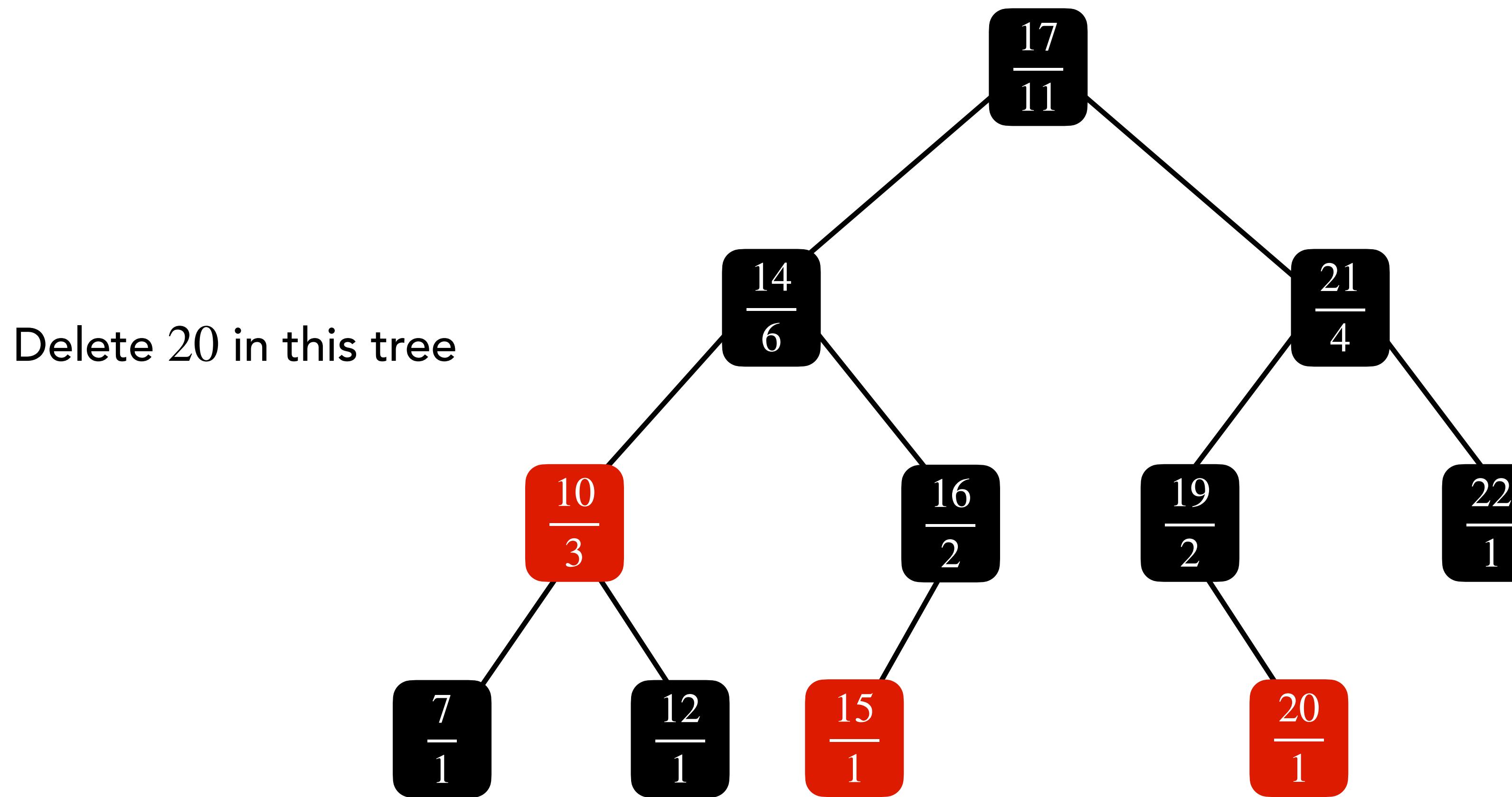
After $\text{Left-rotate}(T, x)$, do the following:

- 1) $y.size = x.size$
- 2) $x.size = x.left.size + x.right.size + 1$

Maintaining Subtree Sizes: Deletion

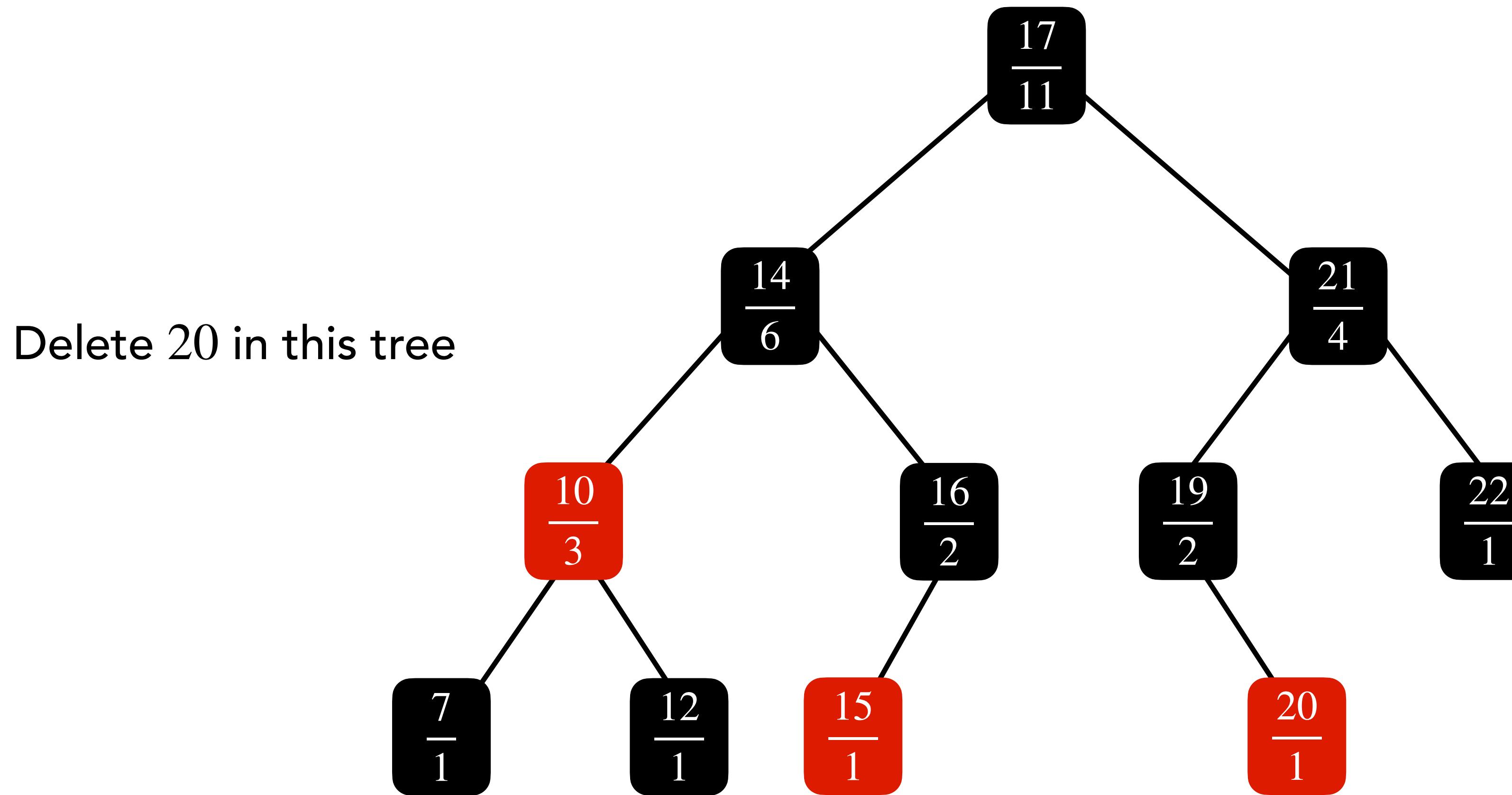


Maintaining Subtree Sizes: Deletion



Maintaining Subtree Sizes: Deletion

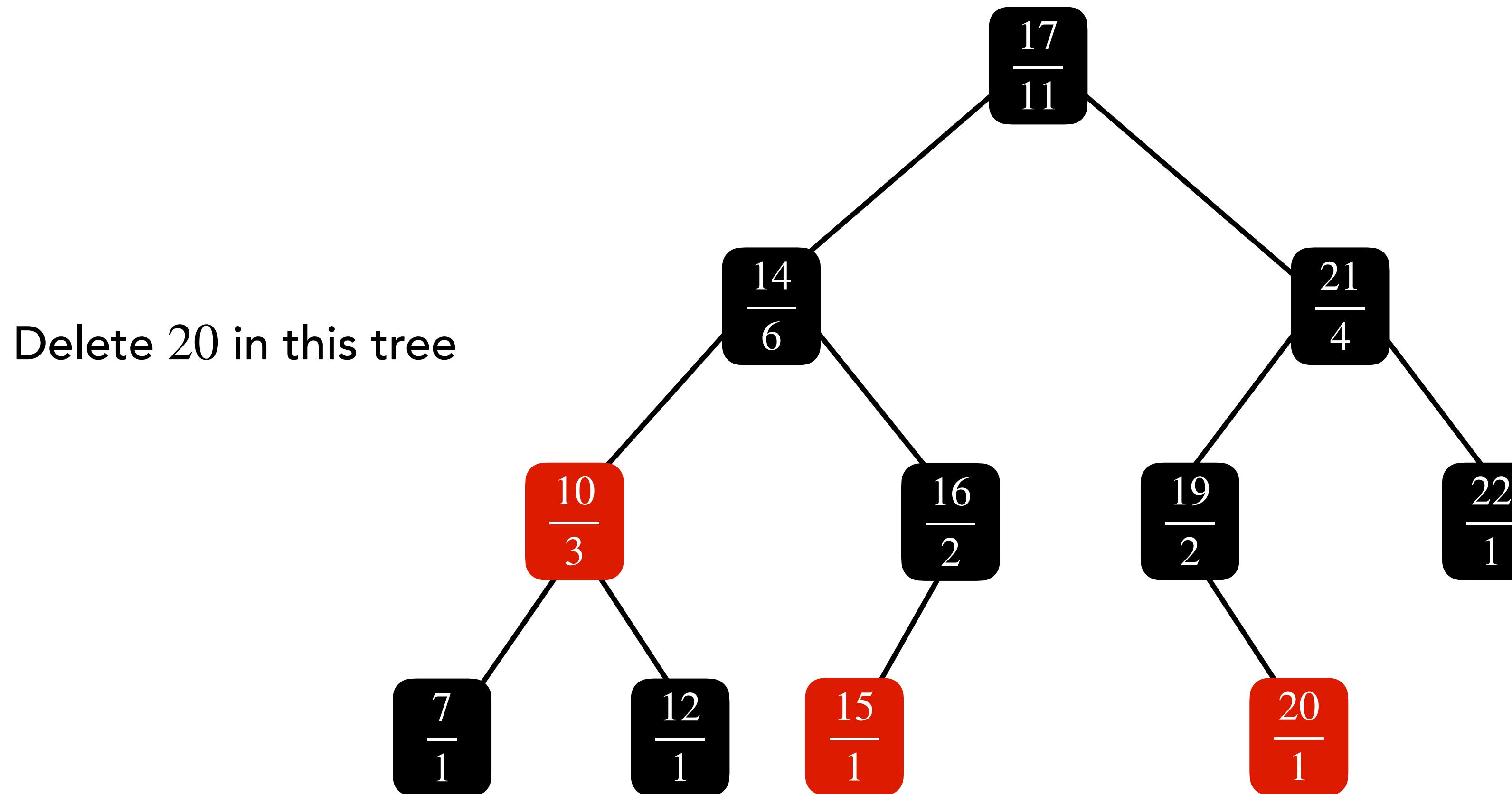
Deletion phase:



Maintaining Subtree Sizes: Deletion

Deletion phase:

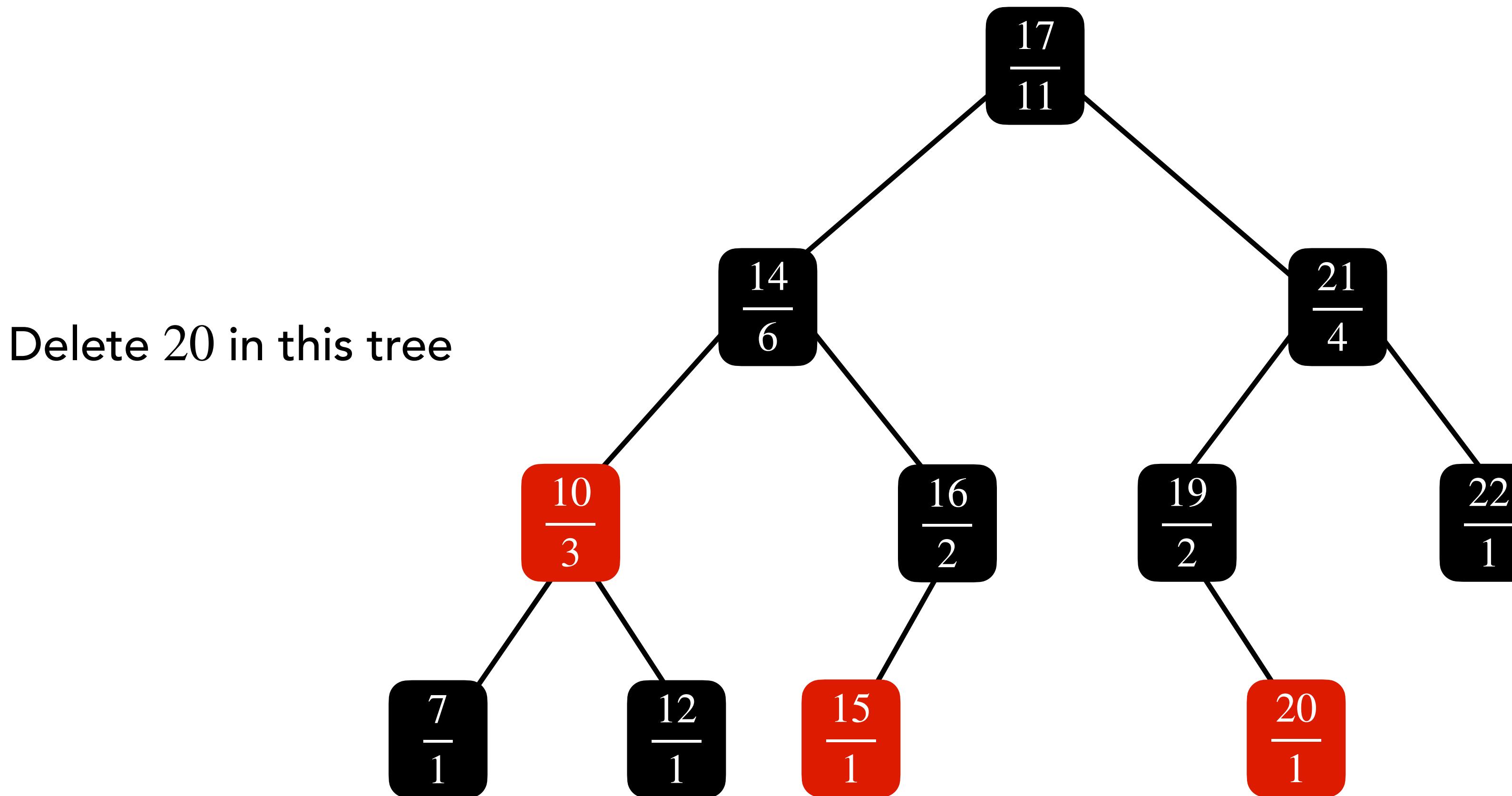
Keep subtracting 1 from the sizes of every node from the parent of the removed node to



Maintaining Subtree Sizes: Deletion

Deletion phase:

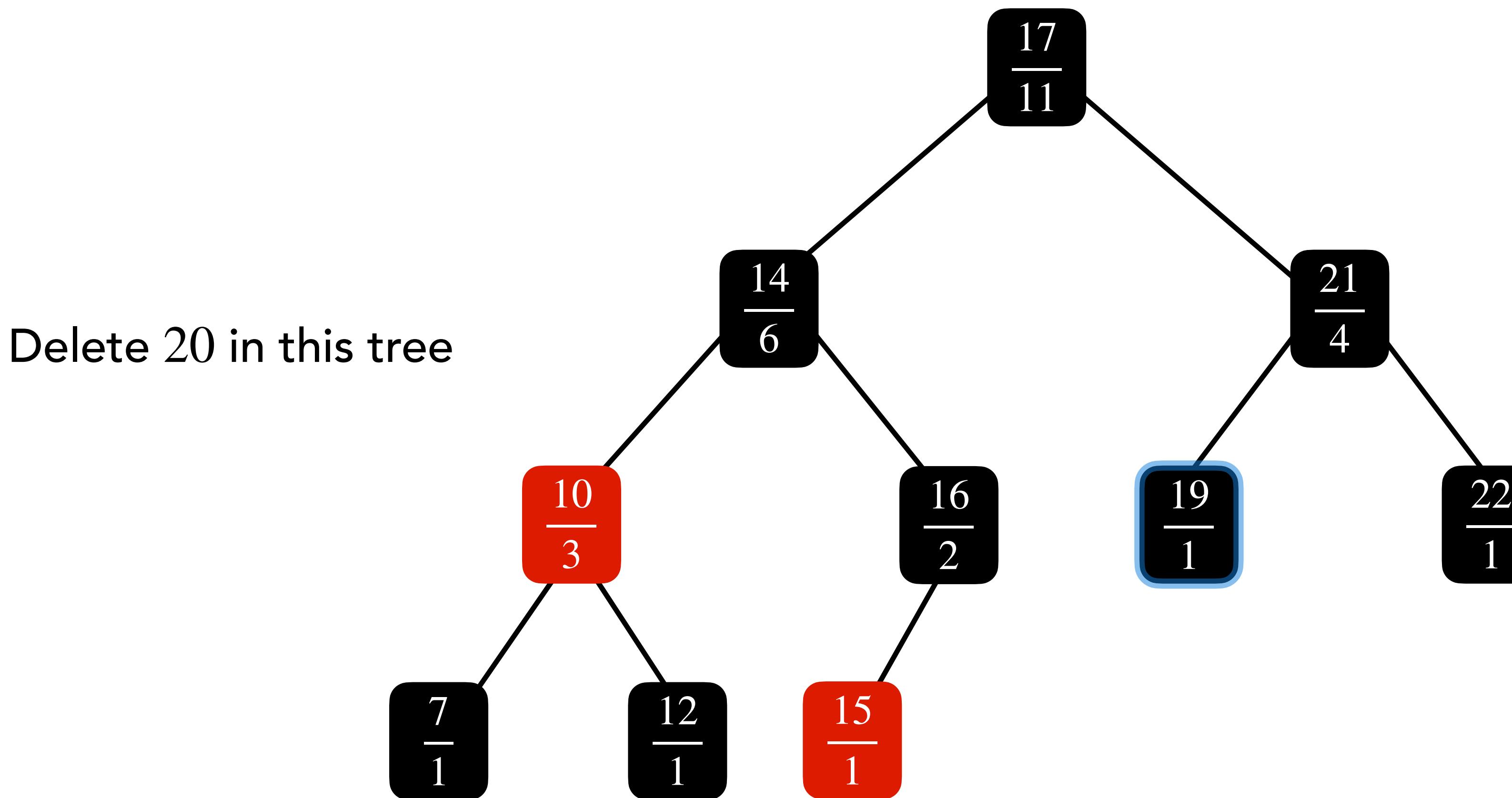
Keep subtracting 1 from the sizes of every node from the parent of the removed node to the root.



Maintaining Subtree Sizes: Deletion

Deletion phase:

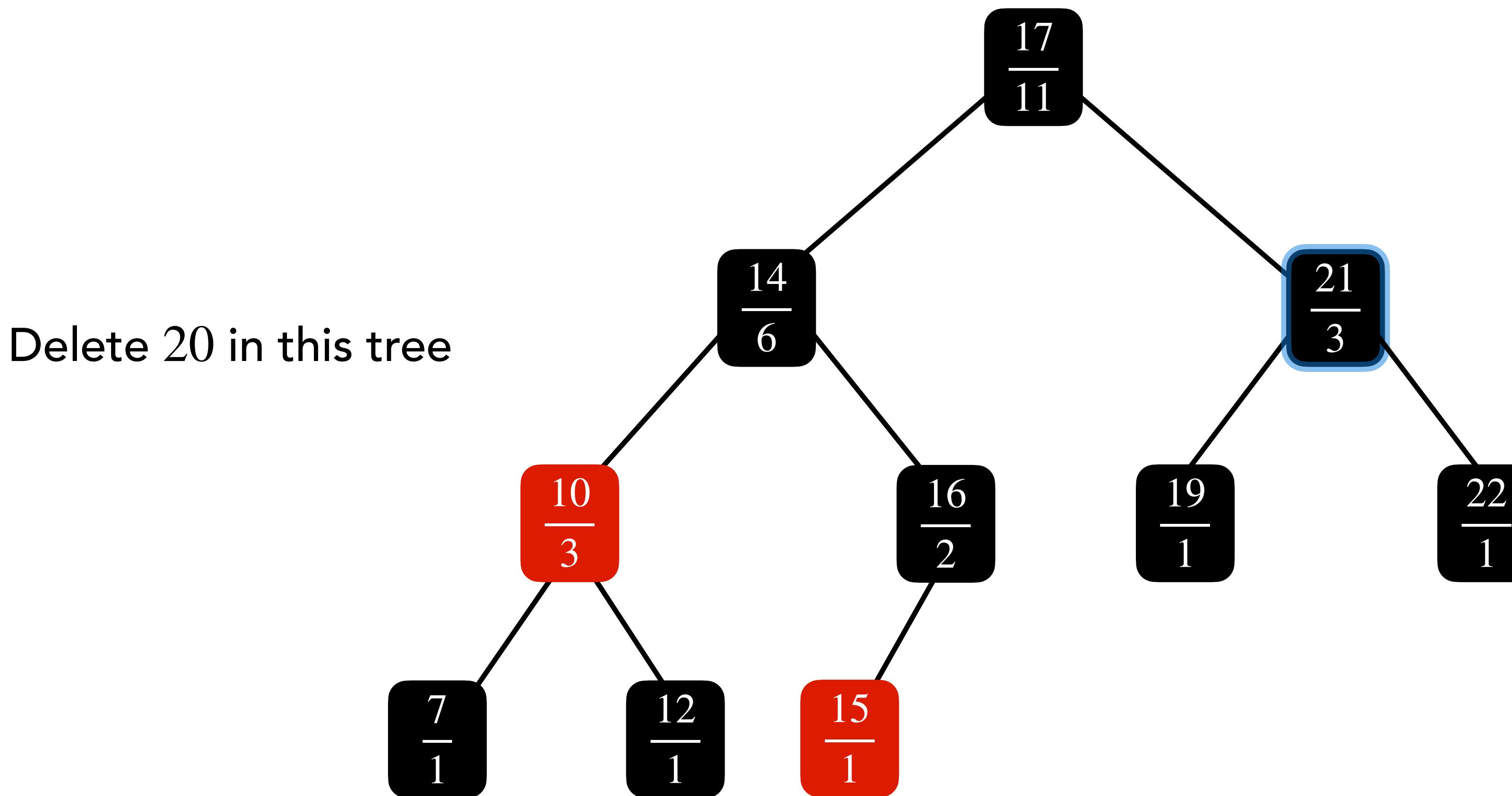
Keep subtracting 1 from the sizes of every node from the parent of the removed node to the root.



Maintaining Subtree Sizes: Deletion

Deletion phase:

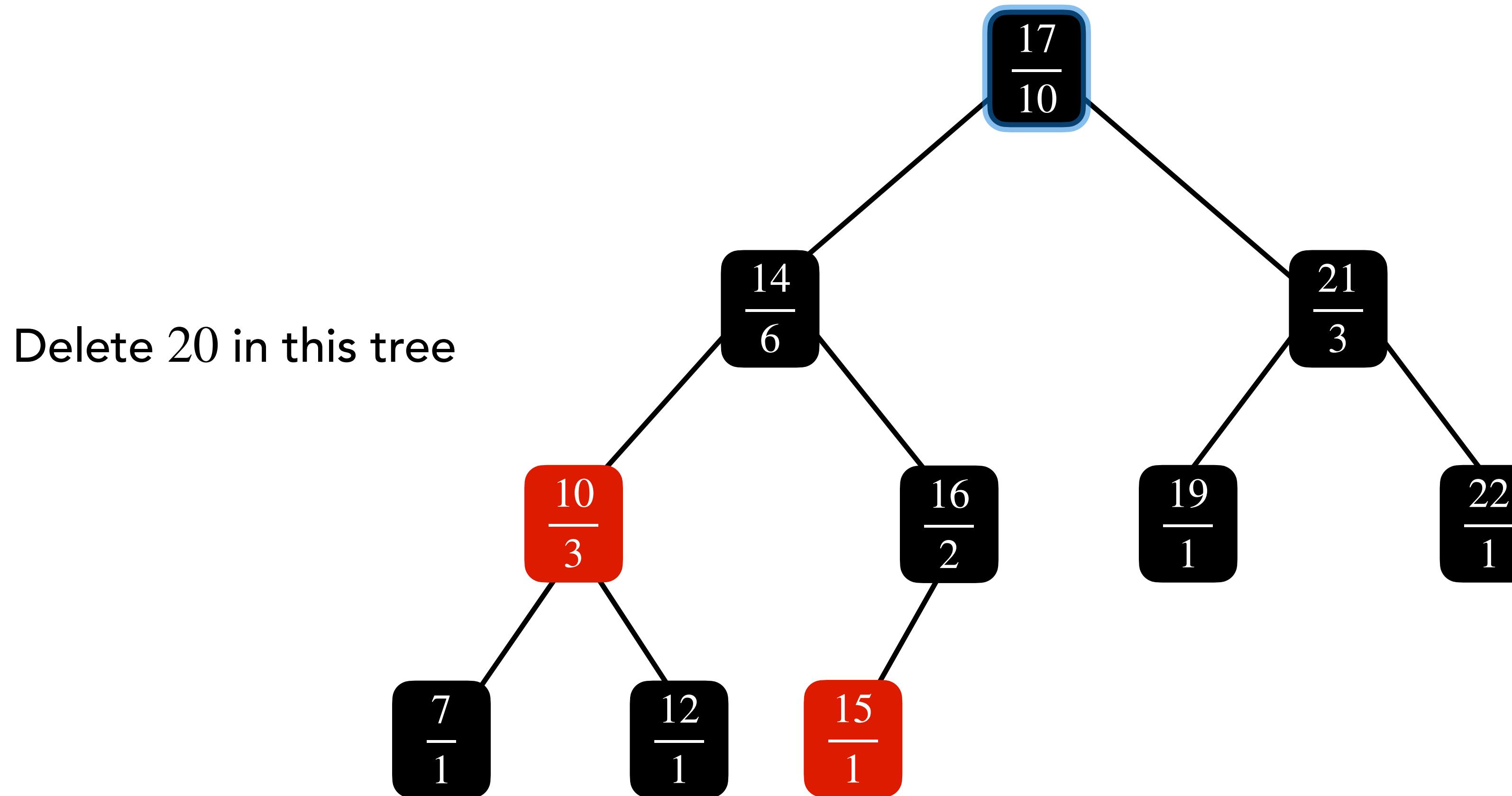
Keep subtracting 1 from the sizes of every node from the parent of the removed node to the root.



Maintaining Subtree Sizes: Deletion

Deletion phase:

Keep subtracting 1 from the sizes of every node from the parent of the removed node to the root.



Maintaining Subtree Sizes: Deletion

Maintaining Subtree Sizes: Deletion

Fix-up phase:

Maintaining Subtree Sizes: Deletion

Fix-up phase:

- Fix-ups involve only rotations and recolouring.

Maintaining Subtree Sizes: Deletion

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